

Priority Queue

a collection of items identified by a key and having an associated priority

	0	1	2	3	4	5
key	BWI	MEX	DEL	JNB	AKL	MNL
priority	20	12	30	40	25	15

enqueue(MNL, 15)

array
amortized
 $O(1)$

enqueue (key, pri)
add a new item with
given key and priority

dequeue ()
remove and return the item
with lowest priority value

change-priority (key, pri)
change the priority of
the item with the given key

Priority Queue

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key	priority	key	index
BWI	20	BWI	0
MNL	15	DEL	2
DEL	30	JNB	3
JNB	10	AKL	4
AKL	25	MNL	1

change-priority (JNB, 10)

enqueue (key, pri)
add a new item with
given key and priority

array
amortized
 $O(1)$

dequeue ()
remove and return the item
with lowest priority value

$\theta(n)$

change-priority (key, pri)
change the priority of
the item with the given key

$O(1)$

Priority Queue

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key	priority	index
0	20	0
1	15	1
2	30	2
3	40 10	3
4	25	4

change-priority (2 , 10)

enqueue (key, pri)
add a new item with
given key and priority

array
amortized
 $O(1)$

dequeue ()
remove and return the item
with lowest priority value

$\theta(n)$

change-priority (key, pri)
change the priority of
the item with the given key

$O(1)$

Priority Queue

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key
priority

0	1	2	3
BWI	MNL	DEL	AKL
20	15	30	25

key index
BWI 0
DEL 2
~~JNB 3~~
AKL 3
MNL 1

array
amortized
 $O(1)$

sorted array

enqueue (key, pri)
add a new item with
given key and priority

dequeue ()
remove and return the item
with lowest priority value

change-priority (key, pri)
change the priority of
the item with the given key

$\theta(n)$

$O(1)$

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key	priority	key	index
MNL	15	BWI	2
BWI	20	DEL	4
AKL	25	AKL	3
DEL	30	MNL	0
		SEA	1

enqueue(SEA, 18)

array
amortized
 $O(1)$

sorted array
worst case
 $\Theta(n)$

enqueue (key, pri)
add a new item with
given key and priority

dequeue ()
remove and return the item
with lowest priority value

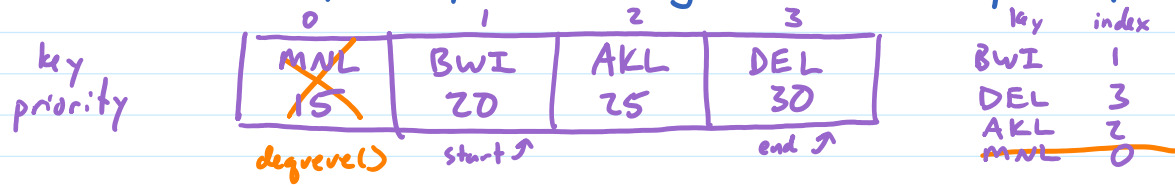
change-priority (key, pri)
change the priority of
the item with the given key

$\Theta(n)$

$O(1)$

Priority Queue

a collection of items identified by a key and having an associated priority



enqueue (key, pri)
 add a new item with given key and priority

array
 amortized
 $O(1)$

sorted array
 worst case
 $\Theta(n)$

dequeve ()
 remove and return the item with lowest priority value

$\Theta(n)$

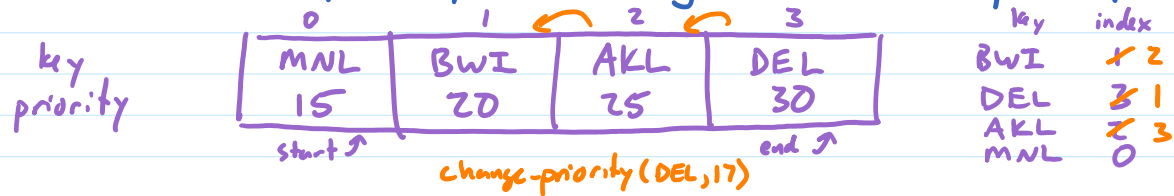
$O(1)$

change-priority (key, pri)
 change the priority of the item with the given key

$O(1)$

Priority Queue

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enqueue (key, pri)
add a new item with given key and priority

