

Course Logistics:

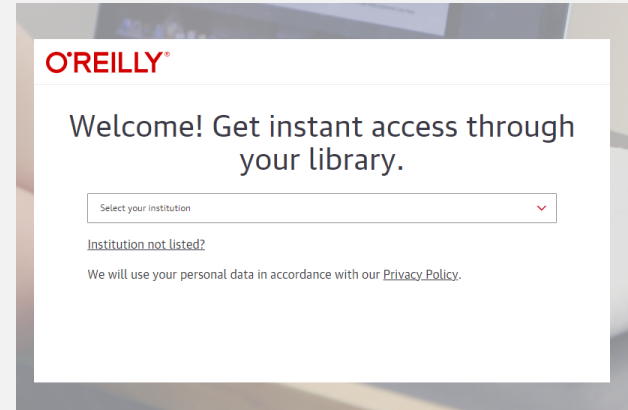
Have you checked your exam?

You have free access to O'Reilly Library

<https://www.oreilly.com/library-access/?orpg>

Good visualization tool:

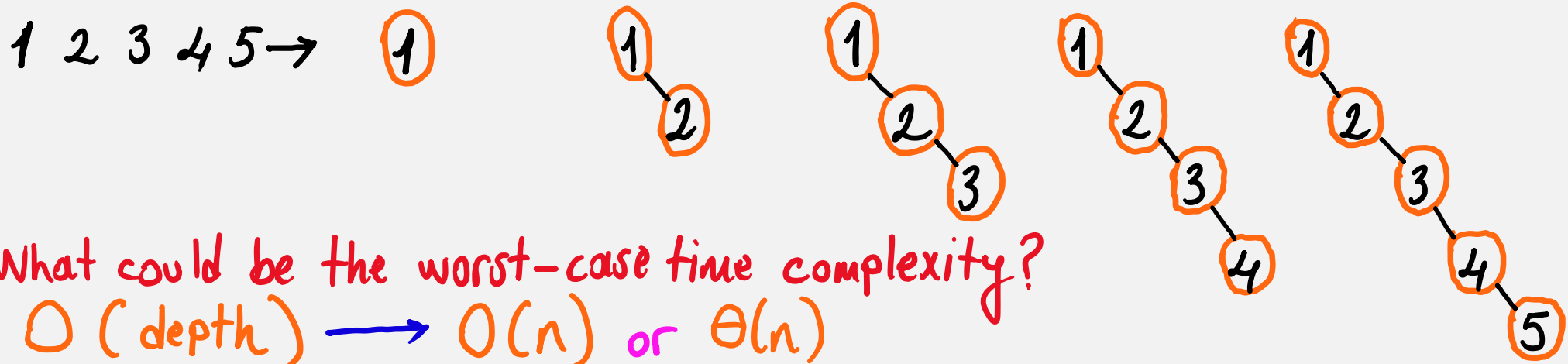
<https://www.cs.usfca.edu/~galles/visualization/Algorithms.html>



Binary Search Trees

How unbalanced could the tree be?

↳ already sorted input, reverse-sorted, nearly sorted ...



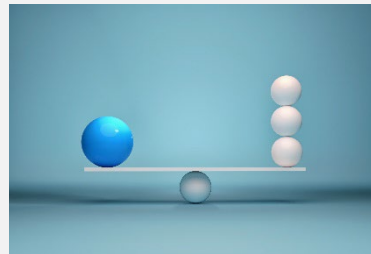
• What could be the worst-case time complexity?

$O(\text{depth}) \rightarrow O(n)$ or $\Theta(n)$

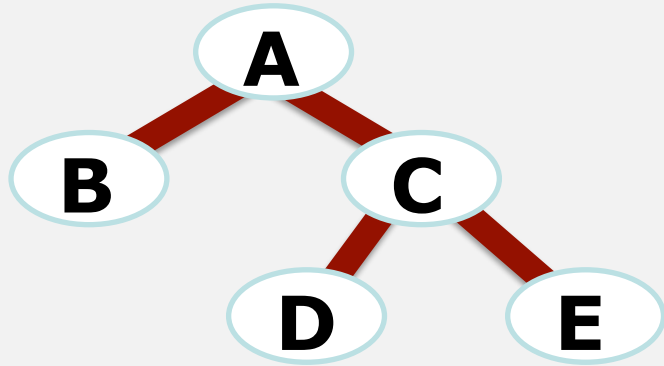
• What would be ideal? $O(\log n)$ or $\Theta(\log n)$

AVL Trees

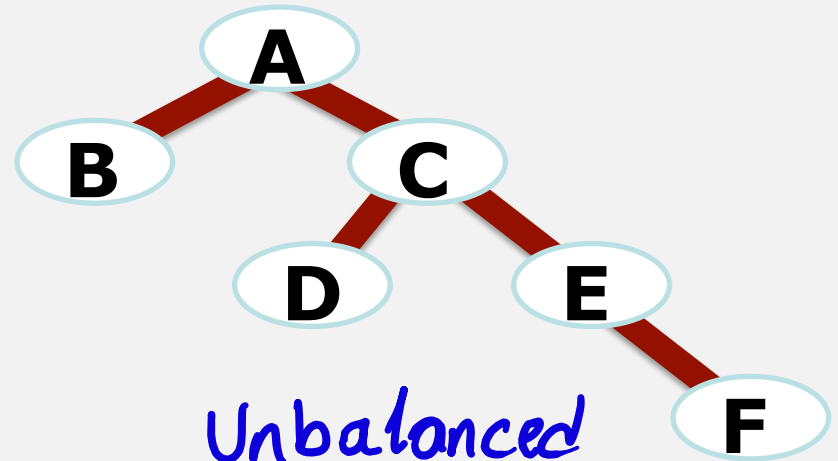
- A data structure, invented by **Adelson-Velsky** and **Landis** in 1962 to balance **Binary Search Trees**.
- Essentially **AVL** trees are **BSTs** with one extra invariant (rule).
 - ↳ tree is always balanced.
- What does **balanced** mean?
 - ↳ What is the criteria for that?



Which one is unbalanced? Who decides that?

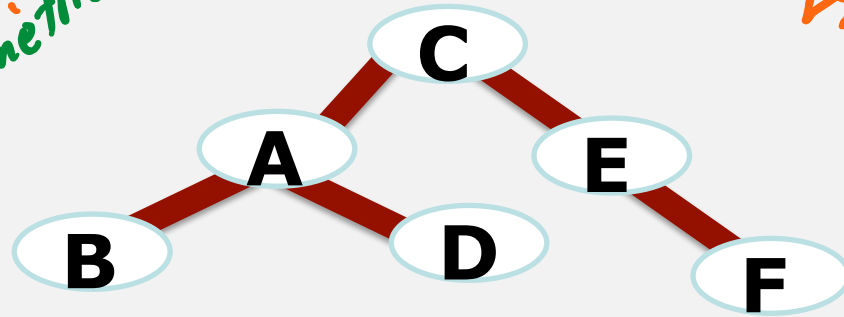


Balanced



Unbalanced

*Is it a symmetry thing?
None of them are symmetric.*

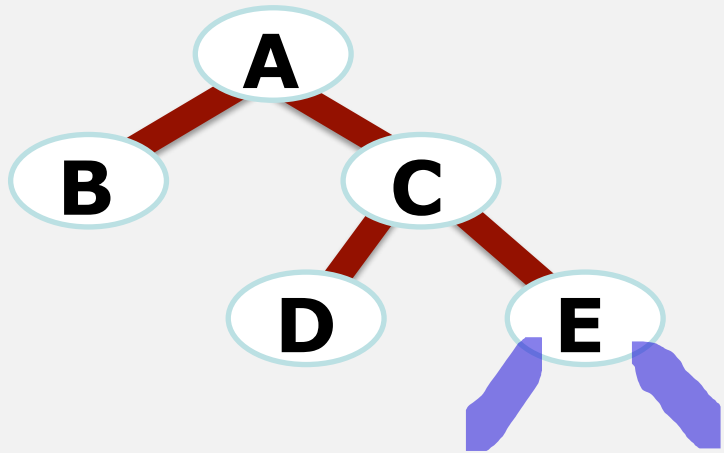


Balanced

*Visually pleasing?
To whom?*

We need a deterministic approach!

