

# Counting Problems

k-Combination of  $\{1, \dots, n\}$  is a size-k subset of  $\{1, \dots, n\}$

"n choose k"  $\binom{n}{k} = \#$  of k-combinations of  $\{1, \dots, n\}$

$\binom{4}{2} = 6$   $\binom{3}{2}$   
 all the 2-combos of  $\{1, 2, 3\}$

1,2	2,3
1,3	2,4
1,4	3,4

all the 1-combos of  $\{1, 2, 3\}$   
 $\binom{3}{1}$  of those

$$= \begin{cases} 0 & \text{if } k > n \text{ or } k < 0 \\ 1 & \text{if } n = k \text{ or } k = 0 \\ \binom{n-1}{k-1} + \binom{n-1}{k} & \text{otherwise} \end{cases}$$

use n      don't use n

CHOOSE(n, k)

if  $k < 0$  or  $k > n$  return 0

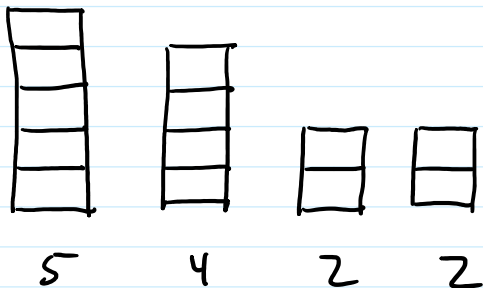
C ← n x n array

C[0,0] = 1

for i = 1 to n  
 C[i,0] ← 1  
 C[i,j] ← 1  
 for j = 1 to i-1

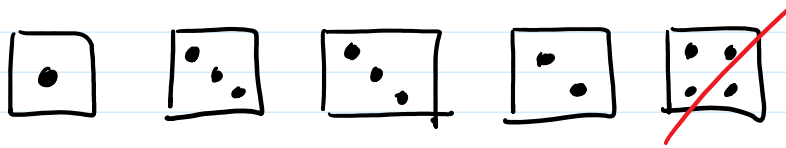
C[i,j] ← C[i-1,j-1] + C[i-1,j]

return C[n,k]



Count(n, h) = # nonincreasing sequences of length n, height h

= {



How many sequences of 5 die rolls sum to 13?

$\text{Count}(n, k) = \# \text{ sequences of } n \text{ rolls that sum to } k$

$O(nk)$  entries

$O(nk)$  total

$$= \begin{cases} 0 & \text{if } k < n \text{ or } k > 6n \\ 1 & \text{if } k = 0 \text{ and } n = 0 \\ \sum_{i=1}^6 \text{count}(n-1, k-i) & \text{otherwise} \end{cases}$$

