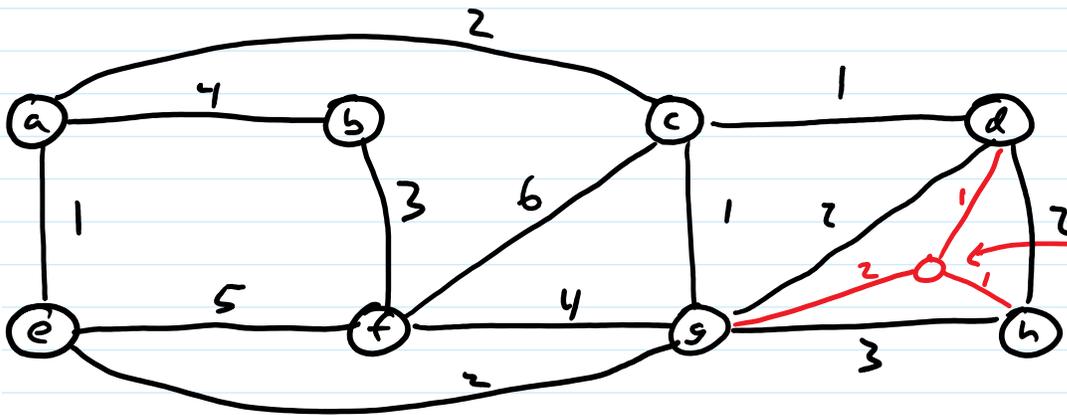
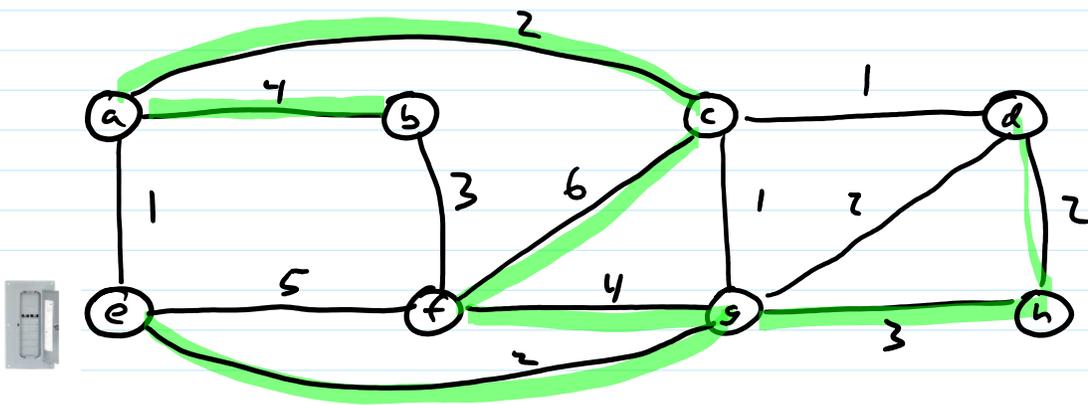


