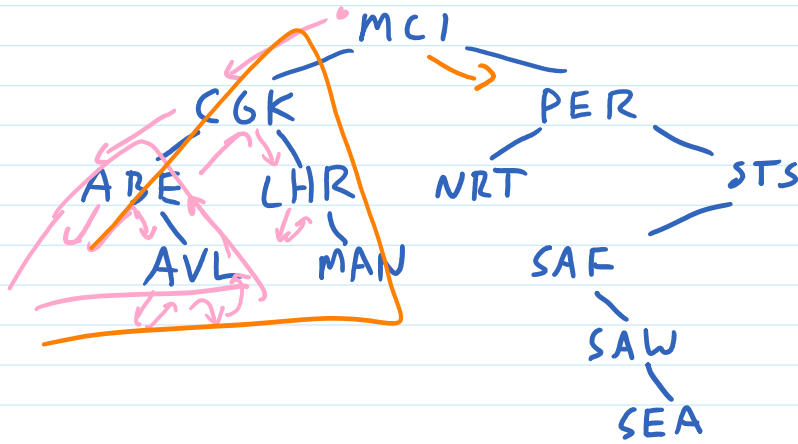


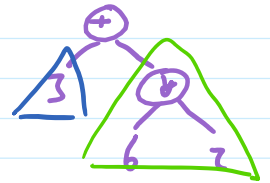
Inorder Traversal

left
root
right



ABE
AVL
CGK
LHR
MAN
MCI
⋮
Sorted!

post order on
expression tree
 $3 + 6 * 2$



post order 3 6 2 +

stack

$$\begin{array}{r} \times \\ 6 \quad 12 \\ 3 \quad 8 \quad 15 \\ + \quad * \end{array}$$

Red-Black Trees

Full Binary Search Trees such that

all nodes
0 or 2
children

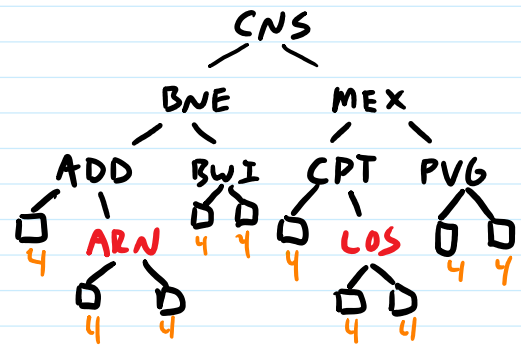
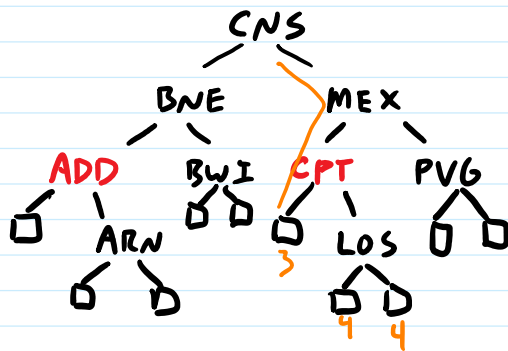
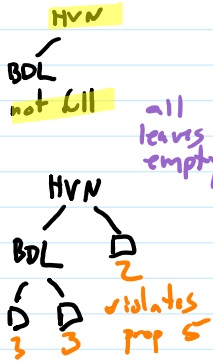
1) each node has a color: red or black

2) root is colored black

3) all leaves are colored black

4) children of red nodes are black

5) every path root \rightarrow leaf has same # of black nodes
black height

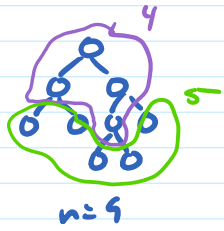
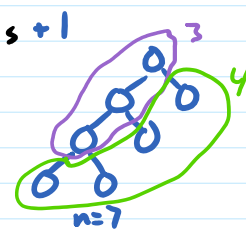


leaves in a non-empty full binary tree = # internal nodes + 1

$k = \text{total nodes}$
 $n = \text{data nodes}$

$$k = 2n + 1$$

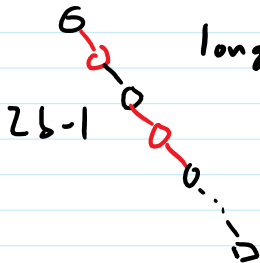
if time was $O(\log 2n+1)$
 $= O(\log n)$