Height



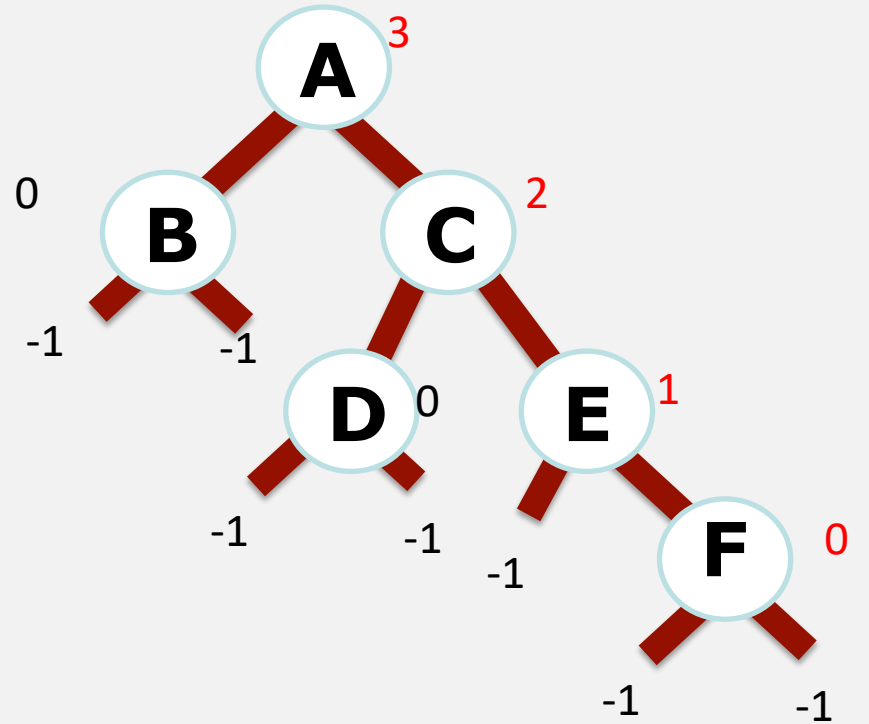
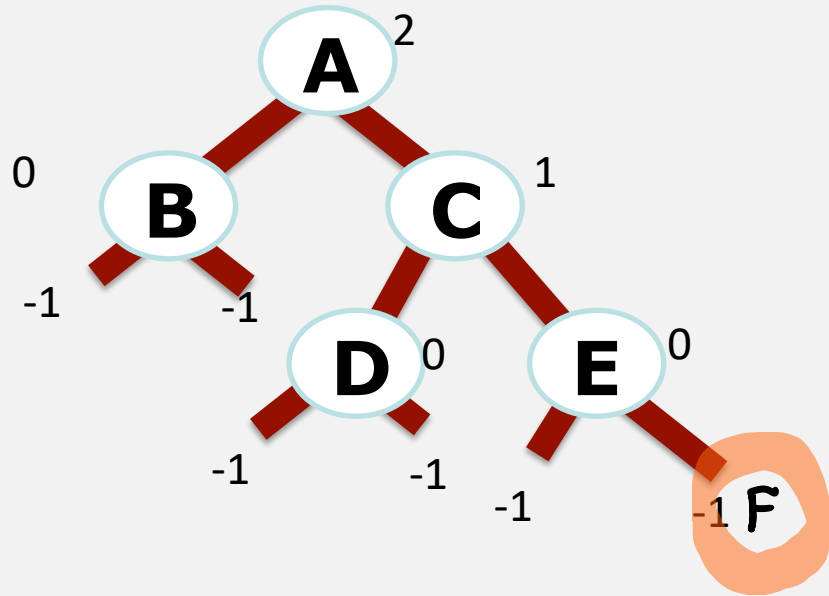
- Height of an empty node
↳ non-existing child of
node E → -1

- Height of a node
↳ longest shortest path to a descendant leaf

example: Node E → 0
Node C → 1
Node A → 2
Node B → 0

- Height of a tree
↳ The leaf which is furthest away from the root.
example: starting from E or D to A → 2

Computing Height:



- Recursion:

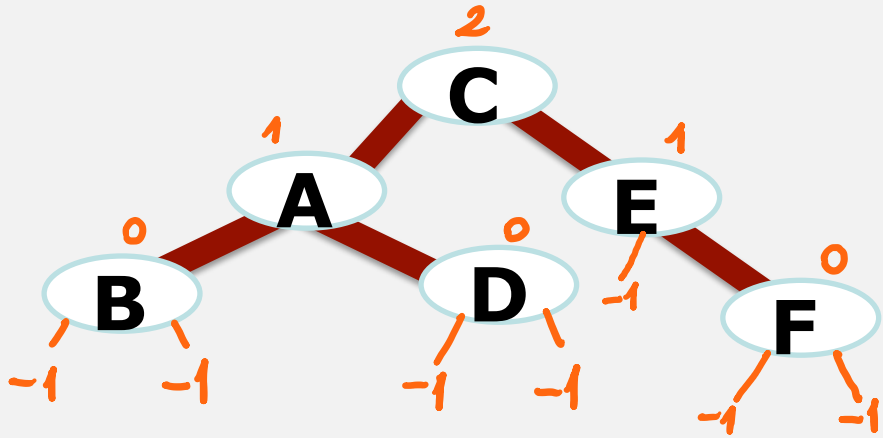
$\text{height}(\text{node}) = \max(\text{height of children}) + 1;$

How can we maintain the height?

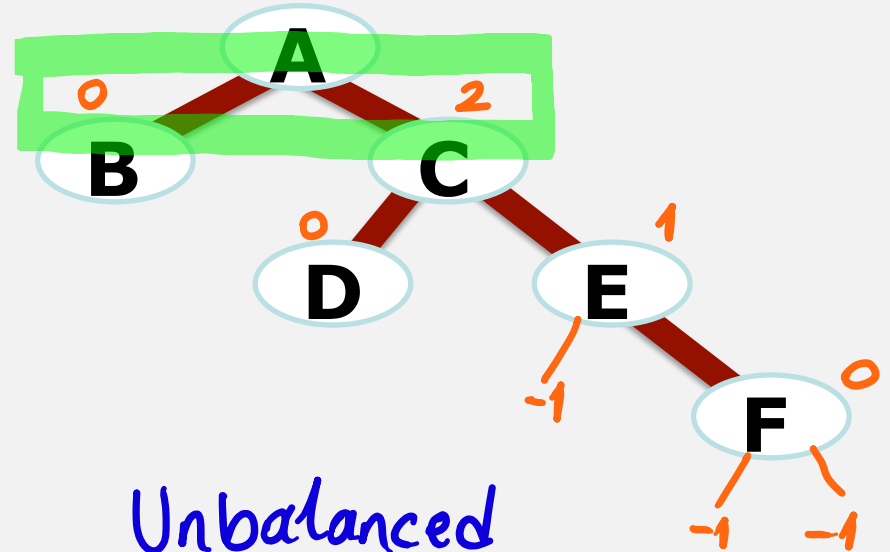
example: Let's add F

- Only nodes in the path to root are affected.
- If we inserted recursively, we can update the tree free.

Which one is unbalanced? Who decides that?
Height difference between siblings should not be more than 2



Balanced



Unbalanced

What kind of tree traversal algorithm do we need?

Pre-order
??

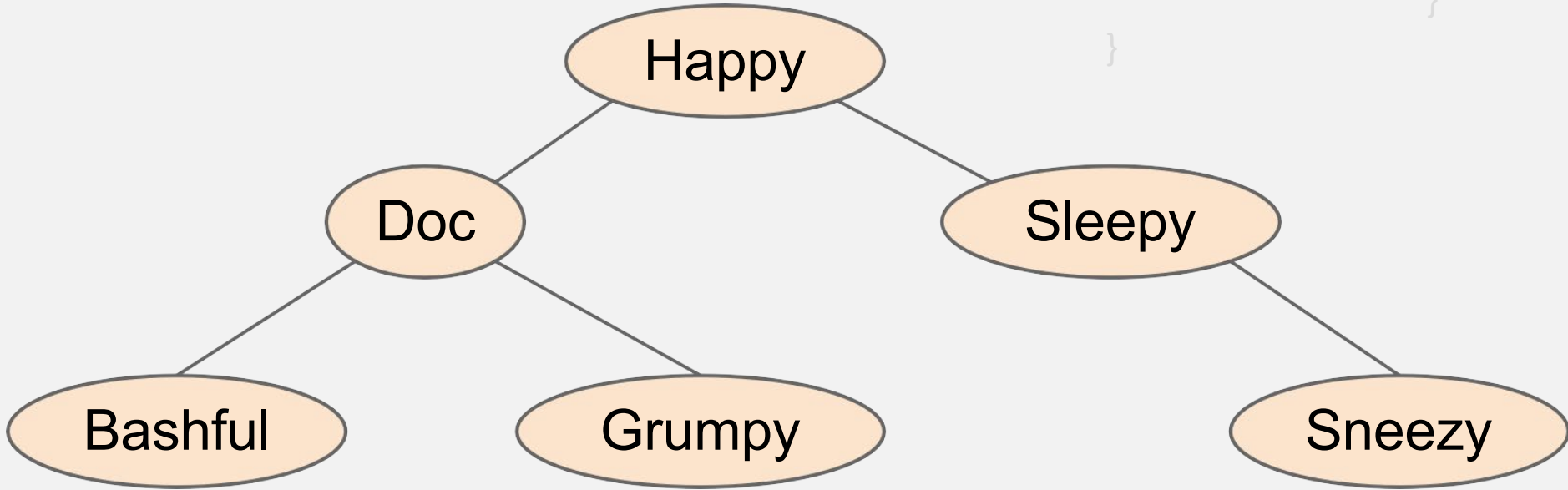
In-order
??

Post-order
??



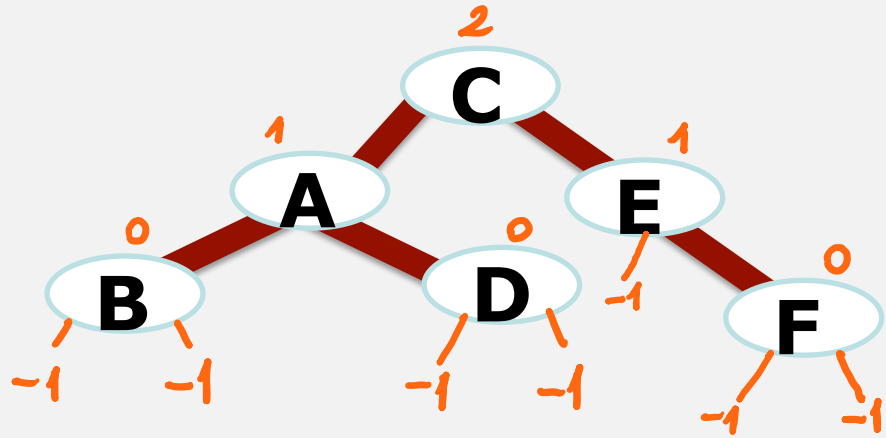
Tree: 6 of the 7 Dwarves

```
void traverse(struct node* root)
{
  if (root != NULL) {
    traverse(root->left);
    traverse(root->right);
    printf("%s", root->key);
  }
}
```

