CPSC 468/568: Lecture 7 (Sept. 20, 2012)

This lecture began with definitions of implicitly logspace computable, logspace reducibility, NL completeness, and read-once certificates and with the proof that PATH is NL-complete. See Chapter 4 of Arora-Barak.

The Immerman-Szelepcsenyi Theorem: NL = coNL

Proof:

Throughout, points that you are encouraged to think through and justify in detail are marked by **"(WHY?)."**

Recall first that PATH is NL-complete and, equivalently, that $\overline{\text{PATH}}$ is coNL-complete. Thus, it suffices to show that $\overline{\text{PATH}}$, the set of triples (G, s, t) in which G is a directed graph that does *not* contain a path from s to t, is in NL.

We will do so by exhibiting a deterministic, logspace verifier that takes as input both an instance (G, s, t) and a certificate of this instance's membership in $\overline{\text{PATH}}$. As usual, the tape on which the instance is written is read-only. New to our discussion of nondeterministic logspace is the requirement that the tape on which the certificate is written is not just read-only but *read-once, left-to-right.* The work/output tapes of this machine are, as usual, read/write, and they are the only tapes that are restricted to logspace.

If $V(G) = \{1, 2, \ldots, n\}$, and G is encoded on the input tape as an $n \times n$ adjacency matrix, then the input is of length $O(n^2)$. Thus, we need certificates of length $poly(n^2) = poly(n)$ and space complexity $O(\log(O(n^2))) = O(\log n)$.

Let $C_i = \{v \in V(G)$ such that v is reachable from s by a path of length at most $i\}$. Note that $C_0 = \{s\}$ and that C_n contains all nodes in G that are reachable from s by any path whatsoever. **(WHY?)** The desired certificate that (G, s, t) is in $\overline{\text{PATH}}$ must therefore certify the fact that $t \notin C_n$. It comprises three types of "subcertificates," as follows.

$\text{CERT}_1(v, i, q_i)$ proves that $v \notin C_i$, given that $|C_i| = q_i$.

$\text{CERT}_2(v, i, q_{i-1})$ proves that $v \notin C_i$, given that $|C_{i-1}| = q_{i-1}$.

$\text{CERT}_3(i, q_i, q_{i-1})$ proves that $|C_i| = q_i$, given that $|C_{i-1}| = q_{i-1}$.

Overall, to prove that $(G, s, t) \notin \overline{\text{PATH}}$, we use the certificate

$$\text{CERT}_3(1, q_1, 1)\text{CERT}_3(2, q_2, q_1) \cdots \text{CERT}_3(n, q_n, q_{n-1})\text{CERT}_1(t, n, q_n).$$

That is, starting with the obvious fact that $|C_0| = 1$, the logspace verifier first checks, for each successive i, $2 \leq i \leq n$, that $|C_i| = q_i$; once it has checked that $|C_n| = q_n$, it checks that t is not one of the q_n nodes in C_n. If each of the constituent subcertificates is polynomial-length and logspace verifiable in a read-once, left-to-right manner, then so is the entire certificate. **(WHY?)**

$\text{CERT}_1(v, i, q_i)$ is a list of q_i paths to all of the nodes reachable from s along paths of length at most i. If we denote by $\ell(1), \ldots, \ell(q_i)$ the lengths of these paths, then this subcertificate has the form:

$$\langle u_1^1 u_2^1 \dots u_{\ell(1)}^1 \rangle \langle u_1^2 u_2^2 \dots u_{\ell(2)}^2 \rangle \cdots \langle u_1^{q_i} u_2^{q_i} \dots u_{\ell(q_i)}^{q_i} \rangle,$$

where $u_1^j = s$, for $1 \leq j \leq q_i$, and $u_{\ell(1)}^1 < u_{\ell(2)}^2 < \cdots < u_{\ell(q_i)}^{q_i}$. That is, all of the paths start at s, and we list them in increasing order of the labels of their terminal vertices.

It suffices for the deterministic, logspace verifier to check the following. **(WHY?)**

- The total number of paths in the subcertificate is q_i.
- v is not in any of the paths.
- For $1 \leq j \leq q_i - 1$, $u_{\ell(j)}^j < u_{\ell(j+1)}^{j+1}$.
- The arcs $u_k^j \to u_{k+1}^j$ are all in $E(G)$.
- $s = u_1^j$, for $1 \leq j \leq q_i$.
- $\ell(j) \leq i$, for $1 \leq j \leq q_i$.

Indeed, all of these conditions can be verified in deterministic logspace in a read-once, left-to-right manner. **(WHY?)**

$\mathrm{CERT}_2(v, i, q_{i-1})$ is the same as $\mathrm{CERT}_1(v, i-1, q_{i-1})$, but its verification procedure contains one more condition: The arc $u_{\ell(j)}^j \to v$ is *not* in $E(G)$ for any j, $1 \leq j \leq q_{i-1}$. That is, if one knows all of the nodes that can be reached by paths of length at most $i-1$, checking that v is not reachable in one step from any of them suffices to show that v cannot be reached by a path of length at most i.

Finally, $\mathrm{CERT}_3(i, q_i, q_{i-1})$ consists of n subcertificates $D_1, \dots, D_n$. If $j \in C_i$, then D_j is a path from s to j of length at most i. If $j \notin C_i$, then $D_j = \mathrm{CERT}_2(j, i, q_{i-1})$. The verifier reads them all left-to-right, checks their validity, and checks that exactly q_i of them certify that the vertex j in question is a member of C_i.