array
amortized
 $O(1)$

sorted array
worst case
 $\Theta(n)$

dequeue ()
remove and return the item with lowest priority value

$\Theta(n)$

$O(1)$

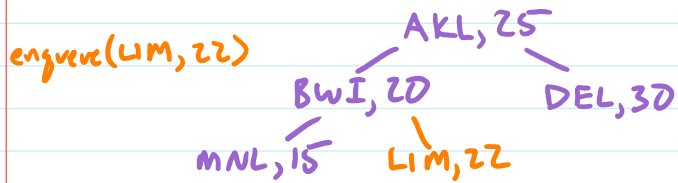
change-priority (key, pri)
change the priority of the item with the given key

$O(1)$

worst case
 $\Theta(n)$

Priority Queue

a collection of items identified by a key and having an associated priority



enqueue (key, pri)
add a new item with
given key and priority

array
amortized
 $O(1)$

sorted array
worst case
 $\Theta(n)$

balanced
binary search tree
worst case
 $\Theta(\log n)$

dequeue ()
remove and return the item
with lowest priority value

$\Theta(n)$

$O(1)$

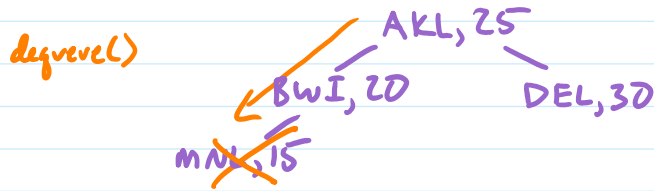
change-priority (key, pri)
change the priority of
the item with the given key

$O(1)$

worst case
 $\Theta(n)$

Priority Queue

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enqueue (key, pri)
add a new item with
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array
amortized
 $O(1)$

sorted array
worst case
 $\Theta(n)$

balanced
binary search tree
worst case
 $\Theta(\log n)$

dequeue ()
remove and return the item
with lowest priority value

$\Theta(n)$

$O(1)$

worst case
 $\Theta(\log n)$

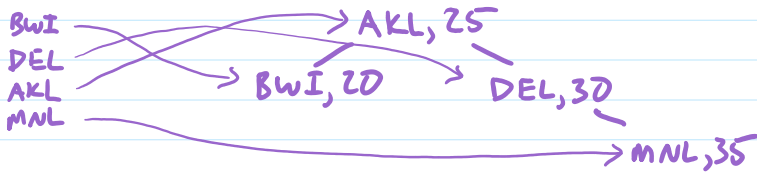
change-priority (key, pri)
change the priority of
the item with the given key

$O(1)$

worst case
 $\Theta(n)$

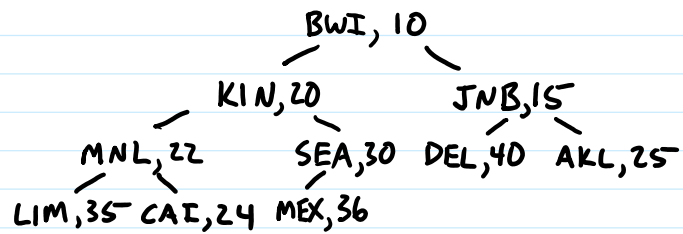
Priority Queue

a collection of items identified by a key and having an associated priority



	array amortized $O(1)$	sorted array worst case $\Theta(n)$	balanced binary search tree worst case $\Theta(\log n)$	heap worst case $\Theta(\log n)$
change-priority (MNL, 35) enqueue (key, pri) add a new item with given key and priority	$O(1)$	$\Theta(n)$	$\Theta(\log n)$	$\Theta(\log n)$ →
dequeue () remove and return the item with lowest priority value	$\Theta(n)$	$O(1)$	worst case $\Theta(\log n)$	worst case $\Theta(\log n)$ →
change-priority (key, pri) change the priority of the item with the given key	$O(1)$	worst case $\Theta(n)$	worst case $\Theta(\log n)$	worst case $\Theta(\log n)$ →
TOTAL n enqueues/dequeues + m change-priority operations	$O(n^2 + m)$ $m \in \Theta(n)$ → $O(n^2)$ $m \in \Theta(n^2)$ → $O(n^2)$	$O(n^2 + nm)$ $O(n^2)$ $O(n^3)$	$O((nm) \log n)$ $O(n \log n)$ $O(n^2 \log n)$	

