

Vendor Scheduling

	Week					
Brooklyn	6	3	10	4	1	<div style="border: 1px solid red; padding: 2px; display: inline-block;">M=3</div> moving cost M=10
Stratford	4	6	3	8	20	
	S	B	B > B	B B	B	
	B	B	B	S	S	44

<https://plav.golang.org/p/YvmAsxoYER>

not opt for 2 week
 but B opt for 2 week
 given end in B

B_i = profit in Brooklyn in week i
 S_i = profit in Stratford in week i

$OPT_B(i)$ = max profit for weeks 1...i ending at B

$OPT_S(i)$ = " " " " " S

$$OPT_B(i) = \begin{cases} B_1 & \text{if } i=1 \\ \max_{\text{stay in B}} (OPT_B(i-1) + B_i, OPT_S(i-1) + B_i - M) \end{cases}$$

$$OPT_S(i) = \begin{cases} S_1 & \text{if } i=1 \\ \max (OPT_S(i-1) + S_i, OPT_B(i-1) + S_i - M) \end{cases}$$

$\Theta(1)$ to compute each entry

$\Theta(n)$ entries ($n = \# \text{ weeks}$)

$\Theta(n)$ total

OPT_B	B_1					
OPT_S	S_1					

Longest Common Subsequence

A C C G T A A C T
 G T C T C T A G A

X: A C C G T A A C T
 Y: G T C T C T A G A

CCTAA is LCS of $\begin{matrix} X, Y \\ X_{1..8}, Y \\ X_{1..7}, Y \end{matrix}$

CCTA is LCS of $X_{1..6}, Y_{1..6}$

length of LCS of $X_{1..i}, Y_{1..j}$ $\geq \begin{cases} 0 & \text{if } i=0 \text{ or } j=0 \\ 1 + \text{lcs-len}(i-1, j-1) & \text{if } X_i = Y_j \\ \max(\text{lcs-len}(i, j-1), \text{lcs-len}(i-1, j)) & \text{otherwise} \end{cases}$

$\theta(n \cdot m)$ entries
 $\theta(1)$ to calc each
 $\theta(n \cdot m)$ total



THM: Let Z be LCS of X and Y

Then if $x_m = y_n$ then $Z_{1..k-1}$ is an LCS of $X_{1..m-1}, Y_{1..n-1}$

and if $x_m \neq y_n$ and $Z_k \neq x_m$ then Z is an LCS of X_{m-1}, Y

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Proof: 1) Suppose Z is LCS of X, Y and $x_m = y_n$

Suppose further that $Z_{1..k-1}$ is not LCS of $X_{1..m-1}, Y_{1..n-1}$
 there is Z' s.t. Z' is LCS of $X_{1..m-1}, Y_{1..n-1}$
 and $\text{len}(Z') > \text{len}(Z) - 1$

Then $Z' \cdot x_m$ is of X, Y

$\text{len}(Z' \cdot x_m) = \text{len}(Z') + 1 > \text{len}(Z)$
 so Z not LCS

So $Z_{1..k-1}$ is LCS of ...

2) Suppose Z is LCS of X, Y , $z_k \neq x_m$ and $x_m \neq y_n$
 Z is CS of $X_{1\dots m-1}, Y$:

Z is LCS of $X_{1\dots m-1}, Y$:

3) similar

LCS Example

		A	C	C	G	T	A	A	C	T
	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	1	1	1	1	1	1
T	0	0	0	0	1	2	2	2	2	2
C	0	0	1	1	1	2	2	2	3	3
T	0	0	1	1	1	2	2	2	3	4
C	0	0	1	2	2	2	2	2	3	4
T	0	0	1	2	2	3	3	3	3	4
A	0	1	1	2	2	3	4	4	4	4
G	0	1	1	2	3	3	4	4	4	4
A	0	1	1	2	3	3	4	5	5	5

The table shows the dynamic programming table for finding the Longest Common Subsequence (LCS) between the strings "GTCAT" (rows) and "ACCTAT" (columns). The values in the cells represent the length of the LCS for the substrings up to that point. Red arrows indicate the path taken to reach the value 5 at the bottom-right cell (row 10, column 10). The path starts at (10, 10) and moves left to (10, 9), then up to (9, 9), then left to (9, 8), then up to (8, 8), then left to (8, 7), then up to (7, 7), then left to (7, 6), then up to (6, 6), then left to (6, 5), then up to (5, 5), then left to (5, 4), then up to (4, 4), then left to (4, 3), then up to (3, 3), then left to (3, 2), then up to (2, 2), then left to (2, 1), and finally up to (1, 1).