

## Mergesort

MERGE-SORT (A)

if  $\text{len}(A) < 2$

return a copy of A

else

$L \leftarrow$  1st half of A

$R \leftarrow$  rest of A

$L \leftarrow$  MERGESORT(L)

$R \leftarrow$  MERGESORT(R)

return MERGE(L, R)

} divide list into two parts

} conquer the subproblems

} combine results

— precondition: L, R are sorted

MERGE(L, R)

A  $\leftarrow$  empty list

$i \leftarrow 0$

$j \leftarrow 0$

while  $i < \text{len}(L)$  and  $j < \text{len}(R)$

if  $L[i] \leq R[j]$

append  $L[i]$  to A

$i \leftarrow i + 1$

else

append  $R[j]$  to A

$j \leftarrow j + 1$

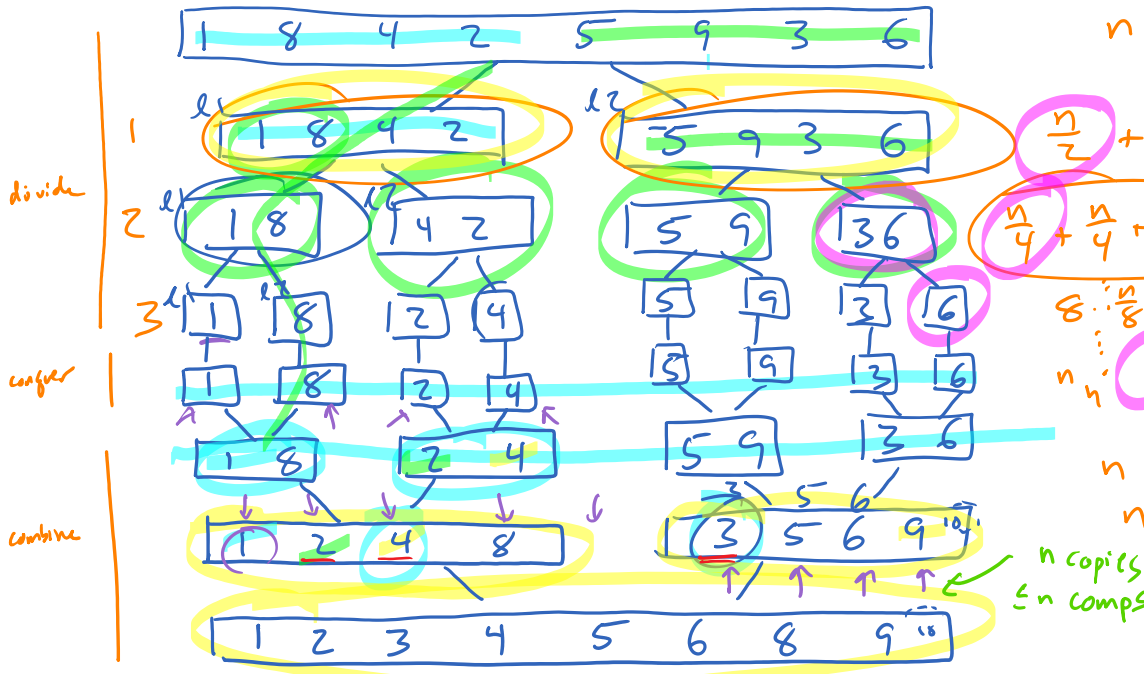
append  $L[i \dots \text{len}(L) - 1]$  to A

