

CPSC 427a: Object-Oriented Programming

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Casts and Conversions

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Casts and Conversions

Casts in C

A C cast changes an expression of one type into another.

Examples:

```
int x;  
unsigned u;  
double d;  
int* p;  
  
(double)x;    // type double; preserves semantics  
(int)u;       // type unsigned; possible loss of information  
(unsigned)d;  // type unsigned; big loss of information  
(long int)p;  // type long int; violates semantics  
(double*)p;   // preserves pointeriness but violates semantics
```

Different kinds of casts

C uses the same syntax for different kinds of casts.

Value casts convert from one representation to another, partially preserving semantics. Often called *conversions*.

- ▶ `(double)x` converts integer `x` to equivalent `double` floating point representation.
- ▶ `(short int)x` converts integer `x` to equivalent `short int`, *if the integer falls within the range of a `short int`.*

Pointer casts leave representation alone but change interpretation of pointer.

- ▶ `(double*)p` treats bits at destination of `p` as the representation of a double.

C++ casts

C++ has four kinds of casts.

1. *Static cast* includes value casts of C. Tries to preserve semantics, but not always safe. Applied at compile time.
2. *Dynamic cast* Applies only to pointers and references to objects. Preserves semantics. Applied at run time. [See demo [18a-Dynamic_cast.](#)]
3. *Reinterpret cast* is like the C pointer cast. Ignores semantics. Applied at compile time.
4. *Const cast* Allows `const` restriction to be overridden. Applied at compile time.

Explicit cast syntax

C++ supports three syntax patterns for explicit casts.

1. C-style: `(double)p`.
2. Functional notation: `double(x); myObject(10);`.
(Note the similarity to a constructor call.)
3. Cast notation:

```
int x; myBase* b; const int c;  
▶ static_cast<double>(x);  
▶ dynamic_cast<myDerived*>(b);  
▶ reinterpret_cast<int*>(p);  
▶ const_cast<int>(c);
```

Implicit casts

General rule for implicit casts: If a type **A** expression appears in a context where a type **B** expression is needed, use a semantically safe cast to convert from **A** to **B**.

Examples:

▶ Assignment: `int x; double d; x=d; d=x;`

▶ Pointer assignment:

```
class A { ... };  
class B : public A { ... };  
A* ap; B* bp; ap = bp;
```

▶ Initialization:

`A a=x;` converts `x` to an `A`, then copies.

▶ Construction:

`A a(x);` calls `A` constructor, possibly casting `x`.

Ambiguity

Can be more than one way to cast from **B** to **A**.

```
class B;
class A { public:
    A(){}
    A(B& b) { cout<< "constructed A from B\n"; }
};
class B { public:
    A a;
    operator A() { cout<<"casting B to A\n"; return a; }
};
int main() {
    A a; B b;
    a=b;
}
```

error: conversion from 'B' to 'const A' is ambiguous

explicit keyword

Not always desirable for constructor to be called implicitly.

Use `explicit` keyword to inhibit implicit calls.

Previous example compiles fine with use of `explicit`:

```
class B;
class A {
public
    A(){}
    explicit A(B& b) { cout<< "constructed A from B\n"; }
};
...
```

Question: Why was an explicit definition of the default constructor not needed?

Operator Extensions

How to define operator extensions

Unary operator `op` is shorthand for `operator op ()`.

Binary operator `op` is shorthand for `operator op (T arg2)`.

Some exceptions: Pre-increment and post-increment.

To define meaning of `++x` on type `T`, define `operator ++()`.

To define meaning of `x++` on type `T`, define `operator ++(int)` (a function of one argument). The argument is ignored.

Other special cases

Some special cases.

- ▶ Subscript: `T& operator [] (S index)`.
- ▶ Arrow: `X* operator ->()` returns pointer to a class `X` to which the selector is then applied.
- ▶ Function call; `T2 operator () (arg list)`.
- ▶ Cast: `operator T()` defines a cast to type `T`.

