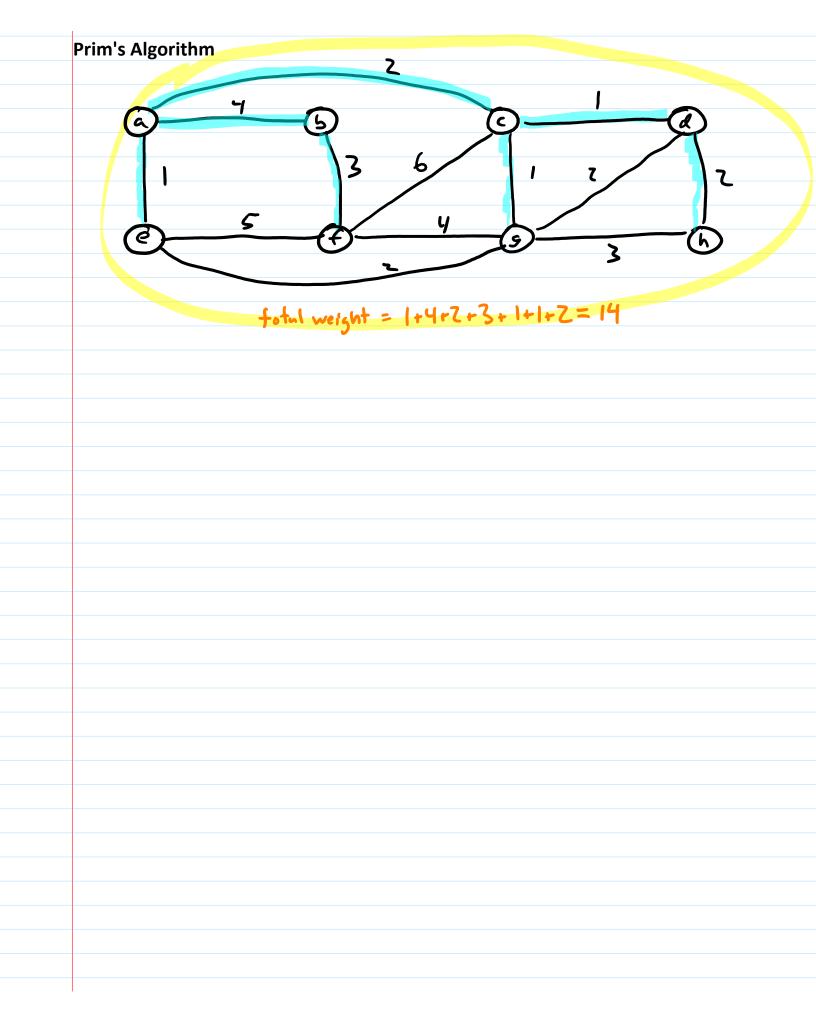
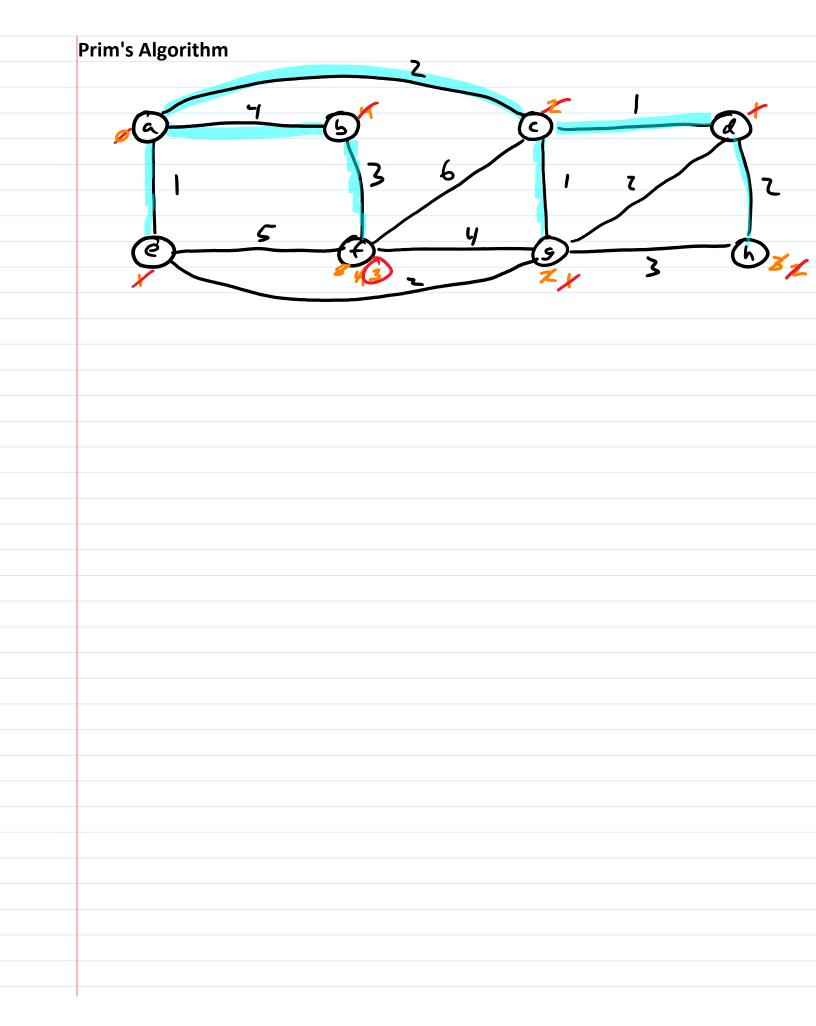
Kruskal's Algorithm 2 connected components C (a ea 3 6 τ 6 cd g y (e h sort edges in order of increasing weight (a,e)(c,g) (c,d) (d,g) (a,c) (d h) (e, ) (g,h) (b,f) (f,g) (a,b) (c,f) for each edge (u,v) in order of sort if find-set(u) ≠ find-set(v) add (u,v) union(u,v)