append  $R[j \dots \text{len}(R) - 1]$  to A

} concatenate remaining items on non-empty list to result

return A

Mergesort Example



$$n = 2^l \implies \log_2 n = l$$

$$\frac{n}{2} + \frac{n}{2}$$

$$\frac{n}{4} + \frac{n}{4} + \frac{n}{4} + \frac{n}{4}$$

$$8 \cdot \frac{n}{8}$$

$$n \cdot 1$$

$n$  copies @ each of  $\log_2 n$  levels

total copies =  $n \log_2 n$  in divide

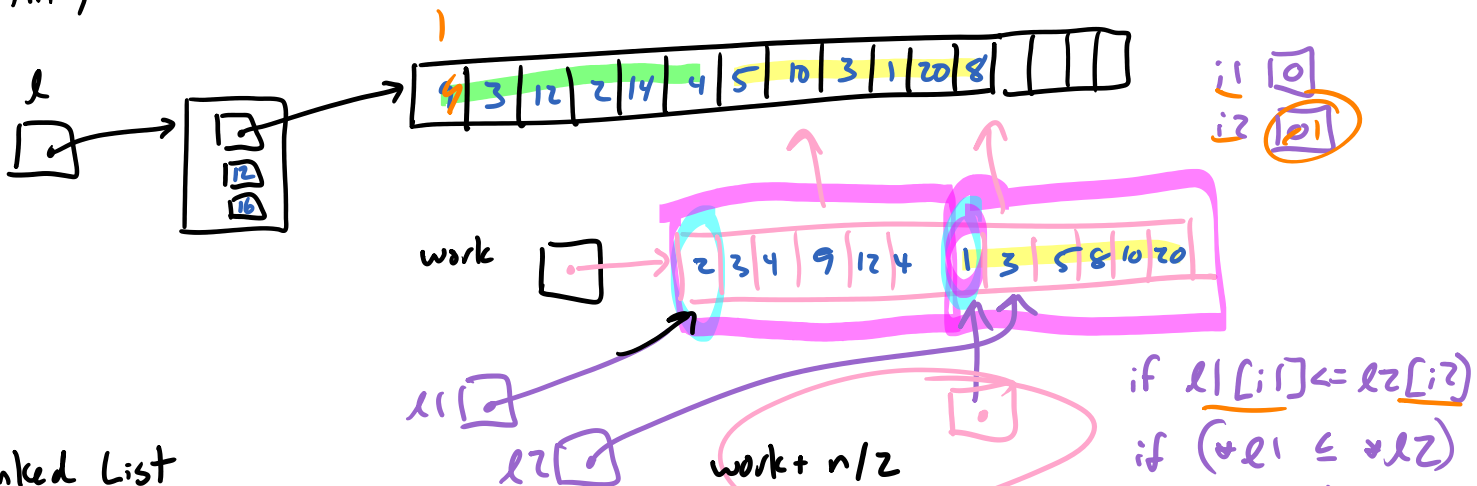
total copies =  $n \log_2 n$  in combine