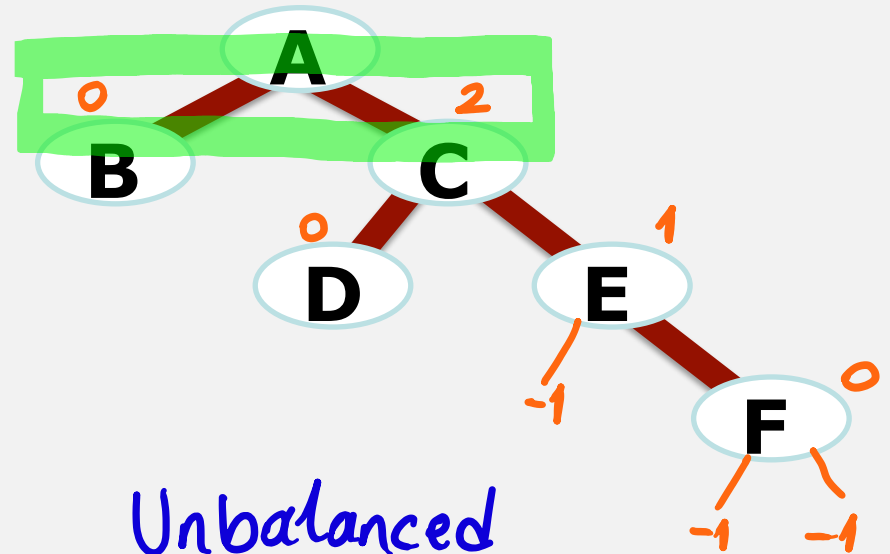


Post-order: Bashful, Grumpy, Doc, Sneezy, Sleepy, Happy

Which one is unbalanced? Who decides that?
Height difference between siblings should not be more than 2



Balanced



Unbalanced

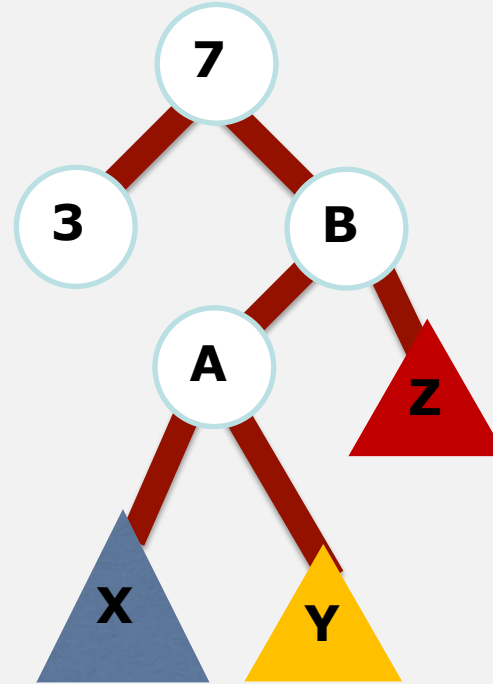
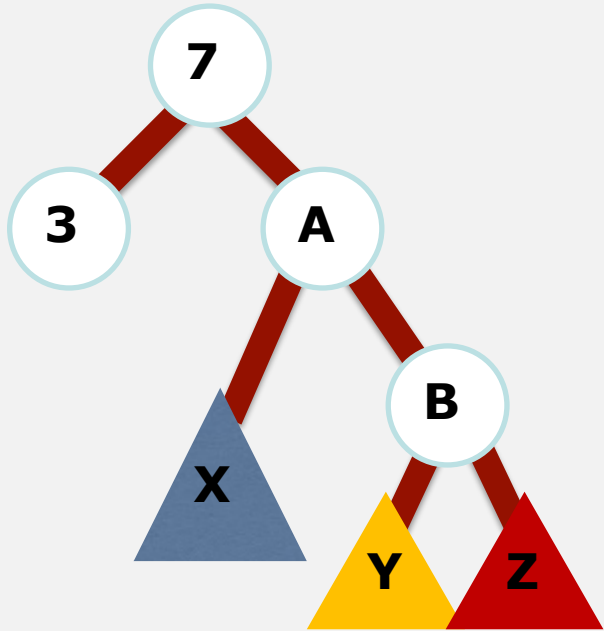
What kind of tree traversal algorithm do we need?

Post-order traversal is what we need here!



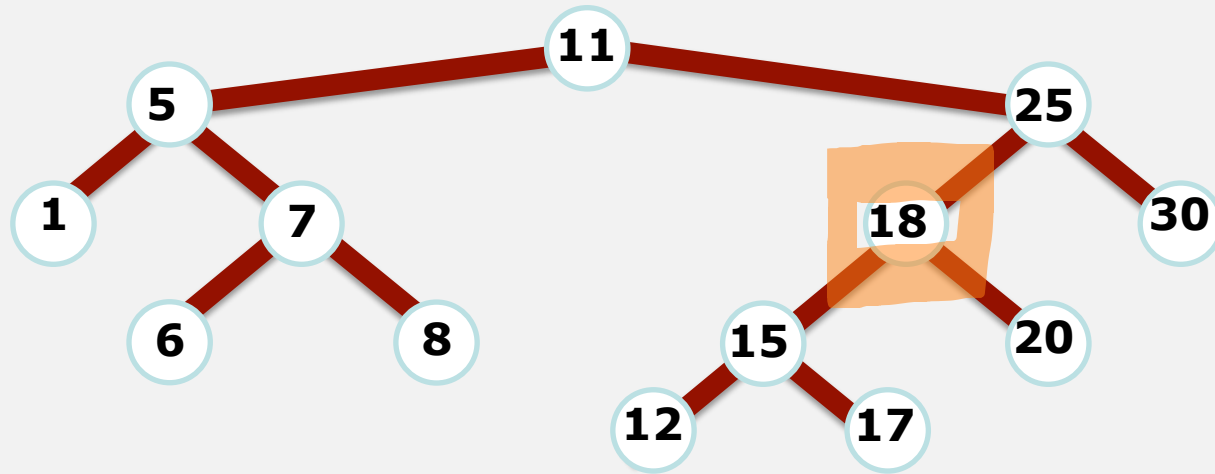
Rotations:

Let's rotate A to the left

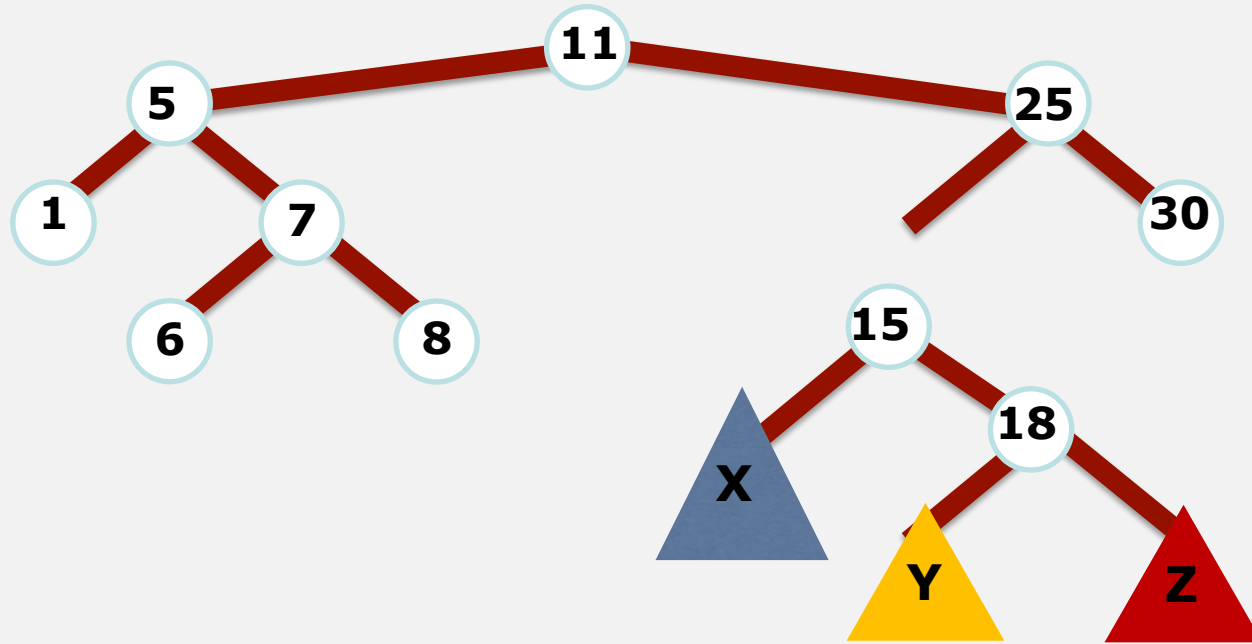
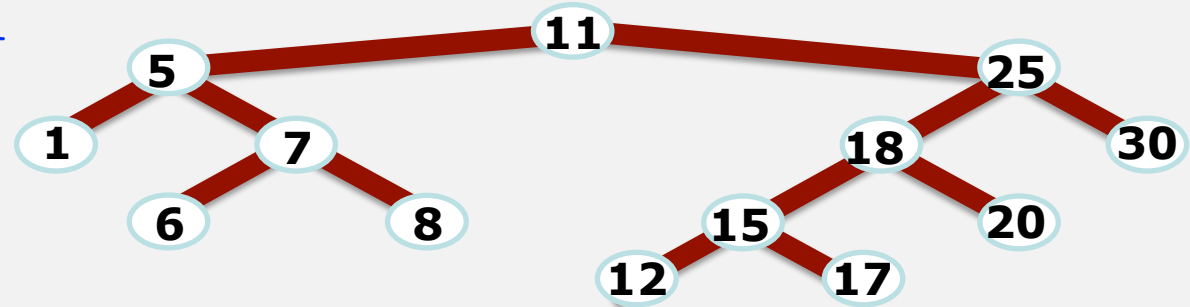