Heap

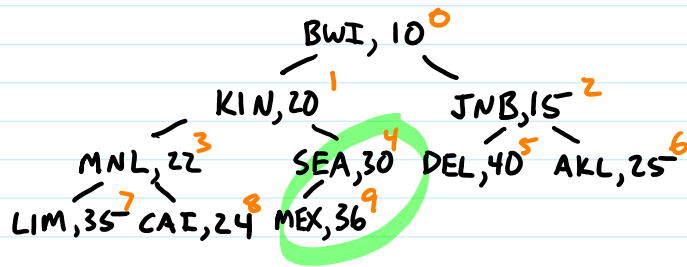


Heap: a binary tree such that

order the value in a node is \leq the value in any children (min-heap)

shape all levels but possibly the last are full and nodes are as far left as possible

Heap

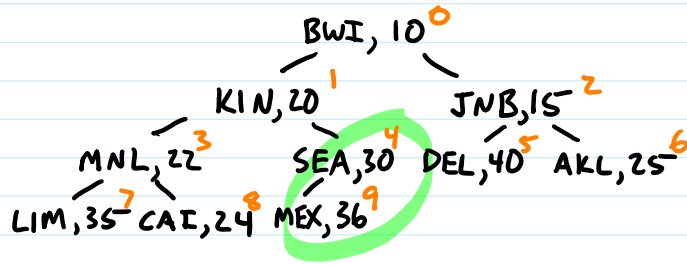


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Heap

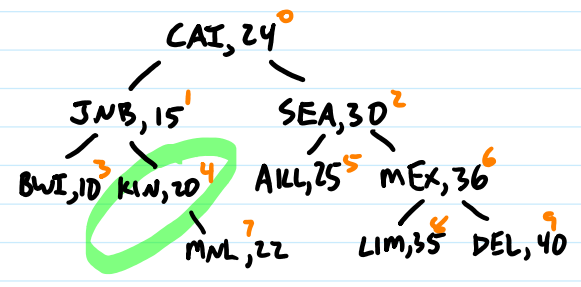


$LEFT(i) = 2i + 1$
 $RIGHT(i) = 2i + 2$

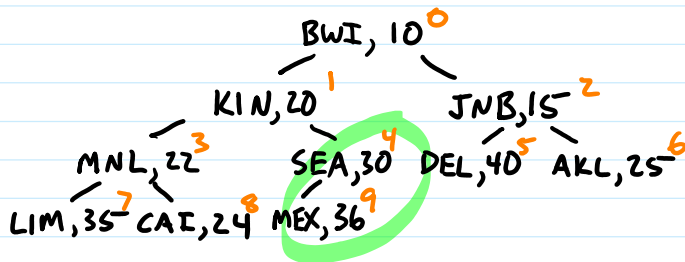
0	1	2	3	4	5	6	7	8	9
BWI	KIN	JNB	MNL	SEA	DEL	AKL	LIM	CAI	MEX
10	20	15	22	30	40	25	35	24	36

Heap: a binary tree such that
 order the value in a node is \leq the value in any children

shape all levels but possibly the last are full and nodes are as far left as possible



Heap



$LEFT(i) = 2i + 1$
 $RIGHT(i) = 2i + 2$
 $PARENT(i) = \lfloor \frac{i-1}{2} \rfloor$

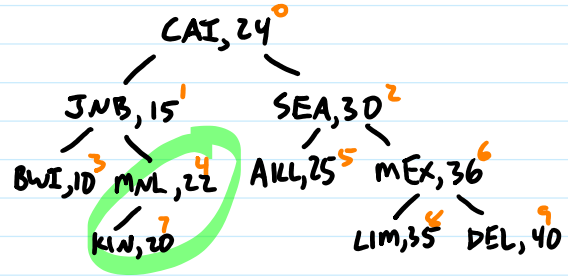
- AKL 6
- BWI 0
- DEL 5
- SEA 4
- LIM 7
- MNL 3
- CAI 8
- MEX 9
- KIN 1
- JNB 2