$O(1)$  time

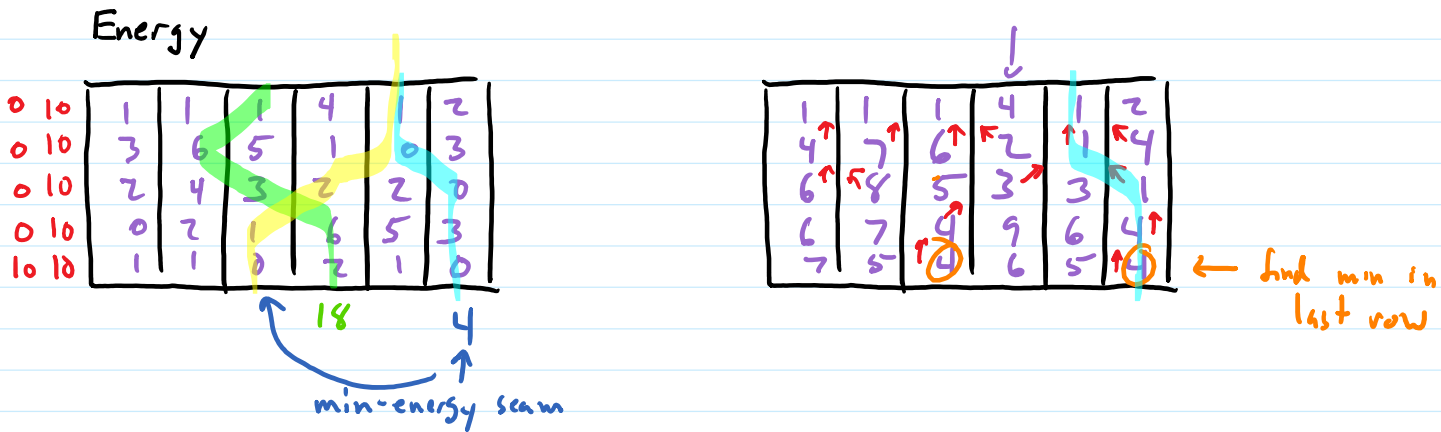
# Seam Carving

<https://www.youtube.com/watch?v=6NcIJXTlugg>

$M(i,j)$  = cost of min-cost seam from bp to row  $i$ , col  $j$

$$O(n \cdot m) \text{ entries} = \begin{cases} E(i,j) & \text{if } i=0 \\ E(i,j) + \min(M(i-1,j-1), M(i-1,j), M(i-1,j+1)) & \end{cases}$$

$O(n \cdot m)$  total  $\leftarrow$   $O(1)$



# Subset Sum

Given set of positive integers  $S$  and sum  $k$ , find  $S' \subseteq S$  to maximize  $\sum_{x \in S'} x$  s.t.  $\sum_{x \in S'} x \leq k$

Example:  $S = \{1, 4, 9, 13, 16, 20\}$   $k = 28$   
 $x_1$   $x_2$   $x_3$  ...

Brute Force: try each subset  $2^n$

Substructure:  $\{1, 4, 9\}$  was subset closest to  $28 - 13 = 15$  w/o going over among  $\{1, 4, 9\}$

$OPT(i, w)$  = highest possible sum  $\leq w$  using 1st  $i$  items

$O(n \cdot w)$  entries

$O(nw)$  total

$O(1)$  per entry

$$= \begin{cases} 0 & \text{if } i=0 \\ \max_{j=1 \dots i} x_j + OPT(j-1, w-x_j) & \text{if } x_i \leq w \\ \max(OPT(i-1, w-x_i) + x_i, OPT(i-1, w)) & \text{otherwise} \end{cases}$$

use element  $i$

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
1	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	
4	0	1	1	1	4	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	
9	0	1	1	1	4	5	5	5	5	9	10	10	10	13	14	14	14	14	14	14	14	14	14	14	14	14	14	14	
13	0	1	1	1	4	5	5	5	5	9	10	10	10	13	14	14	14	17	18	18	18	18	22	23	23	23	26	27	
16	0	1	1	1	4	5	5	5	5	9	10	10	10	13	14	14	14	17	18	18	18	20	21	21	21	25	26	27	
20	0	1	1	1	4	5	5	5	5	9	10	10	10	13	14	14	14	17	18	18	18	20	21	21	21	24	25	26	27

$$S = \{1, 4, 9, 13, 16, 20\} \quad k = 23$$

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23

1  
4  
9  
13  
16  
20

# Vendor Scheduling

	Week					
	1	2	3	4	5	
Brooklyn	6	3	10	4	1	m = 16
Stratford	4	6	3	8	20	

<https://play.golang.org/p/NgB2MxoSmX>

opt S S S S S 41

OPT for 1<sup>st</sup> 4 weeks?

NO: BBBB = 23 > 21

$OPT_B(n)$  = best net through week n ending in B

$OPT_S(n)$  = " " " " ending in S

$$OPT_B(n) = \begin{cases} B_1 & \text{if } n=1 \\ \max(OPT_S(n-1) - m, OPT_B(n-1)) + B_n \end{cases}$$