$X < a < Y < B < Z$

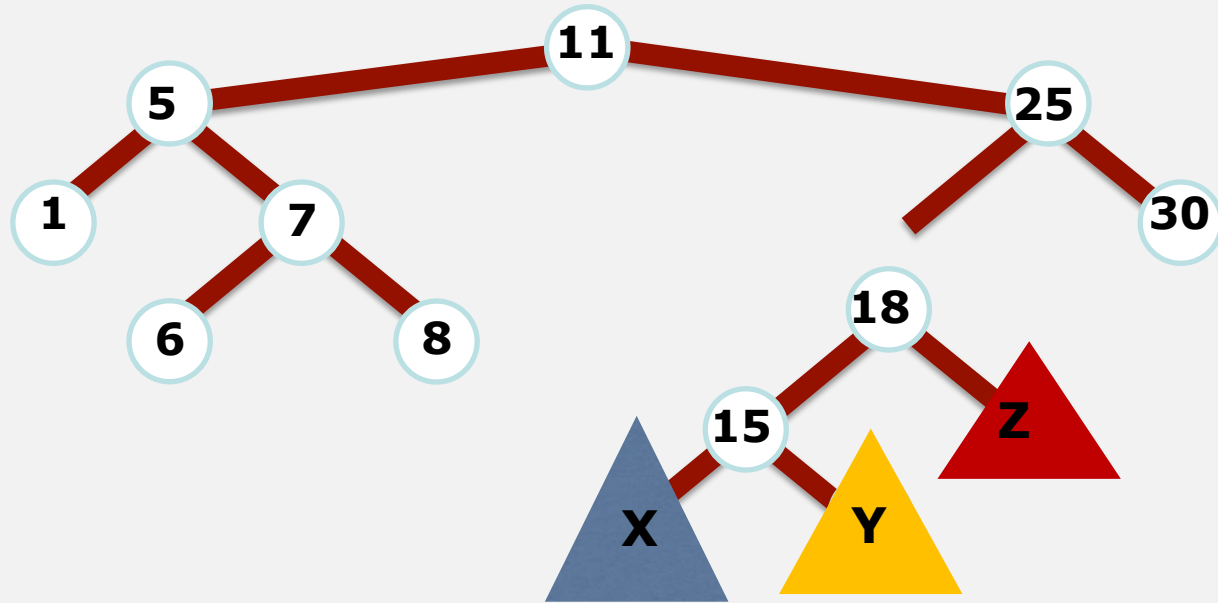
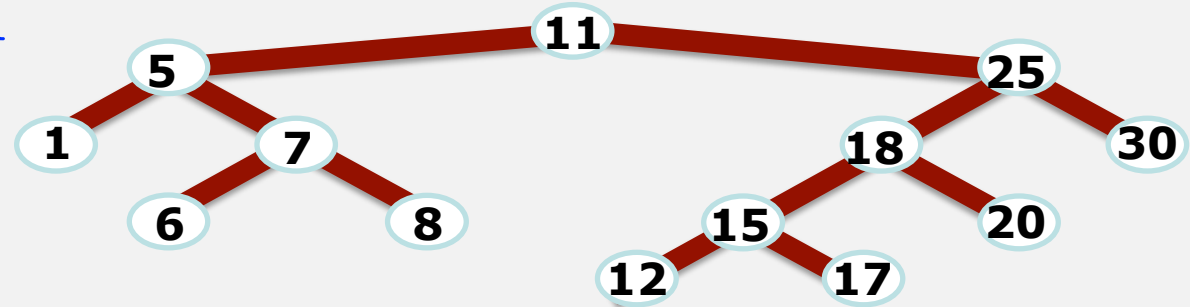
Exercise: Let's rotate 18 right



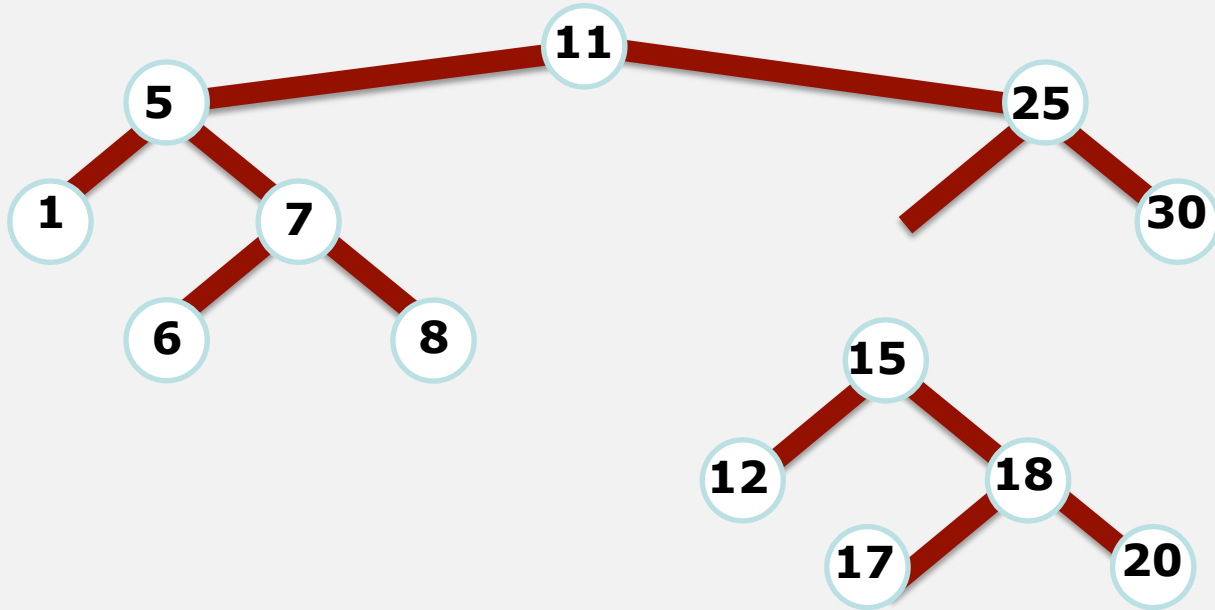
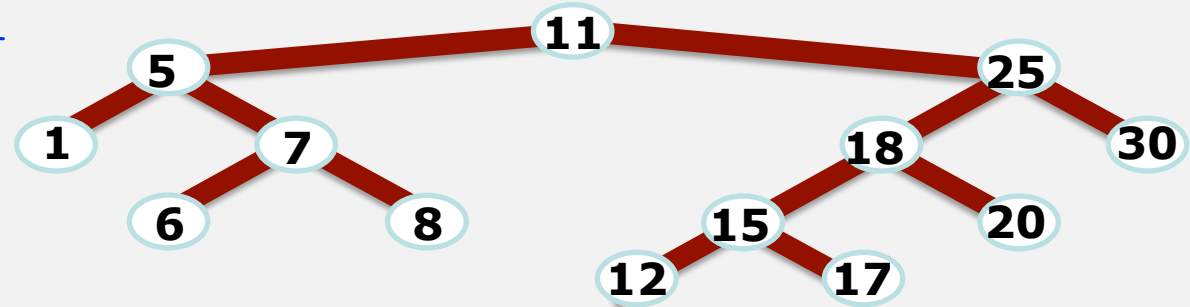
Exercise: Let's rotate 18 right



Exercise: Let's rotate 18 right



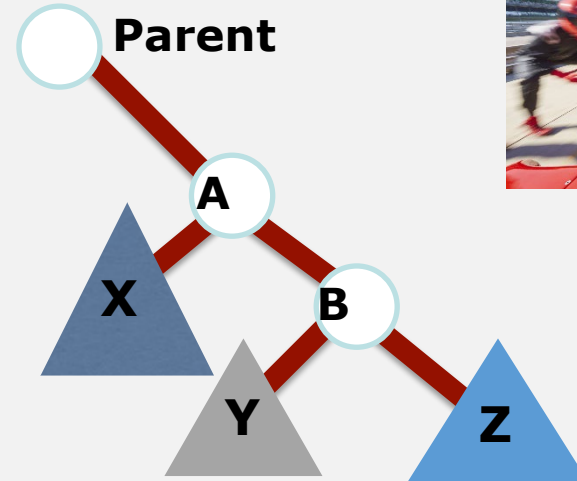
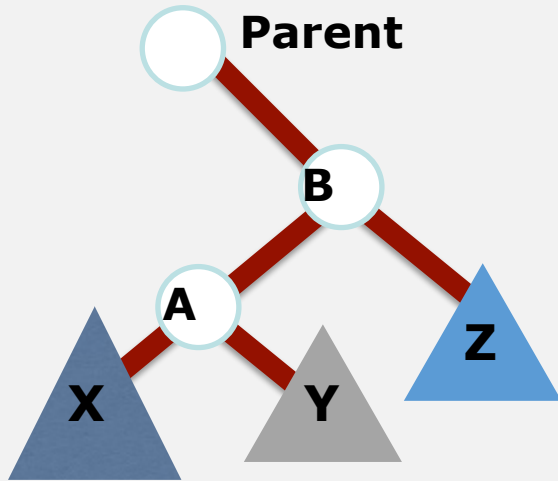
Exercise: Let's rotate 18 right



Implementation:

How fast are rotations?