total comps  $\leq n \log_2 n$

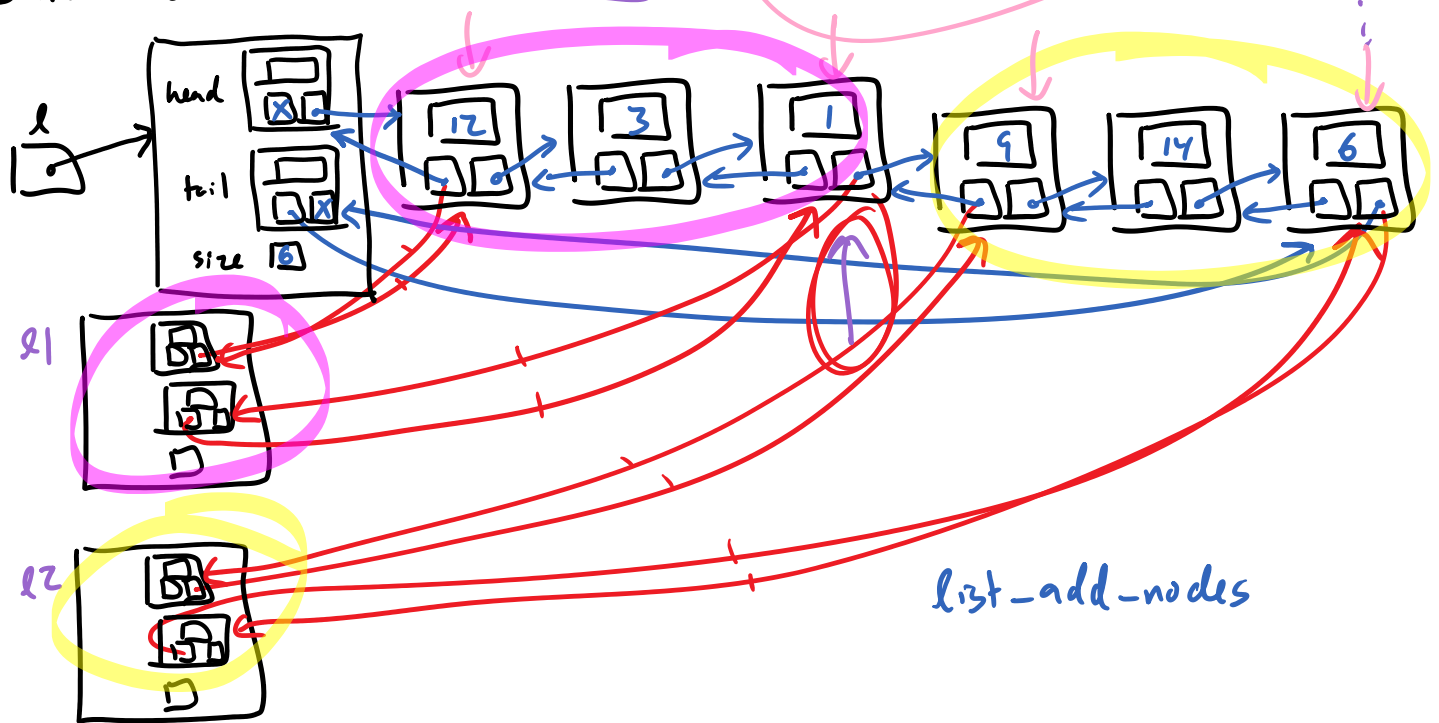
$$O(n \log n)$$

# Splitting

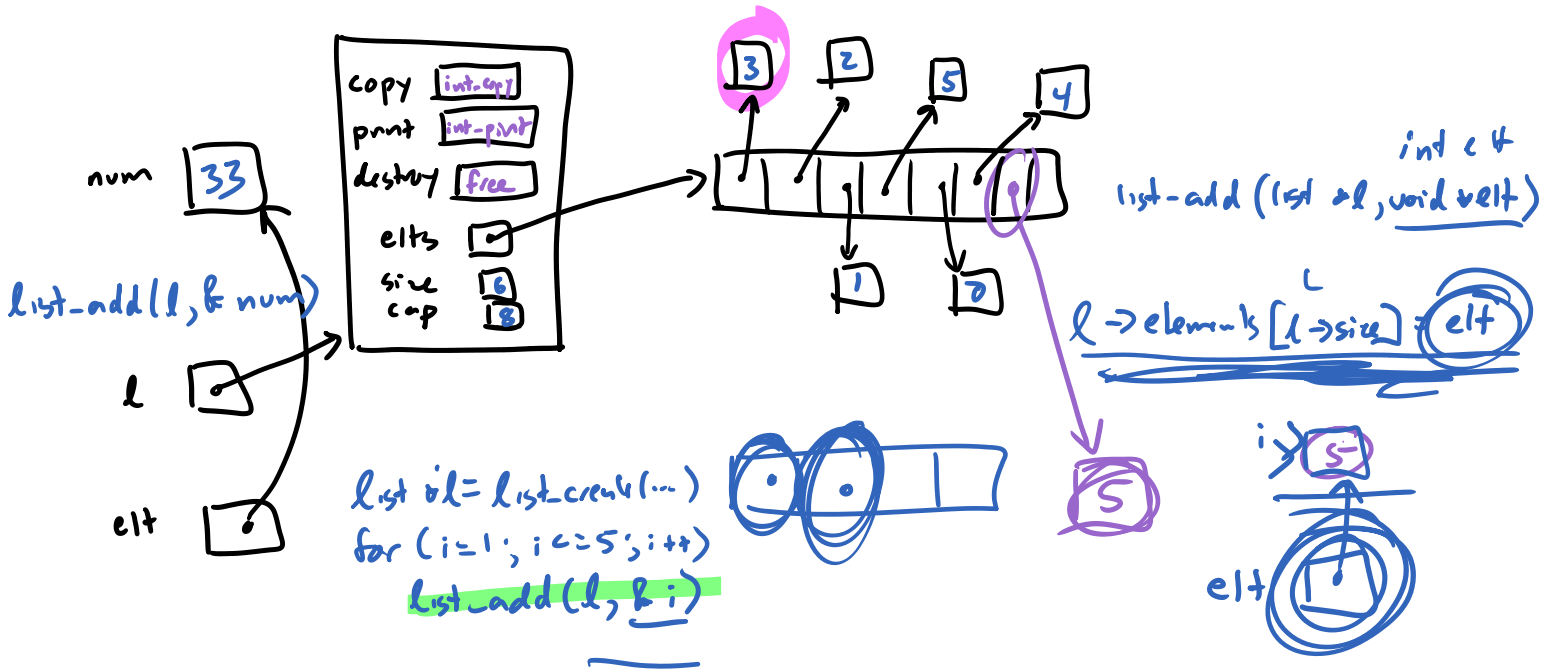
