## YALE UNIVERSITY DEPARTMENT OF COMPUTER SCIENCE

CPSC 467a: Cryptography and Computer Security

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## The Legendre and Jacobi Symbols

Let  $a \ge 0$ ,  $n \in \mathbb{Z}^+$ . Let QR(a, n) hold if (a, n) = 1 and a is a quadratic residue modulo n. Let QNR(a, n) hold if (a, n) = 1 and a is a quadratic non-residue modulo n (i.e., there is no  $y \in \mathbb{Z}_n^*$  such that  $a \equiv y^2 \pmod{n}$ .

For a prime p, the structure of quadratic residues can be fairly easily explained. Let g be a primitive root of  $\mathbf{Z}_p^*$ . Then every element of  $\mathbf{Z}_p^*$  is uniquely expressible as  $g^k$  for some  $k \in \{0, \dots, p-2\}$ .

**Theorem 1** Let p be a prime, g a primitive root of p,  $a \equiv g^k \pmod{p}$ . Then a is a quadratic residue iff k is even.

**Proof:** If k is even, then  $g^{k/2}$  is easily seen to be a square root of a modulo p.

Conversely, suppose  $a \equiv y^2 \pmod{p}$ . Write  $y \equiv g^{\ell} \pmod{p}$ . Then  $g^k \equiv g^{2\ell} \pmod{p}$ . Multiplying both sides by  $g^{-k}$ , we have  $1 \equiv g^0 \equiv g^{2\ell-k} \pmod{p}$ . But then  $\phi(p) | 2\ell - k$ . Since  $2 | \phi(p) = p - 1$ , it follows that 2 | k, as desired.

The following theorem, due to Euler, is now easily proved:

**Theorem 2** (Euler) Let p be an odd prime, and let  $a \ge 0$ , (a, p) = 1. Then

$$a^{(p-1)/2} \equiv \begin{cases} 1 \pmod{p} & \text{if } QR(a,p) \text{ holds;} \\ -1 \pmod{p} & \text{if } QNR(a,p) \text{ holds.} \end{cases}$$

**Proof:** Write  $a \equiv g^k \pmod{p}$ .

If QR(a, p) holds, then a is a quadratic residue modulo p, so k is even by Theorem 1. Write k = 2r for some r. Then  $a^{(p-1)/2} \equiv g^{2r(p-1)/2} \equiv (g^r)^{p-1} \equiv 1 \pmod{p}$  by Fermat's theorem.

If QNR(a, p) holds, then a is a quadratic non-residue modulo p, so k is odd by Theorem 1. Write k = 2r + 1 for some r. Then  $a^{(p-1)/2} \equiv g^{(2r+1)(p-1)/2} \equiv g^{r(p-1)}g^{(p-1)/2} \equiv g^{(p-1)/2} \pmod{p}$ . (mod p). Let  $b = g^{(p-1)/2}$ . Clearly  $b^2 \equiv 1 \pmod{p}$ , so  $b \equiv \pm 1 \pmod{p}$ . Since g is a primitive root modulo p and (p-1)/2 < p-1,  $b = g^{(p-1)/2} \not\equiv 1 \pmod{p}$ . Hence,  $b \equiv -1 \pmod{p}$ .

**Definition:** The Legendre symbol is a function of two integers a and p, written  $\left(\frac{a}{p}\right)$ . It is defined for  $a \ge 0$  and p an odd prime as follows:

$$\left(\frac{a}{p}\right) = \begin{cases} 1 \text{ if } QR(a, p) \text{ holds;} \\ -1 \text{ if } QNR(a, p) \text{ holds;} \\ 0 \text{ if } (a, p) \neq 1. \end{cases}$$

A multiplicative property of the Legendre symbols follows immediately from Theorem 1.

**Observation 3** Let  $a, b \ge 0$ , p an odd prime. Then

$$\left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right) \cdot \left(\frac{b}{p}\right).$$

<sup>&</sup>lt;sup>1</sup>This follows from the fact that  $p|(b^2 - 1) = (b - 1)(b + 1)$ , so either p|(b - 1), in which case  $b \equiv 1 \pmod{p}$ , or p|(b + 1), in which case  $b \equiv -1 \pmod{p}$ .