$O(1)$ } Super fast



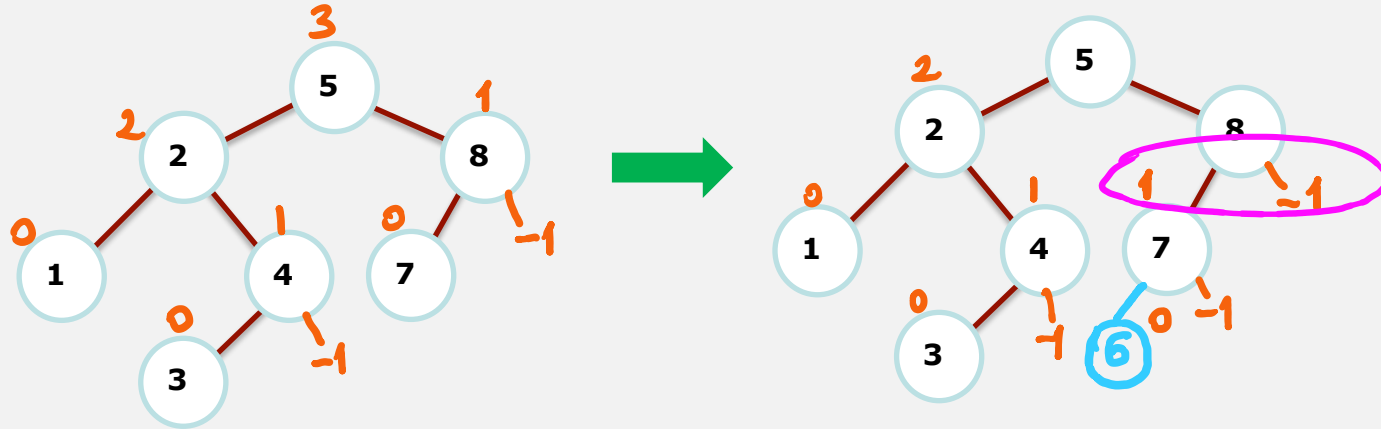
```
void rightRotateRightChild(avlNode *parent){
    avlNode *B;
    avlNode *A;

    B = parent->right;
    A = B->left;
    parent->right=A;           //Right child of parent is A
    B->left = A->right;        //Left child of B is right child of A
    A->right= B;              //A is parent of B now
}
```

Insert 6 into the tree :

↳ Let's check if the tree is a valid BST

↳ Then, let's check if the tree satisfies the AVL invariant



• We need to ask some questions:

Which node is responsible for this? 8

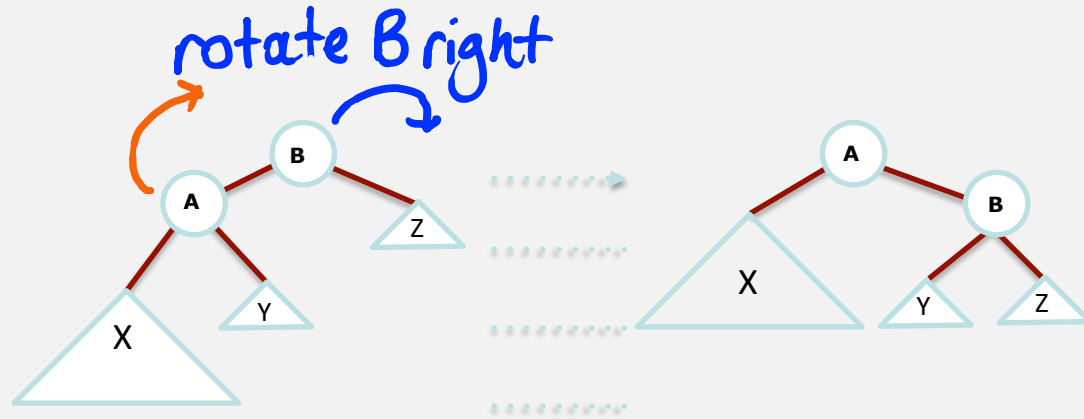
Explain what happened from the eyes of 8?

all possibilities : 8's left child's left subtree → single rotation

8's left child's right subtree → double rotation

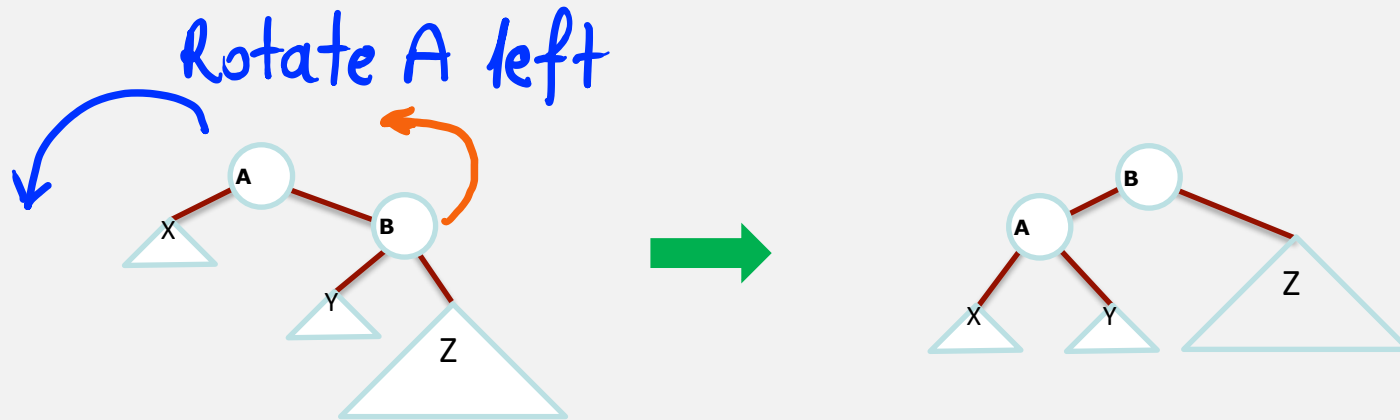
Insertion at left child's left subtree:

- Single rotation is sufficient



Insertion at Right Child's Right Subtree:

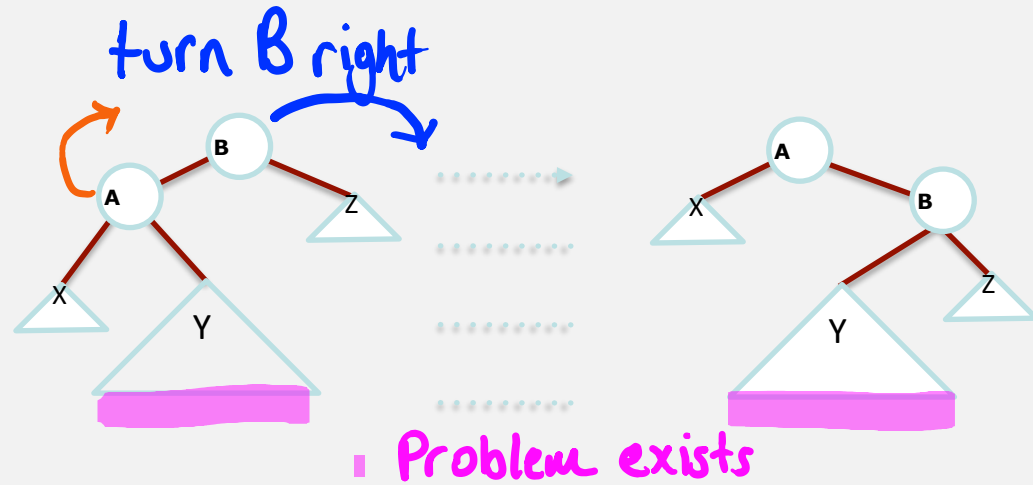
- Single rotation is sufficient



Insertion at Left Child's Right Subtree:

Single rotation will not be enough. Double rotation is needed.

Example: Assume some node is added to Y subtree and let Node B is responsible for this problem.



• Same problem would also occur if the insertion type was right child's left subtree because of the symmetry.

- The task is too difficult for B. Maybe B should seek help from A first. After A rotates to the lighter side, B can rotate right then.

• Exercise: Insert nodes into an AVL tree : 3, 2, 1, 4, 5, 6, 7

• Steps :

• Insert node into a BST \longrightarrow ③

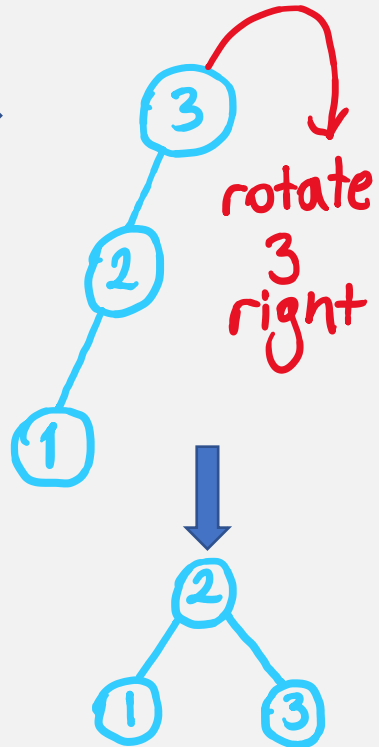
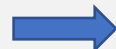
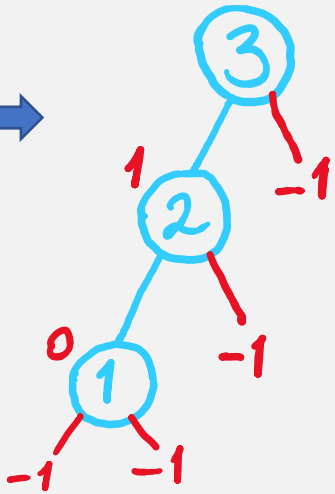
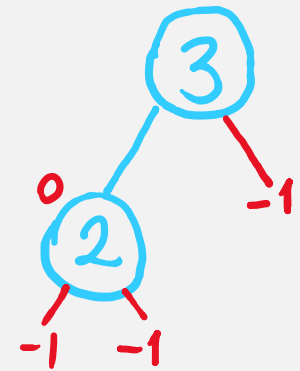
• check height difference of the children \longrightarrow