## Array-based list



## Linked List

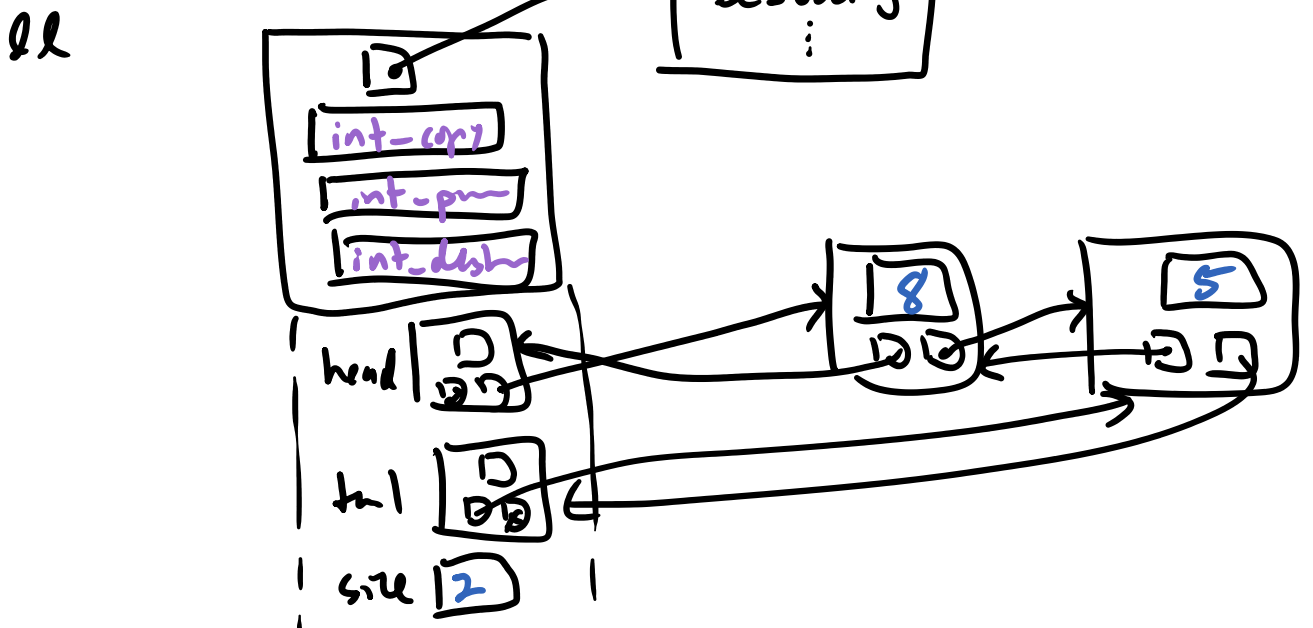
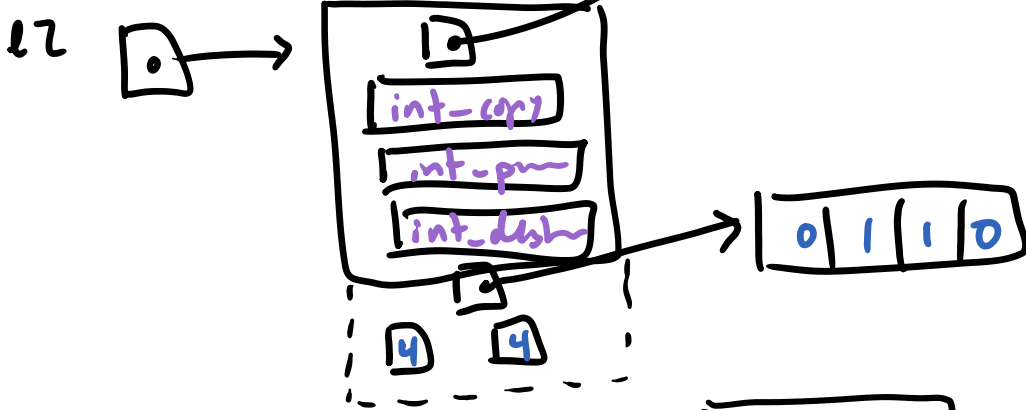