0	1	2	3	4	5	6	7	8	9
BWI	KIN	JNB	MNL	SEA	DEL	AKL	LIM	CAI	MEX
10	20	15	22	30	40	25	35	24	36

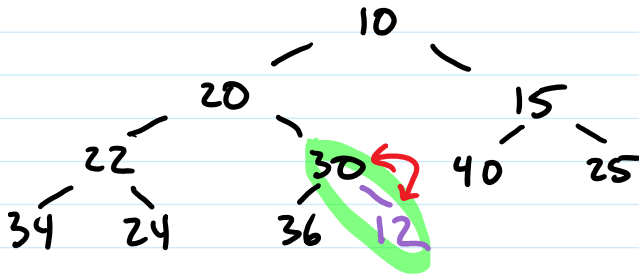
Heap: a binary tree such that

order the value in a node is \leq the value in any children

shape all levels but possibly the last are full and nodes are as far left as possible

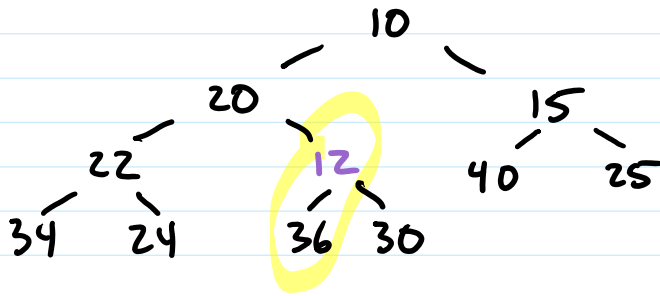


Heap



enqueue(key, pri):
 $i \leftarrow n$
 add (key, pri) at location i
 while priority at i < priority at PARENT(i)
 swmp i , PARENT(i)

Heap



enqueue(key, pri):

$i \leftarrow n$

add (key, pri) at location i

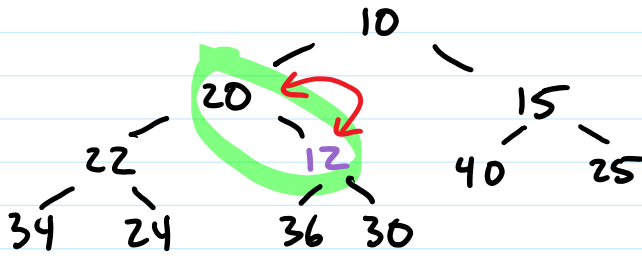
while

priority at $i <$ priority at PARENT(i)

swap i , PARENT(i)

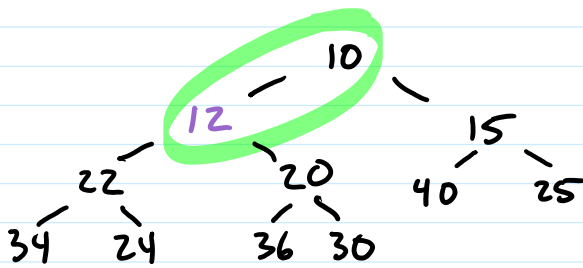
$i \leftarrow$ PARENT(i)

Heap



enqueue(key, pri):
 $i \leftarrow n$
 add (key, pri) at location i
 while $\text{priority at } i < \text{priority at PARENT}(i)$
 ~~swap~~ $i, \text{PARENT}(i)$
 $i \leftarrow \text{PARENT}(i)$

