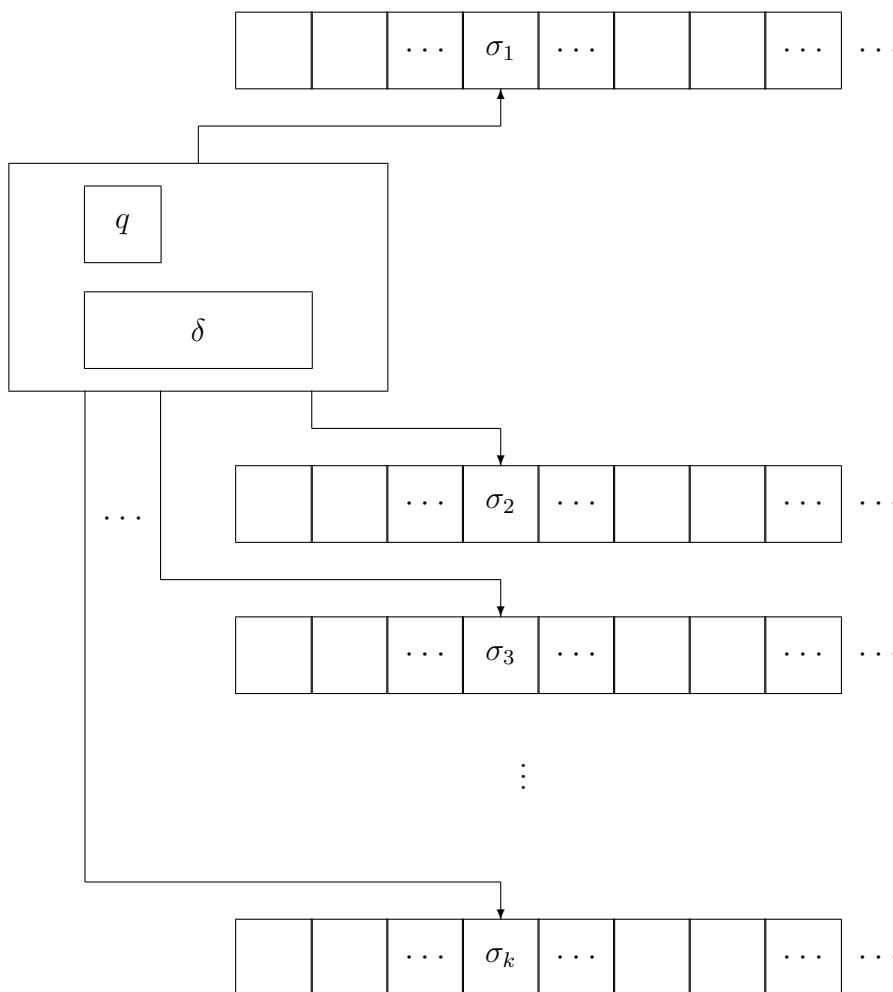


A Turing-Machine model of Computation

Deterministic k -tape Turing machine M .



There is one read-only **input tape** (on top) and $k - 1$ read-write **work/output tapes**. M is a triple Γ, Q, δ that is defined as follows:

- Γ is the **tape alphabet**, a finite set of symbols. Assume \square ("blank" symbol), \triangleright ("start" symbol), 0 and 1 are four distinct elements of Γ .
- Q is the **state set**, a finite set of states that M 's control register can be in. Assume q_{start} and q_{halt} are two distinct states in Q .
- δ is the **transition function**, a finite table that describes the rules (or program) by which M operates:

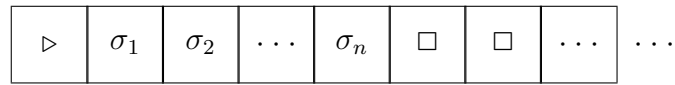
$$\delta : Q \times \Gamma^k \rightarrow Q \times \Gamma^{k-1} \times (L, S, R)^k.$$

$\delta(q, (\sigma_1, \dots, \sigma_k)) = (q', (\sigma'_2, \dots, \sigma'_k), (z_1, \dots, z_k))$ means that, if M is in state q , and the read (or read/write) tape heads are pointing at the cells containing $\sigma_1, \dots, \sigma_k$, then the following "step" of the computation is performed:

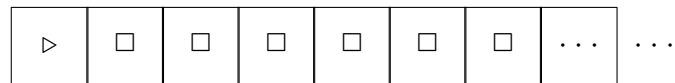
- the read/write tape symbols $\sigma_2, \dots, \sigma_k$ are replaced by $\sigma'_2, \dots, \sigma'_k$;
- tape head i moves left, stays in place or moves right, depending on whether z_i is in L, S or R ;
- the control-register state is changed to q' .

When M starts its execution on input $x = \sigma_1, \dots, \sigma_n$, we have

- $q = q_{\text{start}}$
- input tape



- all other tapes



Meaning of q_{halt} :

$$\delta(q_{\text{halt}}, (\sigma_1, \dots, \sigma_k)) = (q_{\text{halt}}, (\sigma_2, \dots, \sigma_k), S^k) \quad \forall (\sigma_1, \dots, \sigma_k).$$

Designate one of the read/write tapes as "the output tape".

Turing machine M "computes the function f ", if for all $x \in \Gamma^*$ the execution of M on input x eventually reaches the state q_{halt} , and when it does, the contents of M 's output tape is $f(x)$.

M "runs in time T " if for all n and all $x \in \Gamma^n$ M halts after at most $T(n)$ steps.