

Priority Queue

enqueue (elt, pri)

↳ adds elt w/ priority

dequeue ()

↳ removes elt w/ highest max priority
(or lowest) min queue

change-priority (elt, pri)
change priority of existing item

unsorted array

$O(1)$

sorted array

$O(n)$

balanced BST

$O(\log n)$

heap

$O(\log n)$

$O(n)$

$O(1)$

$O(\log n)$

$O(\log n)$

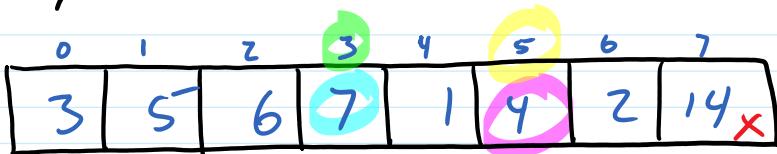
$\underline{\mathcal{O}(1)}$

$O(n)$

$O(\log n)$

$O(\log n)$

unsorted array

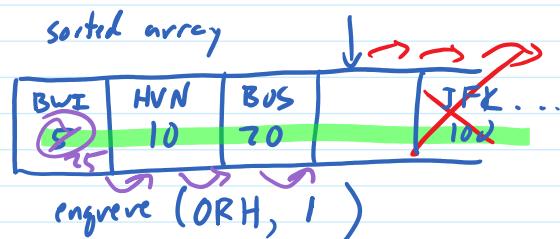


elts are indices $0, \dots, n-1$

enqueue (3, 7)

enqueue (BWI, 4)

key
BWI
value
5



Priority Queue Sort

$O(1) \cdot n$ 1) add each element to priority queue priority = key to sort on $O(\log n) \cdot n$ balanced BST or heap

$O(n) \cdot n$ 2) while p.g. not empty

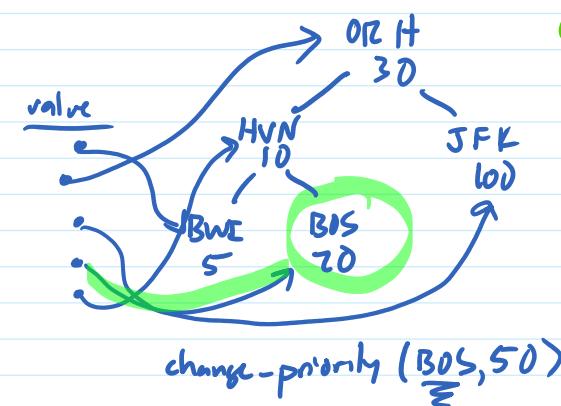
$O(n^2)$

remove min/max

when queue is unsorted array

III selection sort

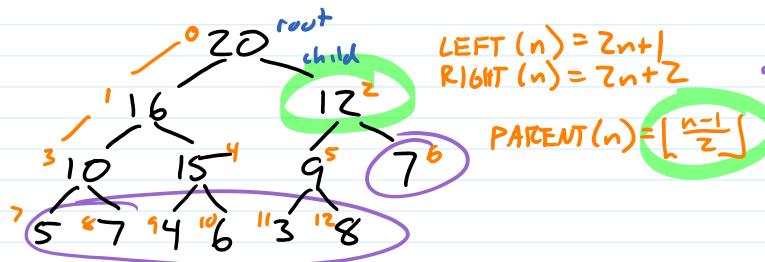
key
BWI
ORH
JFK
BUS
HVN



$O(\log n) \cdot n$

$O(n \log n)$

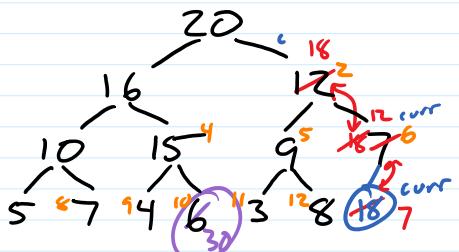
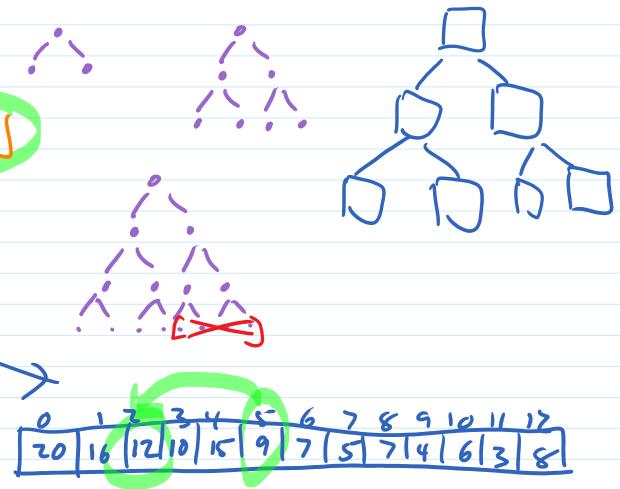
TreeSort
or
HeapSort



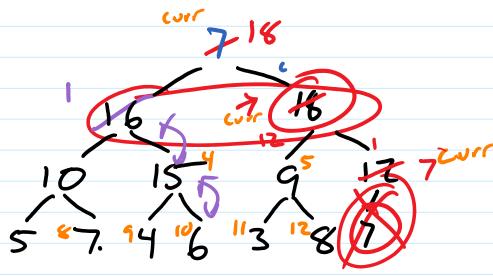
binary tree s.t.

shape tree is nearly complete
 ↳ leaves as far down/left as possible - leaves on last 1 or 2 levels and as far left as possible on bottom level

order value in node is \geq values in children max-heap
 \leq min-heap



enqueue(18, 18)



dequeue

enqueue: add at next available spot
 $O(\log n)$ → while curr not at root and curr > parent
 d1) [swap (curr, parent)
 move curr up one level]

heap-reheap-up

dequeue: get max from root $O(1)$

move last leaf to root $O(1)$

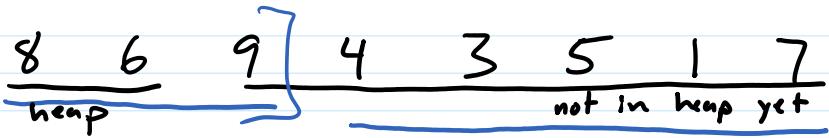
$O(\log n)$ while curr not leaf and curr < one child
 d2) [swap w/ largest child
 move curr down to where swap was]

$O(\log n)$ total

heap-reheap-down

Heapsort

- 1) build heap from array
- 2) extract max repeatedly from array



8
6 9

4 3 5 1

7