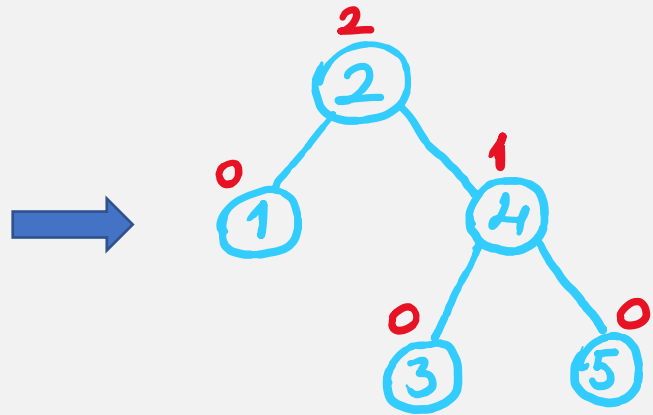
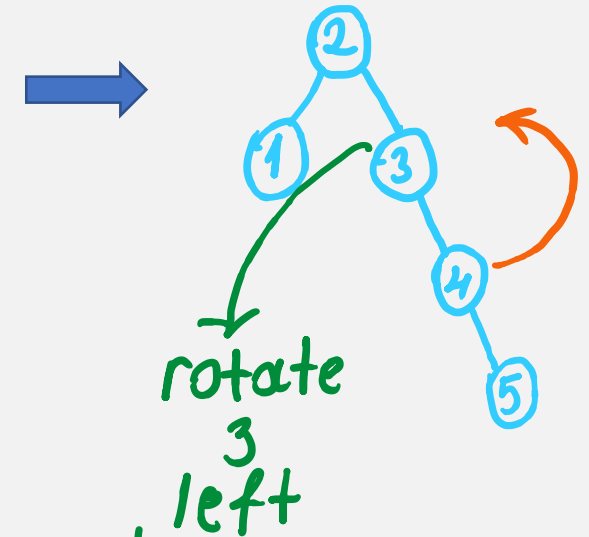
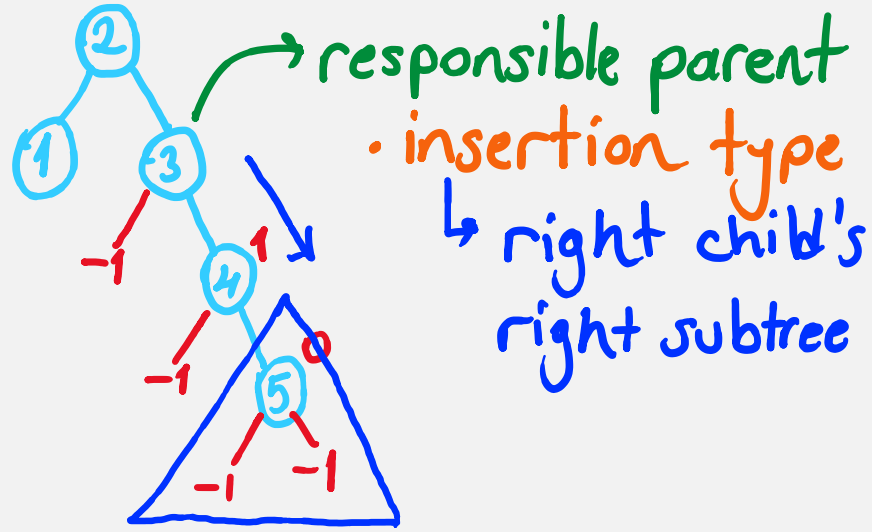
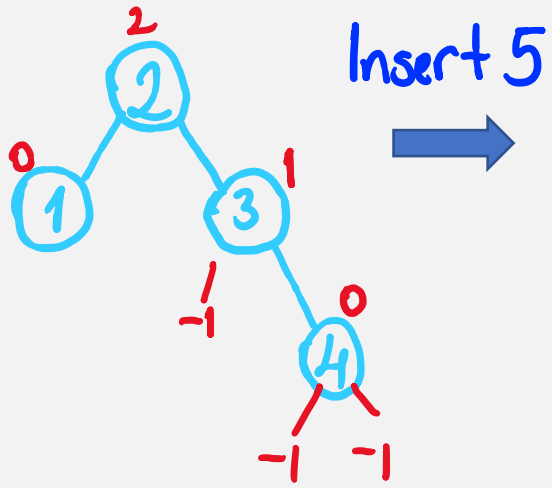
• Steps continued

• find the responsible parent
• analyze the type of insertion (left-right, left-left)??

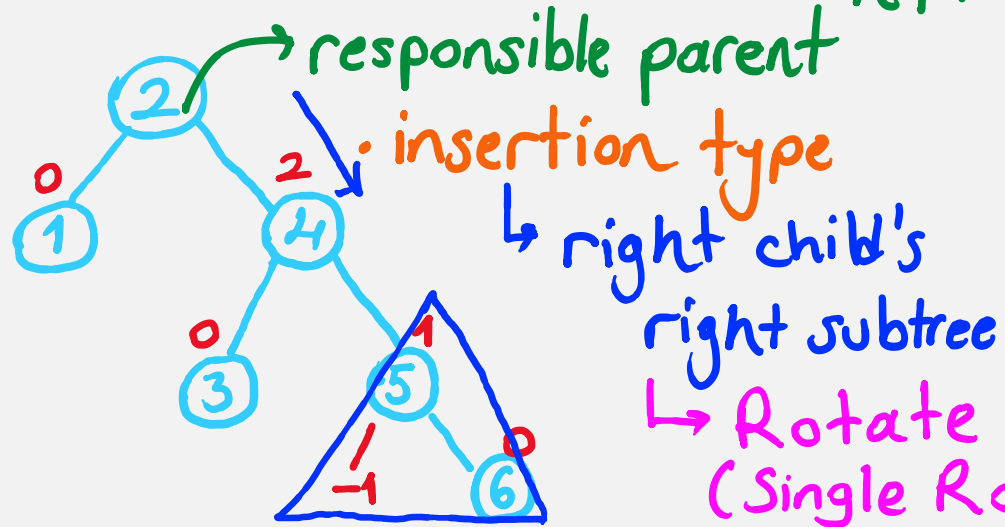
• if the parent needs a single rotation, rotate parent to the opposite side of insertion



• Exercise: Insert nodes into an AVL tree : 3, 2, 1, 4, 5, 6, 7

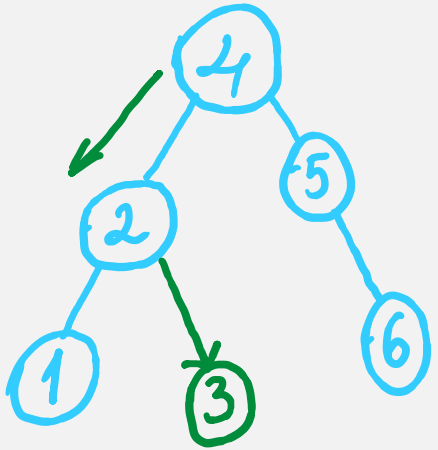
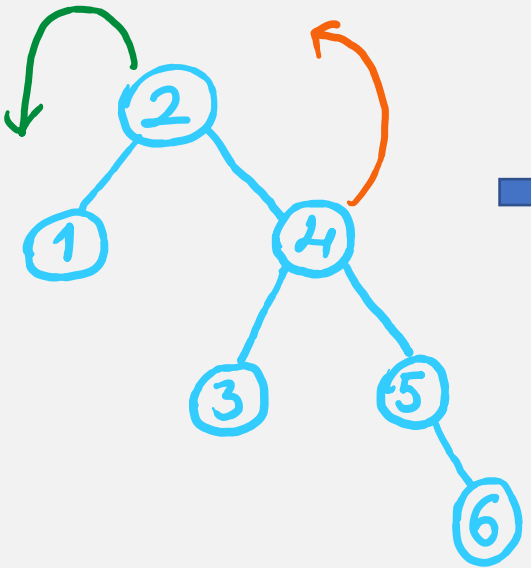


Insert 6



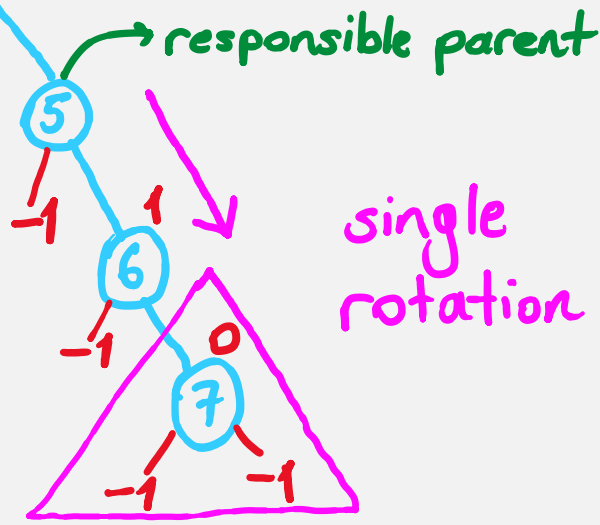
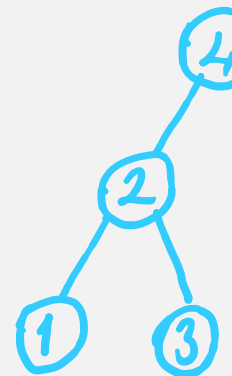
↳ Rotate 2 left (Single Rotation)

• Exercise: Insert nodes into an AVL tree : 3, 2, 1, 4, 5, 6, 7



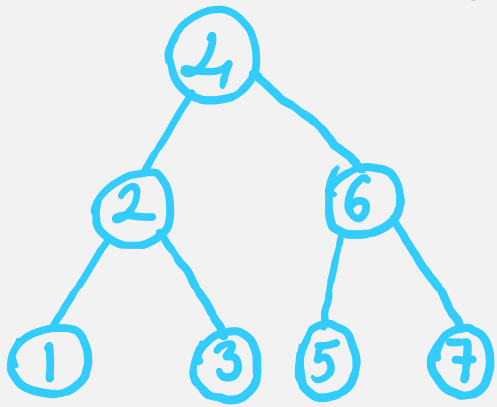
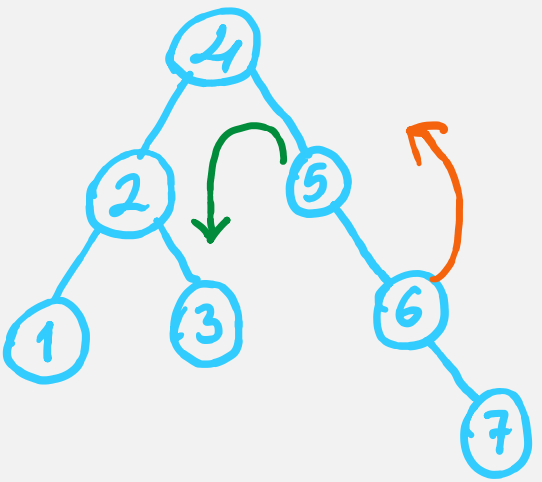
How about 3?