Suppose # black nodes in root \rightarrow leaf path is b

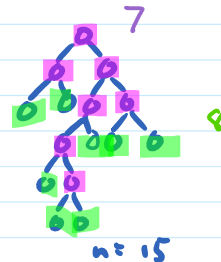
shortest possible path has b nodes

longest possible path has $2b - 1$



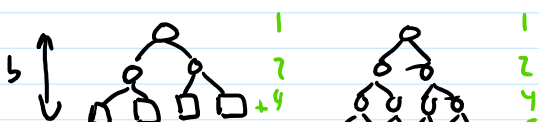
$$h \leq 2b - 1$$

h is $O(b)$



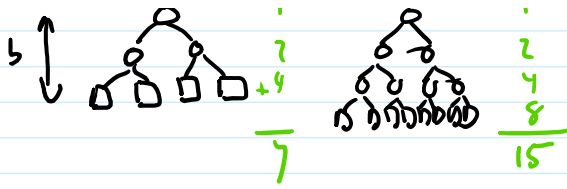
internal (data) nodes = # red nodes
prove by structural induction

Red-black tree with black height b has $\geq 2^b - 1$ nodes



$$k \geq 2^b - 1$$

$$k + 1 \geq 2^b$$



$$k \geq 2^h - 1$$

$$k+1 \geq 2^h$$

$$\log_2 k+1 \geq h$$

h is $O(\log k)$

and $O(\log n)$

1) Do normal BST insert, color new node **red** (with empty black leaves)

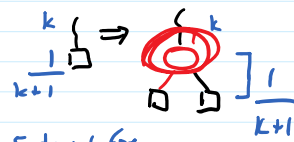
2) If new node's parent exists and is black, DONE

Else

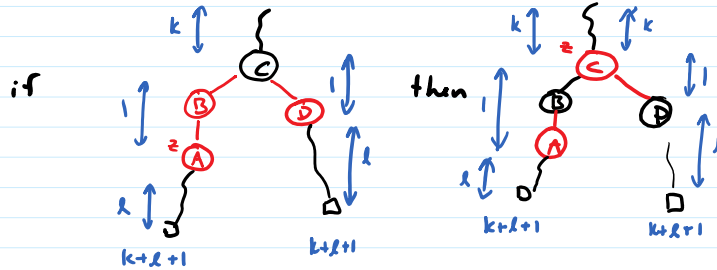
$z \leftarrow$ node added

red node we're worried about

while z is not root and $z \rightarrow \text{parent} \rightarrow \text{color} == \text{red}$

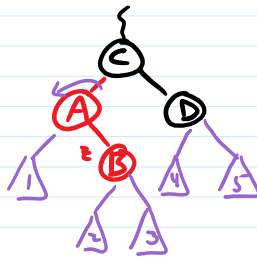


prop 5 tree before \rightarrow tree after

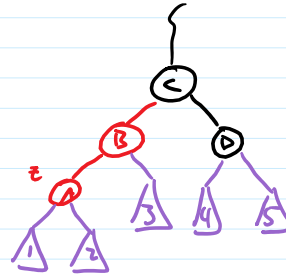


else

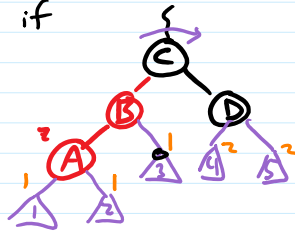
if



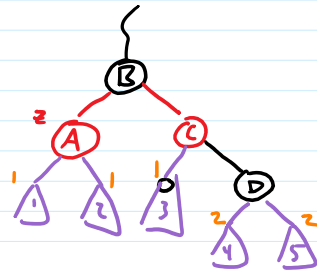
then



if



then



STOP!
(no more red links)

[other cases symmetric]

$C \rightarrow \text{root} \rightarrow \text{color} \leftarrow \text{black}$