As an easy corollary of Theorem 2, we have:

**Corollary 4** Let  $a \ge 0$  and let p be an odd prime. Then

$$\left(\frac{a}{p}\right) \equiv a^{(p-1)/2} \pmod{p}$$

The Jacobi symbol extends the domain of the Legendre symbol.

**Definition:** The Jacobi symbol is a function of two integers a and n, written  $\left(\frac{a}{n}\right)$ , that is defined for all  $a \ge 0$  and all odd positive integers n. Let  $\prod_{i=1}^{k} p_i^{e_i}$  be the prime factorization of n. Then

$$\left(\frac{a}{n}\right) = \prod_{i=1}^{k} \left(\frac{a}{p_i}\right)^{e_i}.$$

Here  $\left(\frac{a}{p_i}\right)$  denotes the previously-defined Legendre symbol. (Note that by this definition,  $\left(\frac{0}{1}\right) = 1$ , and  $\left(\frac{0}{n}\right) = 0$  for odd  $n \ge 3$ .)

We have seen that if (a, p) = 1 and p is prime, then the Legendre symbol  $\left(\frac{a}{p}\right) = 1$  iff a is a quadratic residue modulo p. It is *not* true for the Jacobi symbol that  $\left(\frac{a}{n}\right) \equiv 1 \pmod{n}$  implies that a is a quadratic residue modulo n. For example,  $\left(\frac{8}{15}\right) = 1$ , but 8 is not a quadratic residue modulo 15. However, the converse does hold:

**Observation 5** If  $\left(\frac{a}{n}\right) \not\equiv 1 \pmod{n}$ , then a is not a quadratic residue modulo n.

The usefulness of the Jacobi symbol  $\left(\frac{a}{n}\right)$  stems from its ability to be computed efficiently, even without knowning the factorization of either *a* or *n*. The algorithm is based on the following theorem, which is stated without proof.

**Theorem 6** Let n be an odd positive integer,  $a, b \ge 0$ . Then the following identities hold:

- $1. \quad \left(\frac{0}{n}\right) = \begin{cases} 1 & \text{if } n = 1; \\ 0 & \text{if } n > 1 \end{cases}$   $2. \quad \left(\frac{2}{n}\right) = \begin{cases} 1 & \text{if } n \equiv \pm 1 \pmod{8}; \\ -1 & \text{if } n \equiv \pm 3 \pmod{8} \end{cases}$   $3. \quad \left(\frac{a}{n}\right) = \left(\frac{b}{n}\right) \text{ if } a \equiv b \pmod{n}.$   $4. \quad \left(\frac{ab}{n}\right) = \left(\frac{a}{n}\right) \cdot \left(\frac{b}{n}\right)$
- 5. (Quadratic reciprocity). If a is odd, then  $\begin{pmatrix} \frac{a}{n} \end{pmatrix} = \begin{cases} -\binom{n}{a} & \text{if } a \equiv n \equiv 3 \pmod{4}; \\ \binom{n}{a} & \text{otherwise.} \end{cases}$

Theorem 6 leads directly to a recursive algorithm for computing  $\left(\frac{a}{n}\right)$ :

```
int jacobi(int a, int n)
/* Precondition: a, n >= 0; n is odd */
{
 int ans;
 if (a == 0)
   ans = (n==1) ? 1 : 0;
 else if (a == 2) {
   switch (n%8) {
   case 1:
   case 7:
     ans = 1;
     break;
   case 3:
   case 5:
    ans = -1;
     break;
   }
  }
 else if (a \ge n)
   ans = jacobi(a%n, n);
 else if (a\%2 == 0)
   ans = jacobi(2,n) * jacobi(a/2, n);
 else
   ans = (a%4 == 3 && n%4 == 3) ? -jacobi(n,a) : jacobi(n,a);
 return ans;
}
```