

Complexity

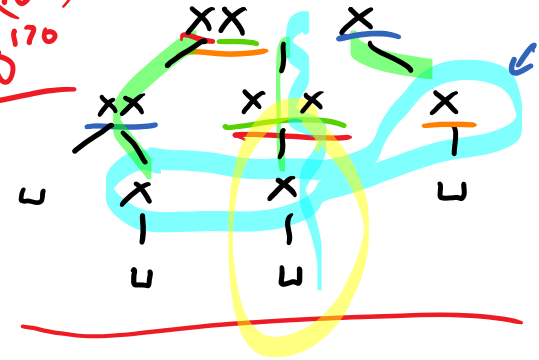
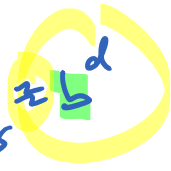
[Game complexity - Wikipedia](#)

State Space - # of arrangements of pieces

checks $\leq 5^{32} \approx 10^{20}$
50 $\approx 3^{361} \approx 10^{170}$

10^{10} ok

Game tree size - # of nodes in game tree
sequences of moves



Branching factor
 $b = \#$ of moves available

Depth
 $d =$ length of game

Backtracking

for finite puzzles (so no cycles)

```
find_solution(s)
```

```
  if s is solved position return []
```

```
  for every possible move m
```

```
    s' <- result of move m at state s    make move m on s
```

```
    solution = find_solution(s')
```

```
    if solution is not NIL
```

```
      return [m] + solution
```

```
    else
```

```
      pass undo move m on s
```

```
  return NIL
```

```

g[s], f[s] <- infinity for all states s
g[initial] <- 0, f[initial] <- h(initial)
Q <- priority queue(states, f[])

```

$g(s)$ = cost of best path so far $init \rightarrow s$
 $f(s)$ = est of best path $init \rightarrow s \rightarrow goal$

```

while Q not empty
  curr <- Q.extractMin()
  if curr is goal
    return goal, pi[goal], pi[pi[goal]], ..., initial
  else
    for every state n reachable from curr
      d <- g[curr] + cost(curr, n)
      if d < g[n] then
        if g[n] = infinity then
          f[n] <- d + h(n)
        else
          f[n] <- f[n] - (g[n] - d)
          g[n] <- d
          pi[n] = curr
        if not Q.contains(n) then
          f[n] <- d + h(n)
          Q.add(n, f[n])
        else
          Q.decreasePriority(n, f[n])
return NIL