Insert 7



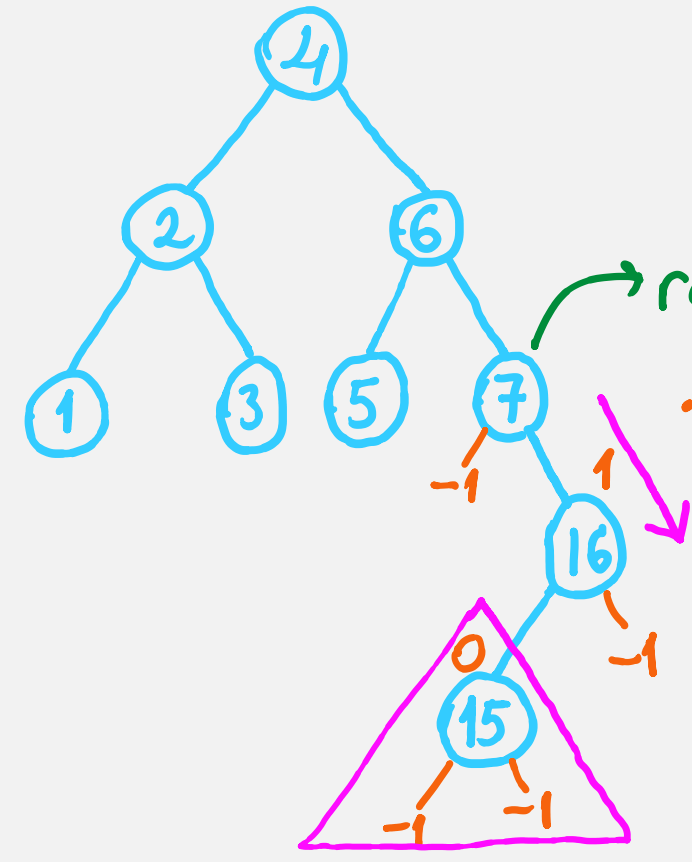
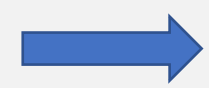
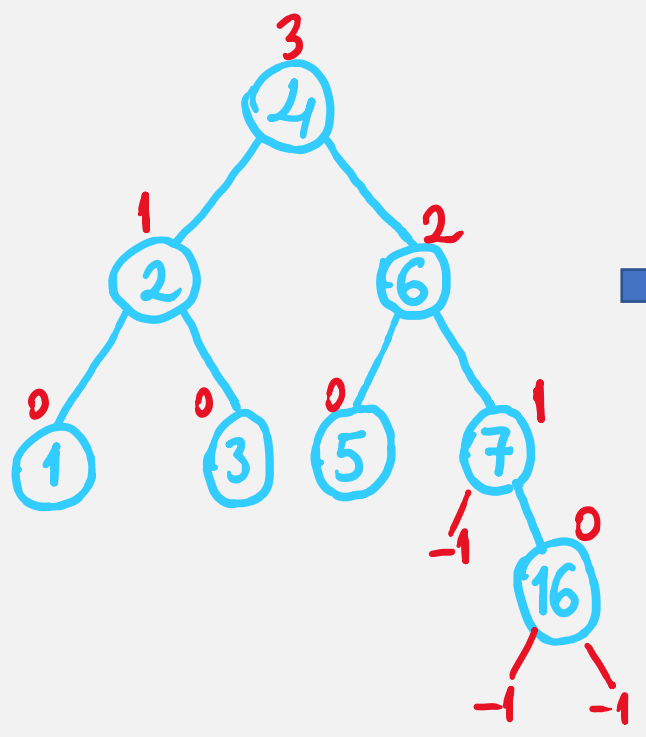
single rotation

Rotate 5 left



Ideal balance

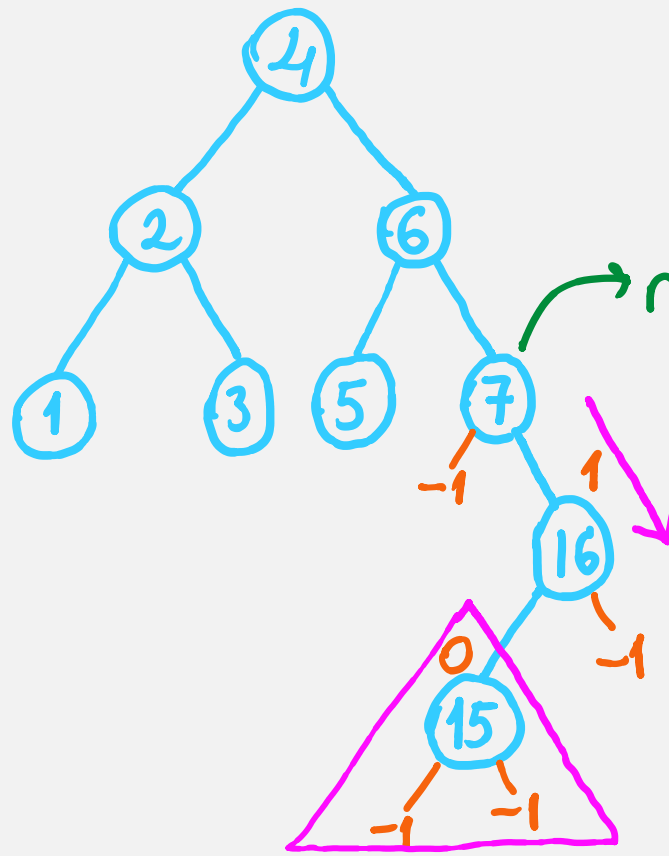
Exercise: Double Rotations: Add more nodes: 16, 15, 14



responsible parent

- insertion type
- needs double rotation

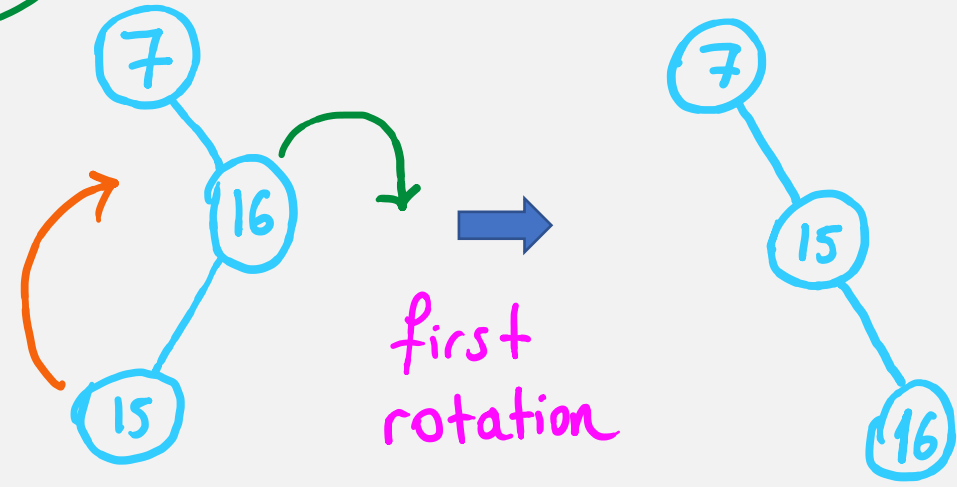
Exercise: Double Rotations: Add more nodes: 16, 15, 14



responsible parent

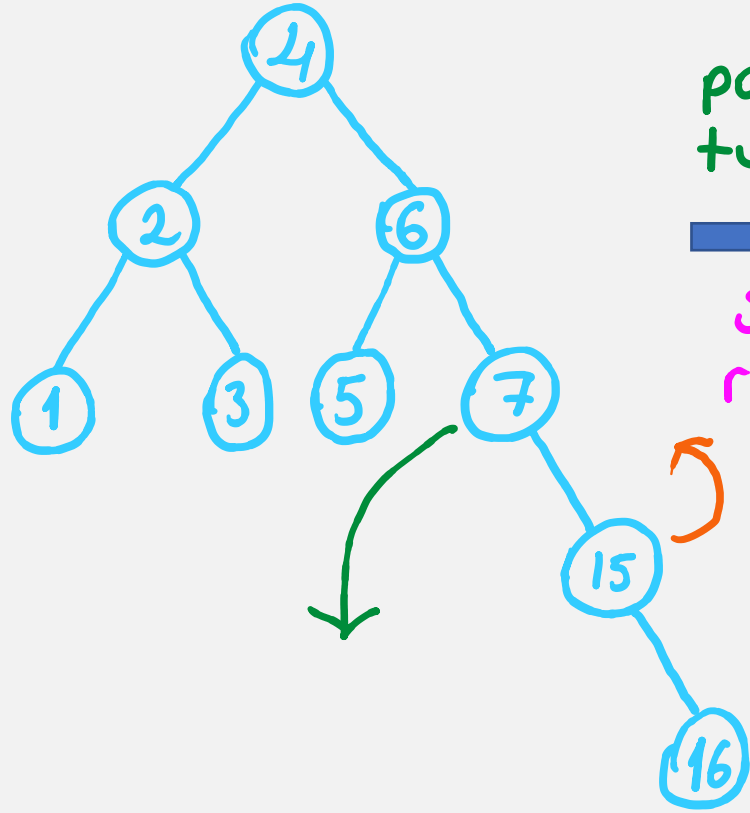
- insertion type
- needs double rotation

