

0-1 Knapsack

Inner tube: 106g	26 troy dr \$8
trading cards: 38g	10 dr \$25
drink coupons: 2g	1 dr \$1
book: 246g	50dr \$5
tp: 164g	41dr \$50

Jordan and Jamie

carbon fiber reinforced titanium magnesium rods

Avery and Ashley

0-1 Knapsack

Problem: Given 1) n items with weights w_1, \dots, w_n
values v_1, \dots, v_n } ← positive integers
2) weight bound W } ←

Find subset of n items of highest total value
and total weight $\leq W$

Brute force $O(2^n)$
⚡

	weight (tray dr)	value (B)
1) inner tube	26	8
2) trading cards	10	25
3) book	50	5
4) drink coupons	1	1
5) toilet paper	41	50

$$W = 50$$

0-1 Knapsack Optimal Substructure

item	weight	value
✓ 1	1	2
x 2	2	3
✓ 3	4	4
x 4	6	5
x 5	12	18
✓ 6	15	22

total = 24

W = 20

the highest valued subset of 1st 5 items with total weight ≤ 5

→ Suppose that {1,3,6} is optimal for $n=6, W=20$ but {1,3} is not optimal for $n=5, W=5$

Then we could do better than 6 for $n=5, W=5$ and better than $6+22$ for $n=6, W=20$

So {1,3,6} would not be optimal

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✓ 1	1	2
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x 5	12	18
✓ 6	15	22

total = 24

W = 20

the highest valued subset of 1st 5 items with total weight ≤ 5

$$OPT[6, 20] = v_6 + OPT[5, 20 - w_6]$$

$OPT[i, j] = \max$ possible value of subset of 1st i items with total weight $\leq j$

$$= \begin{cases} 0 & \text{if } i=0 \text{ or } j=0 \\ OPT[i-1, j] & \text{if } w_i > j \\ \max(v_i + OPT[i-1, j-w_i], OPT[i-1, j]) & \text{otherwise} \end{cases}$$

includes (pointing to $v_i + OPT[i-1, j-w_i]$)
don't include (pointing to $OPT[i-1, j]$)

```

for j = 0 to W
  OPT[0, j] = 0
for i = 0 to n
  OPT[i, 0] = 0
for i = 1 to n
  for j = 1 to W
    OPT[i, j] =
  
```

$O(nW)$ entries
 $O(1)$ each

$O(nW)$

OPT

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	2	3	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
3	0	2	3	5	5	6	7	9	9	9	9	9	9	9	9	9	9	9	9	9	9
4	0	2	3	5	5	6	7	9	9	10	10	11	12	14	14	14	14	14	14	14	14
5	0	2	3	5	5	6	7	9	9	10	10	11	18	20	21	23	23	24	25	27	27
6	0	2	3	5	5	6	7	9	9	10	10	11	18	20	21	23	24	25	27	27	29

pseudopolynomial : $O(n \cdot 2^{\text{#bits in } w})$

1 3 6

```

for i = 1 to n
  for j = 1 to W
    if OPT[i, j] = OPT[i-1, j]
      CHOOSE[i, j] = N
    else
      CHOOSE[i, j] = Y
  
```

CHOOSE

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	
1	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y
2	N	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y
3	N	N	N	N	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y
4	N	N	N	N	N	N	N	N	N	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y
5	N	N	N	N	N	N	N	N	N	N	N	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y
6	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	Y	Y	Y	Y	Y	Y