

# CPSC 427: Object-Oriented Programming

Michael J. Fischer

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## Polymorphic Derivation

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## Some uses for derived classes.

- ▶ **Code reuse.** A base class can contain one copy of code that is be used by several derived variants through inheritance.
- ▶ **Modularity.** The functionality provided by a base class can be extended in a derived class. Example: `BSquare` extends `Square` by adding board coordinates and clusters.
- ▶ **Generic programming and isolation.** Demo 17-Craps-extended contains a simulator for the gambling game “craps” that can use different dice implementations.
- ▶ **Polymorphic collections.** A company has different kinds of employees with different rules for calculating their pay, each represented by a derived class with its own `calculatePay` function appropriate to that kind of employee.

## Type Hierarchies

Consider following simple type hierarchy:

```
class B      { public: int f(); ... };  
class U : B { int f(); ... };  
class V : B { int f(); ... };
```

We have a base class **B** and derived classes **U** and **V**.  
A different method **f()** is defined in each.

Relationships: A **U** is a **B** (and more). A **V** is a **B** (and more).

A **U** can be used wherever a **B** is expected.

Example: Definition **f(B& x) ...** ; call **U z; f(z);**

Inside of **f()**, only the **B**-part of **z** is visible. This is called **slicing**.

## Pointers and slicing

Declare `B* bp; U* up = new U; V* vp = new V.`

Can write `bp = up;` or `bp = vp;`.

Why does this make sense?

- ▶ `*up` has an embedded instance of `B`.
- ▶ `*vp` has an embedded instance of `B`.

If `bp = up`, then `bp` points to the embedded `B`-instance of object `*up`. The rest of `*up` is inaccessible because of object slicing.

## Ordinary derivation

In our previous example

```
class B      { public: int f(); ... };  
class U : B { int f(); ... };  
class V : B { int f(); ... };  
B* bp;
```

`bp` can point to objects of type `B`, type `U`, or type `V`.

Want `bp->f()` to refer to `U::f()` if `bp` points to a `U` object.

Want `bp->f()` to refer to `V::f()` if `bp` points to a `V` object.

However, with ordinary derivation, `bp->f()` always refers to `B::f()`.

## Polymorphic derivation

The keyword `virtual` allows for polymorphic derivation.

```
class B      { public: virtual int f(); ... };  
class U : B { virtual int f(); ... };  
class V : B { virtual int f(); ... };  
B* bp;
```

A virtual function is dispatched at run time to the class of the actual object.

`bp->f()` refers to `U::f()` if `bp` points to a `U`.

`bp->f()` refers to `V::f()` if `bp` points to a `V`.

`bp->f()` refers to `B::f()` if `bp` points to a `B`.

Here, the type refers to the allocation type.



## Unions and type tags

We can regard `bp` as a pointer to the union of types `B`, `U` and `V`.

To know which of `B::f()`, `U::f()` or `V::f()` to use for the call `bp->f()` requires runtime **type tags**.

If a class has **virtual** functions, the compiler adds a type tag field to each object.

This takes space at run time.

The compiler also generates a **vtable** to use in dispatching calls on virtual functions.

## Virtual destructors

Consider `delete bp;`, where `bp` points to a `U` but has type `B*`.

The `U` destructor will *not* be called unless destructor `B::~~B()` is declared to be `virtual`.

Note: The base class destructor is always called, *whether or not it is virtual*.

In this way, destructors are different from other member methods.

**Conclusion:** If a derived class has a non-empty destructor, the *base class* destructor should be declared `virtual`.

## Uses of polymorphism

Some uses of polymorphism:

- ▶ To define an extensible set of representations for a class.
- ▶ To allow containers to store mixtures of different but related types of objects.
- ▶ To support run-time variability of within a restricted set of related types.

## Multiple representations

Might want different representations for an object.

Example: A point in the plane can be represented by either Cartesian or Polar coordinates.

A `Point` base class can provide abstract operations on points. E.g., `virtual int quadrant() const` returns the quadrant of `*this`.

For Cartesian coordinates, quadrant is determined by the signs of the  $x$  and  $y$  coordinates of the point.

For polar coordinates, quadrant is determined by the angle  $\theta$ .

Both `Cartesian` and `Polar` derived classes should contain a method for `int quadrant() const`.

## Heterogeneous containers

One might wish to have a stack of `Point` objects.

The element type of the stack would be `Point*`.

The actual values would have type either `Cartesian*` or `Polar*`.

The automatically generated type tags and dynamic dispatching obviates the need to cast the result of `pop()` to the correct type.

Example:

```
Stack st; Point* p;  
p = st.pop(); // no need to cast result  
p->quadrant(); // automatic dispatch
```

## Run-time variability

Two types are closely related; differ only slightly.

Example: Company has several different kinds of employees.

- ▶ `Employee` base class has a large and complicated payroll function.
- ▶ Payroll is same for all kinds of employees except for a function `pay()` that computes the actual weekly pay.
- ▶ Each employee kind has its own `pay()` function.
- ▶ Big payroll function is in base class.
- ▶ It calls `pay()` to get the actual pay for this `Employee`.

## Pure virtual functions

Suppose we don't want `B::f()` and never create instances of `B`. We make `B::f()` into a **pure virtual function** by writing `=0`.

```
class B      { public: virtual int f()=0; ... };  
class U : B { virtual int f(); ... };  
class V : B { virtual int f(); ... };  
B* bp;
```

A pure virtual function is sometimes called a **promise**. It tells the compiler that a construct like `bp->f()` is legal. The compiler requires every derived class to contain a method `f()`.

## Abstract classes

An **abstract class** is a class with one or more pure virtual functions.

An abstract class cannot be instantiated.

It can only be used as the base for another class.

The destructor can never be a pure virtual function but will generally be **virtual**.

A **pure abstract class** is one where all member functions are pure virtual (except for the destructor) and there are no data members,

Pure abstract classes define an **interface** à la Java.

An interface allows user-supplied code to integrate into a large system.