Prim's Algorithm 2 I C Ĺ 9 d 3 6 7 ٢ e 9 h 3 5  $\binom{\min}{(u,v)}$  wlusb),  $\binom{\min}{(u,v)}$  wlusf),  $\binom{\min}{(u,v)}$  wlusd),  $\binom{\min}{(u,v)}$  wlush) ues ues ues w(n,v)min Ŀ min (u,v) uts V¢s  $\min_{\substack{(u,f)\\u\in S}} \omega(u,f) = \min_{\substack{u\in S}} \omega(u,f) = \max_{\substack{u\in S}} \omega(u,f) =$ min wlu,f), w(g,f) (u,f) ue {u,e,c}

Prim's Algorithm 2 1 5 C d [a 3 6 ٢ 1 Z 5 Y e 9 h 3 5 min w(u,b), (u,v) w(u,f), (u,v) w(u,d), (u,v) (u,v) ues ues ues ues wlu,h) min (u,v) uts V#s w(n,v)min Ľ min w(u,f) = min(u,f) $u \in S$ min w (u,f) ue{u,e,c} w(g,f wlu,f



Dijkstra's Algorithm PRE: no negative weight edges Prim POST: T is a minimum spanning tree for each v  $-color[v], pred[v], d[v] \leftarrow IN QUEUE, NIL, \infty$ d[s] ← 0 s ⊢ Ø using adjacency list + unsarted array worst case  $\Theta(n^2)$ binary heap worst case  $\Theta(m \log n)$ Fiberneci heap worst case  $O(m + n \log n)$ Q ~ new PriorityQueue(d) while Q not empty degreve amortized O (lug n) change priority americad O(1) to make priority ( americad O(1)  $u \leftarrow dequeue(Q)$  $S \leftarrow S \cup \{u\}$ for each outneighbor v of u if color[v] = IN\_QUEUE and  $d[v] > \frac{d[u]}{d[u]} + w(u, v)$ change priority(Q, v, d[u] + w(u, v))  $d[v] \leftarrow \overline{a[u]} + w(u, v)$  $pred[v] \leftarrow u$  $color[u] \leftarrow DONE$  $T \leftarrow \{(u, pred[u]) \mid pred[u] \neq NIL\}$