Heap

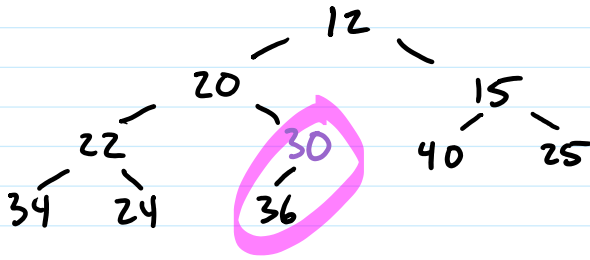


enqueue(key, pri):
worst case $\Theta(\log n)$
 ≤ 1 iteration per level
 $\Theta(\log n)$ levels

$i \leftarrow n$
add (key, pri) at location i
while $i > 0$ and priority at i < priority at PARENT(i)
 swap i , PARENT(i)
 $i \leftarrow$ PARENT(i)
 $n \leftarrow n + 1$

} $O(1)$ per iteration

Heap

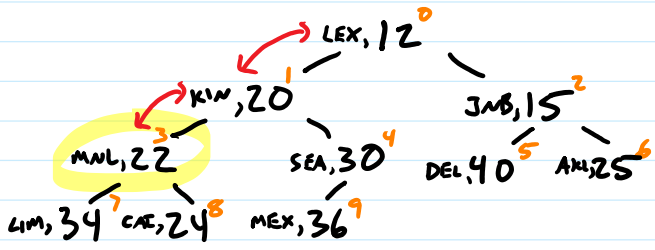


`dequeue()`: remember item at root
worst case $\Theta(\log n)$ move last node to root
 $n \leftarrow n-1$
 $i \leftarrow 0$

worst case $\Theta(\log n \text{ iterations})$ while $\text{LEFT}(i) < n$ and priority at $i >$ priority at one of i 's children
swap i with smallest child
 $i \leftarrow$ former index of smallest child $O(1)$ per iteration

change-priority: find index i of item to change
use loop from dequeue if priority value \uparrow , loop from enqueue if \downarrow

Heap



KIN	1
JNB	2
AKL	6
MEX	9
CAI	8
MNL	3
DEL	5
LIM	7
SEA	4
LEX	0

change-priority (MNL, 5)

dequeue(): remember item at root
 worst case $\Theta(\log n)$ move last node to root
 $n \leftarrow n - 1$
 $i \leftarrow 0$

worst case $\Theta(\log n)$ iterations) while LEFT(i) < n and priority at i > priority at one of i's children
 swap i with smallest child
 $i \leftarrow$ former index of smallest child
 $O(1)$ per iteration

change-priority: find index i of item to change
 worst case $\Theta(\log n)$ use loop from dequeue if priority value \uparrow , loop from enqueue if \downarrow

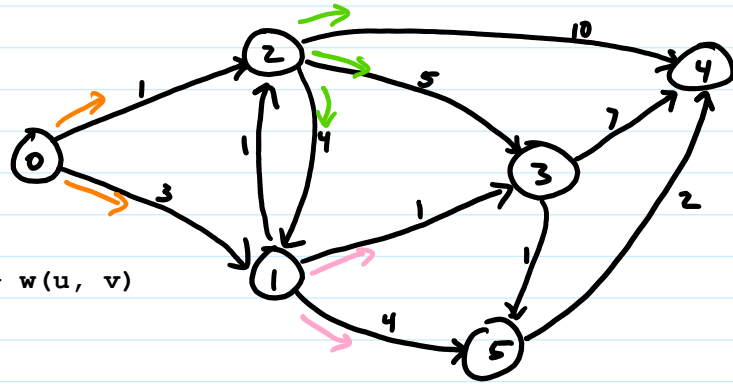
Dijkstra's Algorithm

```

for each v
  color[v], pred[v], d[v] ← IN_QUEUE, NIL, ∞
d[s] ← 0

Q ← new PriorityQueue(d)

while Q not empty
  u ← dequeue(Q)
  for each outneighbor v of u
    if color[v] = IN_QUEUE and d[v] > d[u] + w(u, v)
      change_priority(Q, v, d[u] + w(u, v))
      d[v] ← d[u] + w(u, v)
      pred[v] ← u
  color[u] ← DONE
  
```



vertex	0	1	2	3	4	5
priority(u)	0	∞ 3	∞ 1	∞ 6	∞ 11	∞ 7
predecessor		0	0	2	2	1