Can also extend the `new`, `delete`, and `,` (comma) operators.

Virtue Demo

Virtual virtue

```
class Basic {
public:
    virtual void print(){cout <<"I am basic.  "; }
};
class Virtue : public Basic {
public:
    virtual void print(){cout <<"I have virtue.  "; }
};
class Question : public Virtue {
public:
    void print(){cout <<"I am questing.  "; }
};
```

Main virtue

What does this do?

```
int main (void) {
    cout << "Searching for Virtue\n";
    Basic* array[3];
    array[0] = new Basic();
    array[1] = new Virtue();
    array[2] = new Question();
    array[0]->print();
    array[1]->print();
    array[2]->print();
    return 0;
}
```

See demo [18b-Virtue!](#)

Linear Data Structure Demo

Using polymorphism

Similar data structures:

- ▶ Linked list implementation of a stack of items.
- ▶ Linked list implementation of a queue of items.

Both support a common **interface**:

- ▶ `void push(Item*)`
- ▶ `Item* pop()`
- ▶ `Item* peek()`
- ▶ `ostream& print(ostream&)`

They differ only in where `push()` places a new item.

The demo [18c-Virtual](#) (from Chapter 15 of textbook) shows how to exploit this commonality.

Interface file

We define this common interface by the abstract class.

```
class Container {
    public:
        virtual void    put(Item*)      =0;
        virtual Item*  pop()            =0;
        virtual Item*  peek()           =0;
        virtual ostream& print(ostream&) =0;
};
```

Any class derived from it is required to implement these four functions.

We could derive `Stack` and `Queue` directly from `Container`, but we instead exploit even more commonality between these two classes.

Class Linear

```
class Linear: public Container {
protected:
    Cell* head;
private:
    Cell* here; Cell* prior;
protected:
    Linear();
    virtual ~Linear ();
                void reset();
                bool end() const;
                void operator ++();
    virtual void insert( Cell* cp );
    virtual void focus() = 0;
                Cell* remove();
                void setPrior(Cell* cp);
public:
    void put(Item * ep);
                Item* pop();
                Item* peek();
    virtual ostream& print( ostream& out );
};
```

Example: Stack

```
class Stack : public Linear {
public:
    Stack(){}
    ~Stack(){}
    void insert( Cell* cp ) { reset(); Linear::insert(cp); }
    void focus(){ reset(); }

    ostream& print( ostream& out ){
        out << " The stack contains:\n";
        return Linear::print( out );
    }
};
```

Example: Queue

```
class Queue : public Linear {
private:
    Cell*   tail;

public:
    Queue() { tail = head; }
    ~Queue(){ }

    void insert( Cell* cp ) {
        setPrior(tail); Linear::insert(cp); tail=cp; }
    void focus(){ reset(); }
};
```

Class structure

Class structure.

- ▶ `Container` specifies the common interface.
- ▶ `Linear` contains the bulk of the code. It is derived from `Container`.
- ▶ `Stack` and `Queue` are both derived from `Linear`.
- ▶ `Cell` is a “helper” class that is aggregated by `Linear`.
- ▶ `Item` is the base type for the container elements. It is defined by a `typedef` here but would normally be specified by a template.
- ▶ `Exam` is a non-trivial item type used by `main` to illustrate stacks and queues.

C++ features

The demo illustrates several C++ features.

1. [Container] Pure abstract class.
2. [Cell] Friend functions.
3. [Cell] Printing a pointer in hex.
4. [Cell] Operator extension `operator Item*()`.
5. [Linear] Virtual functions and polymorphism.
6. [Linear] Scanner pairs (prior, here) for traversing a linked list.
7. [Linear] Operator extension `operator ++()`
8. [Linear, Exam] Use of `private`, `protected`, and `public` in same class.

#include structure

Getting `#include`'s in the right order.

Problem: Making sure compiler sees symbol definitions before they are used.

Partial solution: Make dependency graph. If not cyclic, each `.hpp` file includes the `.hpp` files just above it.