- 7 needs help from the child 16
- 16 should turn right because the insertion was to its left



first rotation

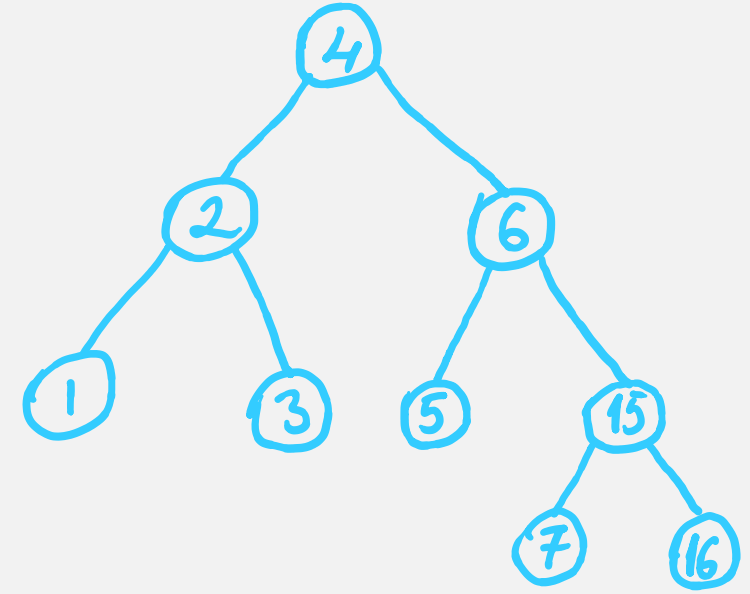
Exercise: Double Rotations: Add more nodes: 16, 15, 14



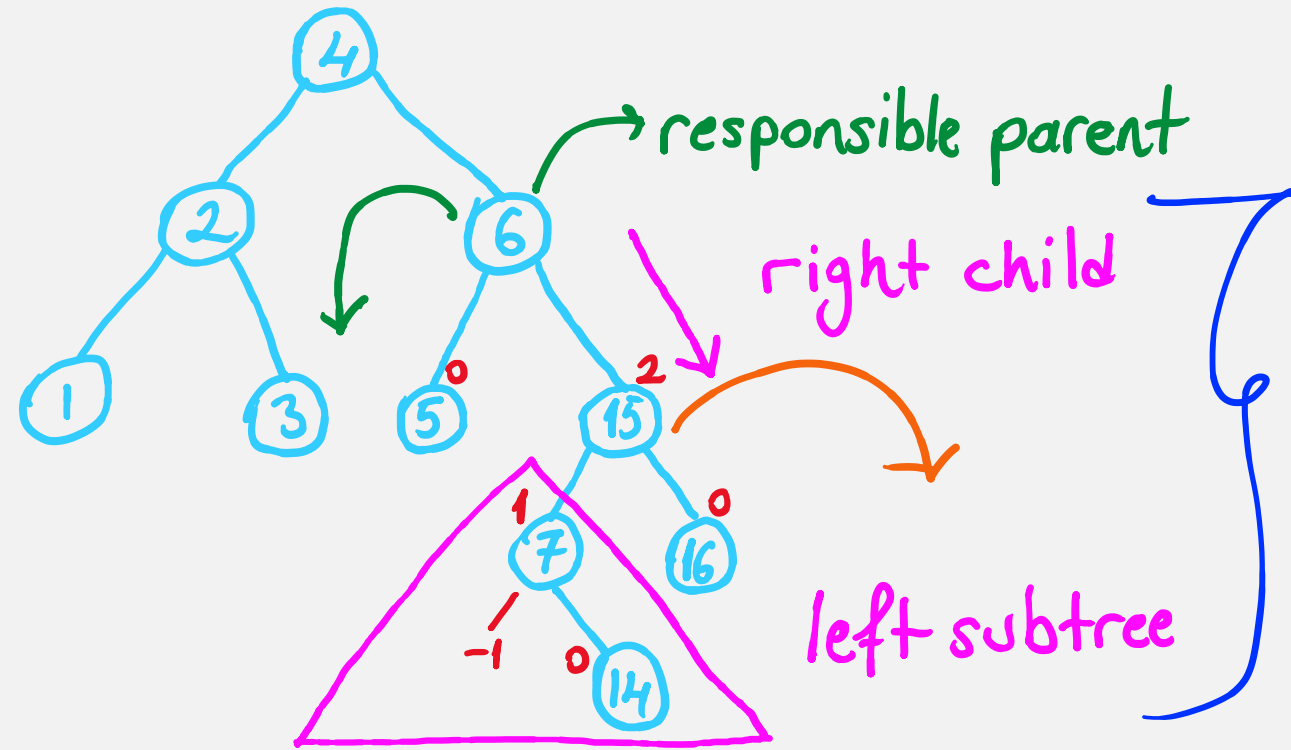
parent 7
turn left



second
rotation

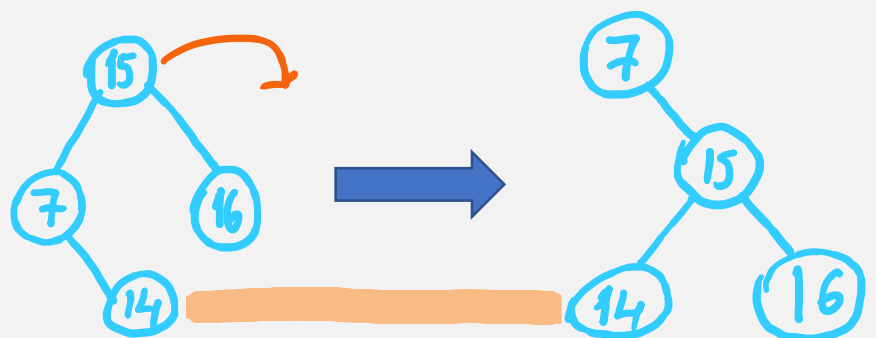


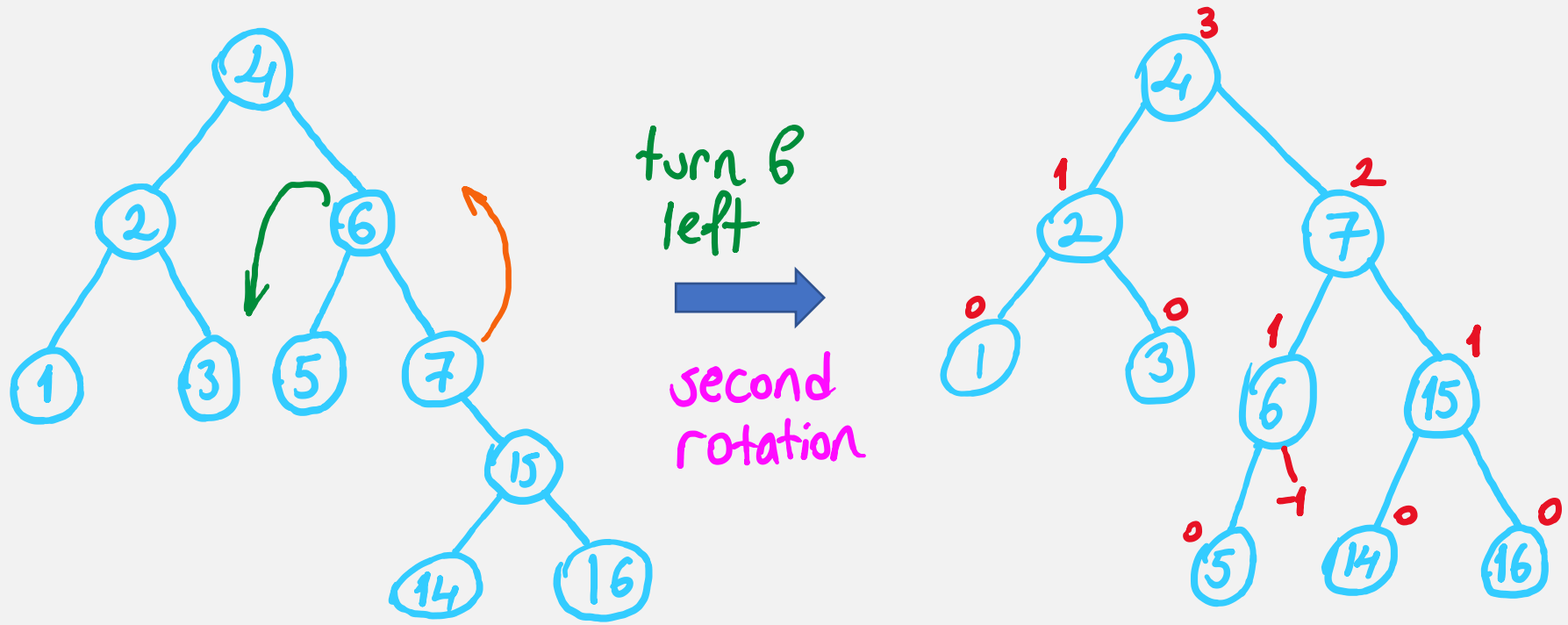
Exercise: Double Rotations: Add more nodes: 16, 15, 14



- Double rotation
1. 6 gets help from 15
 2. 15 turns right
 3. 6 turns left

first rotation





Practice Homework: Add 13 to the tree!