Minimum Spanning Tree



another hole in the ceiling is

Minimum Spanning Tree

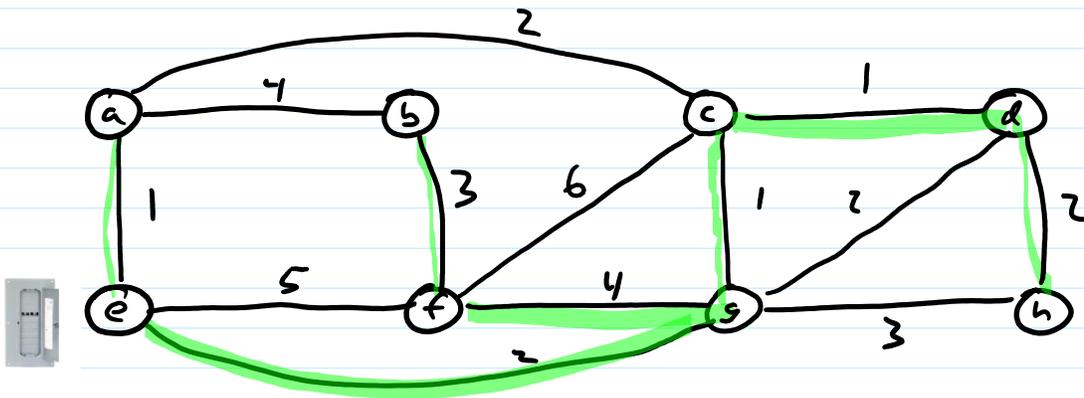


$$4 + 2 + 6 + 4 + 2 + 3 + 2 = 23$$

spanning tree: a subset of edges that forms a tree and connects all vertices

minimum spanning tree: a spanning tree of minimum total weight

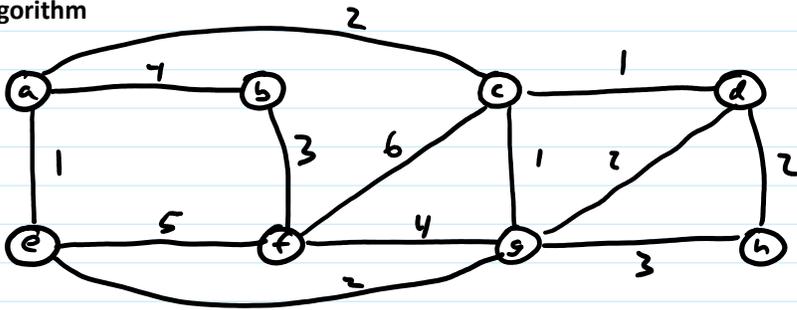
Minimum Spanning Tree



spanning tree: a subset of edges that forms a tree and connects all vertices

minimum spanning tree: a spanning tree of minimum total weight

Kruskal's Algorithm

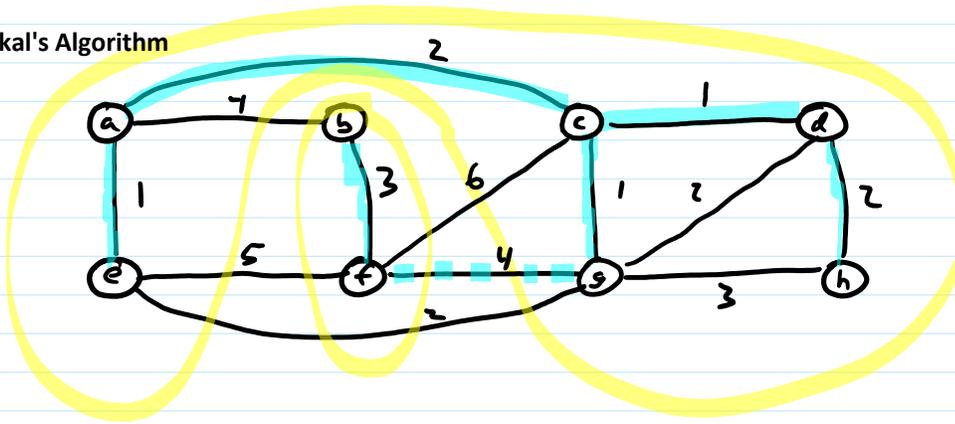


sort edges in order of increasing weight (a,e) (c,g) (c,d) (d,g) (a,c) (d,h) (e,g) (g,h) (b,f) (f,g) (a,b) (c,f)

for each edge (u,v) in order of sort

if u and v aren't connected using previously selected edges
add (u,v)

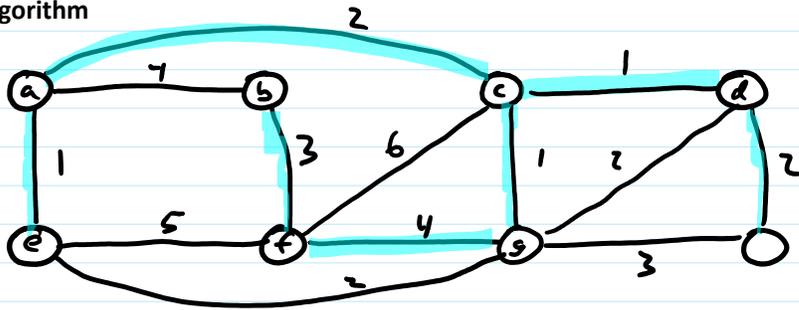
Kruskal's Algorithm



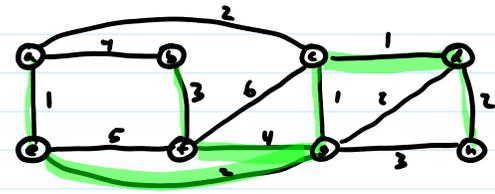
sort edges in order of increasing weight (a,e) (c,g) (c,d) (d,g) (a,c) (d,h) (e,g) (g,h) (b,f) (f,g) (a,b) (c,f)

for each edge (u,v) in order of sort
if u and v are connected using previously selected edges
add (u,v)