```

$$IND: f[n] = g[n] + h(n)$$

$$BASIS: f[i] = h(i) = h(i) + 0 = h(i) + g[i]$$

$$f[s] = \infty$$

$$g[s] = \infty = \infty + h(s)$$

IND: outer if $F \rightarrow$ no change

1st inner if $T \rightarrow$

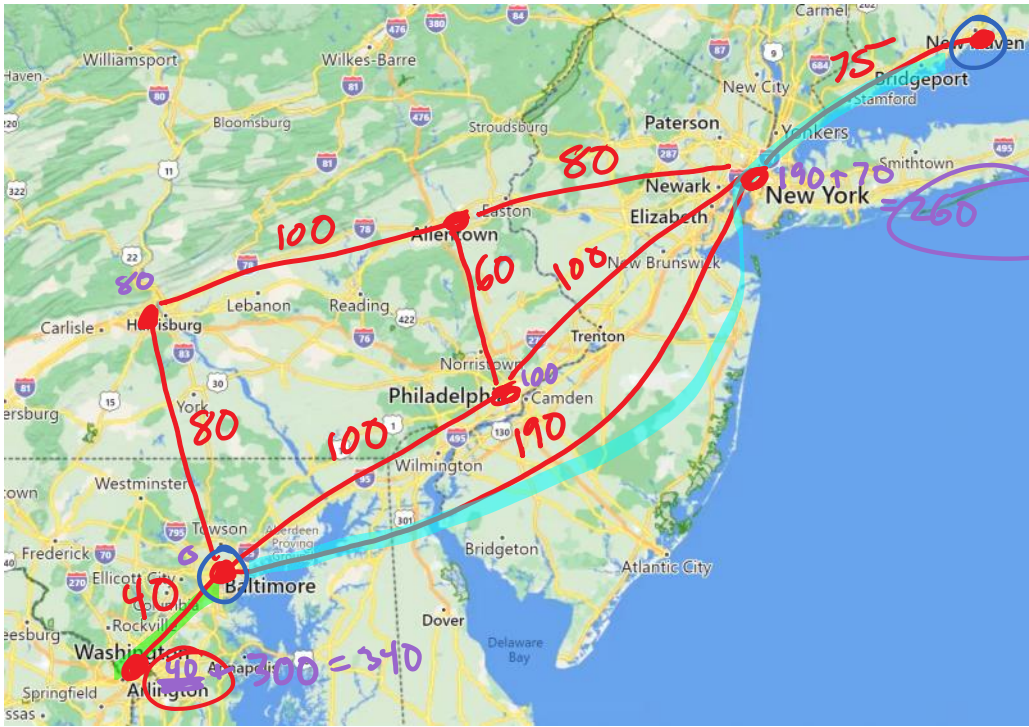
$$f[n] = d + h(n) = g[n] + h(n)$$

1st inner if $F \rightarrow$

$$\begin{aligned}
f_{new}[n] &= f_{old}[n] - g_{old}[n] + d \\
&= g_{old}[n] + h(n) - g_{old}[n] + d \\
&= d + h(n) \\
&= g_{new}[n] + h(n)
\end{aligned}$$

heuristic h is admissible if it never overestimates

if h is admissible then A^* returns optimal path

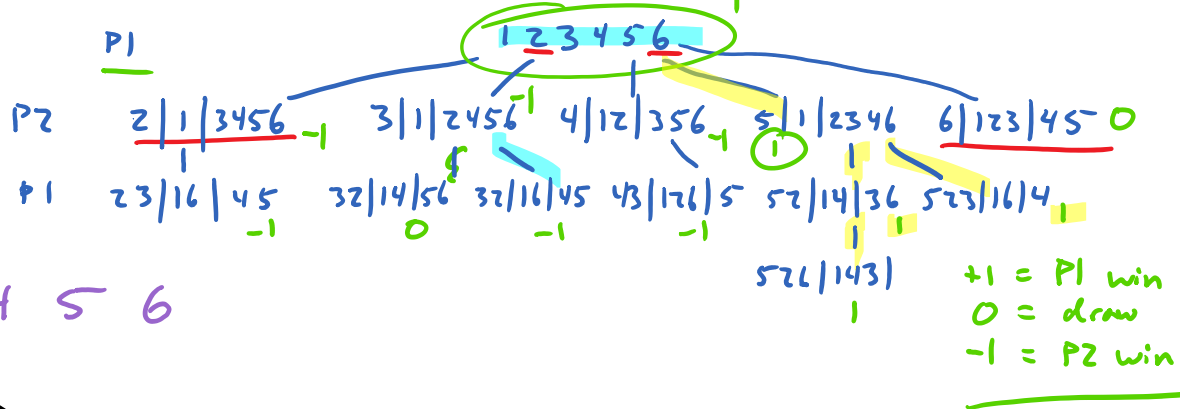


$h(s)$ = great circle distance to goal

15-puzzle: $h(s) = \sum_{t \neq s} \text{distance from cur loc of } t \text{ to its position in goal}$

Finite Combinatorial Games

Divisors : Start with 1...n, players take turns taking a number with remaining divisors; opponent gets all the remaining divisors. (Game is over when no moves remain; winner is player with higher sum (draw if =))



Minimax (s)

if s is end of game
return value according to rules
else

S ← states reachable in 1 move from s

if s is P1's max return $\max_{s' \in S} \text{Minimax}(s')$
else return $\min_{s' \in S} \text{Minimax}(s')$

Graph Game

Graph: take turns coloring a vertex in a graph with your color
player who covers the most edges wins (draw if =)