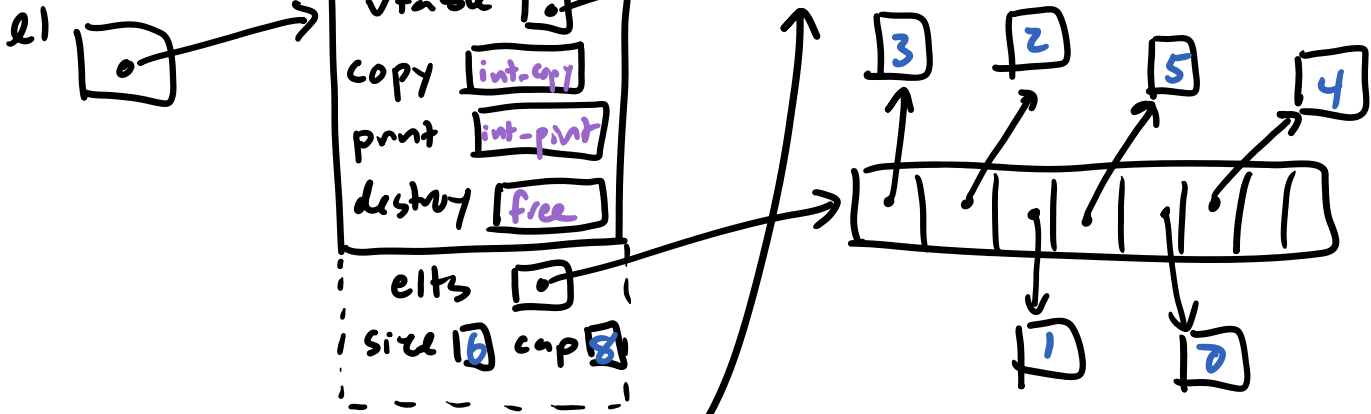


# Array-based list



# Polymorphic Lists

Array-based list



size 12