Kruskal's Algorithm



$$1+2+1+4+3+1+2=14$$



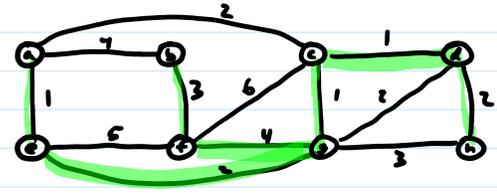
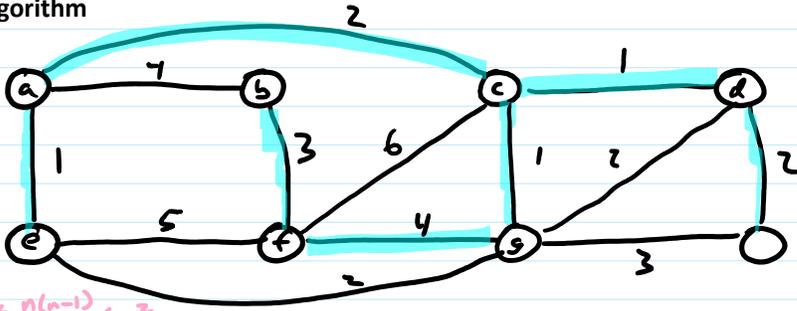
$$1+2+1+4+3+1+2=14$$

sort edges in order of increasing weight (a,e) (c,g) (c,d) (d,g) (a,c) (d,h) (e,g) (g,h) (b,f) (f,g) (a,b) (c,f)

for each edge (u,v) in order of sort

if u and v are connected using previously selected edges
add (u,v)

Kruskal's Algorithm



$$n-1 \leq m \leq \frac{n(n-1)}{2} \leq n^2$$

$$\log n-1 \leq \log m \leq 2 \log n \quad 1+2+1+4+3+1+2=14$$

$$\log m \in \Theta(\log n)$$

$$1+2+1+4+3+1+2=14$$

PRE: input is a connected, undirected graph with positive weights

POST: output is an MST

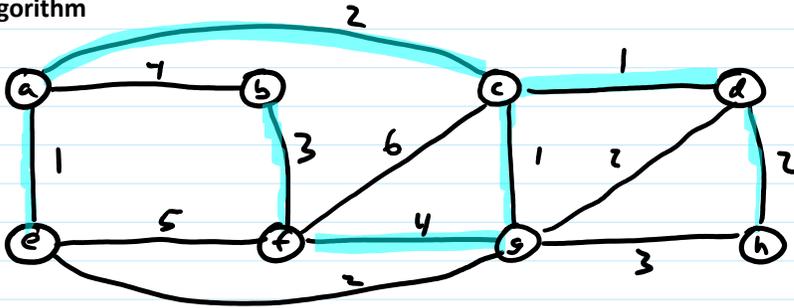
sort edges in order of increasing weight (a,e) (c,g) (c,d) (d,g) (a,c) (d,h) (e,g) (g,h) (b,f) (f,g) (a,b) (c,f)

→ worst case $\Theta(m \log n)$
 for each edge (u,v) in order of sort worst case m iterations

if u and v are connected using previously selected edges BFS/DFS: worst case $\Theta(n)$
 add (u,v) ($\leq n$ edges selected)

total: worst-case $\Theta(nm)$

Kruskal's Algorithm



connected components

e a c g d h b f

initialize (V)

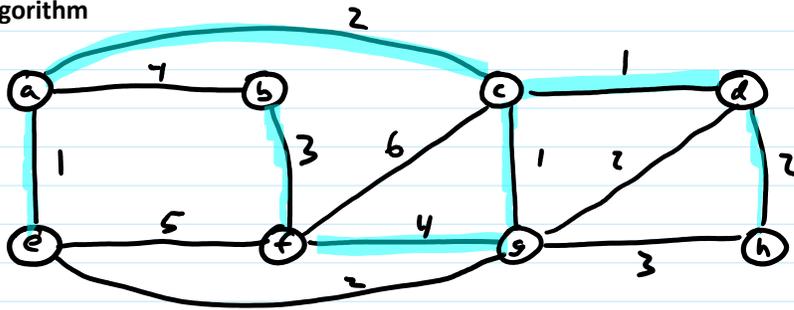
sort edges in order of increasing weight (a,e) (c,g) (c,d) (d,g) (a,c) (d,h) (e,g) (g,h) (b,f) (f,g) (a,b) (c,f)

for each edge (u,v) in order of sort
 if find-set(u) ≠ find-set(v)
 add (u,v)
 union(u,v)

worst-case $\Theta(m \log n)$

UNION/FIND ADT (or disjoint set, or partition) :
 initialize(S): make singletons for each set S
 union(u,v): merge sets containing u, v
 find-set(u): find representative of set containing u

Kruskal's Algorithm



connected components

e a c g d h b f

sort edges in order of increasing weight (a,e) (c,g) (c,d) (d,g) (a,c) (d,h) (e,g) (g,h) (b,f) (f,g) (a,b) (c,f)

for each edge (u,v) in order of sort
if $\text{find-set}(u) \neq \text{find-set}(v)$
 add (u,v)
 union (u,v)