

Simultaneous Play Games

	R	P	S
Rock	0, 0	-1, 1	1, -1
Paper	1, -1	0, 0	-1, 1
Scissors	-1, 1	1, -1	0, 0

$$A = \begin{pmatrix} 0 & -1 & 1 \\ 1 & 0 & -1 \\ -1 & 1 & 0 \end{pmatrix} \quad B = \begin{pmatrix} 0 & 1 & -1 \\ -1 & 0 & 1 \\ 1 & -1 & 0 \end{pmatrix}$$

a_{ij} = payout for I when I chooses i , II chooses j
 b_{ij} = payout for II

		goal keeper	
		L	R
Penalty Kick	Shooter L	$\frac{1}{10}, \frac{1}{10}$	$\frac{1}{2}, -\frac{1}{2}$
	R	$\frac{1}{2}, -\frac{1}{2}$	$\frac{1}{10}, \frac{1}{10}$

zero-sum : $a_{ij} + b_{ij} = 0$
 (so only need A)

		blitz	balanced	pass	run
		W	X	Y	Z
off	run	$\frac{3}{4}$	$\frac{1}{2}$	$\frac{3}{4}$	$\frac{3}{4}$
	short pass	$-\frac{1}{10}$	$\frac{1}{3}$	$\frac{1}{4}$	$-\frac{1}{2}$
	long pass	$\frac{1}{2}$	0	$-\frac{1}{10}$	$\frac{1}{2}$

↑ expected change in win probability

constant-sum : $a_{ij} + b_{ij} = C$

const-sum game w/ constant C

- 1) play zero-sum game A-C
- 2) sword player I C

so strategy for constant-sum = strategy for zero-sum

	S	H
Stag	2, 2	0, 1
Hare	1, 0	1, 1

non-constant sum game

	R	P	S	row min
Rock	0	-1	1	-1
Paper	1	0	-1	-1
Scissors	-1	1	0	-1

$v^- = 1$ $v^+ = -1$

v^- = amount I guaranteed to win
 $\max_i \min_j a_{ij}$
 no saddle points

	L	R
L	$\frac{1}{2}$	1
R	1	$\frac{3}{4}$

	W	X	Y	Z	row min
A	$\frac{3}{4}$	$\frac{1}{2}$	$\frac{3}{4}$	$\frac{1}{4}$	-1
B	-1	$\frac{1}{3}$	0	-1	-1
C	$\frac{1}{2}$	0	$\frac{3}{4}$	$\frac{3}{4}$	-1

v^+ = ceiling on what I loses
 $\min_j \max_i a_{ij}$
 $v^- = -1$

saddle point

For any constant-sum game, $v^- \leq v^+$

$\forall i, j \quad \min_j a_{ij} \leq a_{ij}$ (min in row i always \leq any entry in that row)

$\forall j \quad \max_i \min_j a_{ij} \leq \max_i a_{ij}$ (max of smaller things \leq max of larger things)

$\min_j \max_i a_{ij} \leq \min_j \max_i a_{ij}$ (min of smaller things \leq min of larger things)

$\max_i \min_j a_{ij} \leq \min_j \max_i a_{ij}$ (constant w/ respect to j ; rename inner var $j' \rightarrow j$)

$v^- \leq v^+$ (def v^-, v^+)

Saddle point: where neither player has incentive to unilaterally change strategy (equilibrium)
 (i^*, j^*) is saddle point means $a_{i^*j^*} \leq a_{ij^*} \leq a_{i^*j}$

each player chooses a single row/column

A constant-sum game A has a saddle point in pure strategies if and only if $v^- = v^+$

\Rightarrow : Suppose A has a saddle point in pure strategies.

Then $\exists i^*, j^*$ s.t. $\forall i, j \quad a_{i^*j^*} \leq a_{ij^*} \leq a_{i^*j}$

so $\max_i a_{ij^*} \leq a_{i^*j^*} \leq \min_j a_{i^*j}$ (max of small things small; min of big things big)

and $\min_j \max_i a_{ij} \leq \max_i \min_j a_{ij}$ $\min_j \max_i a_{ij} \leq \max_i \min_j a_{ij}$
 (min includes $j=j$; max includes $i=i$)
 (min \leq each term; max \geq each term)

so $v^+ \leq \max_i \min_j a_{ij} \leq \min_j \max_i a_{ij} \leq v^-$ (def v^+, v^- ; transitivity)

but $v^- \leq v^+$ so all \leq are = (squeeze)

\Leftarrow Suppose $v^- = v^+$. Let i^* be the i s.t. $v^- = \max_i \min_j a_{ij}$
 j^* be the j s.t. $v^+ = \min_j \max_i a_{ij}$

$a_{i^*j^*} \geq \min_j a_{i^*j} = v^- = v^+ = \max_i a_{ij} \geq a_{i^*j^*}$
 (min is \leq each term; max is \geq each term)

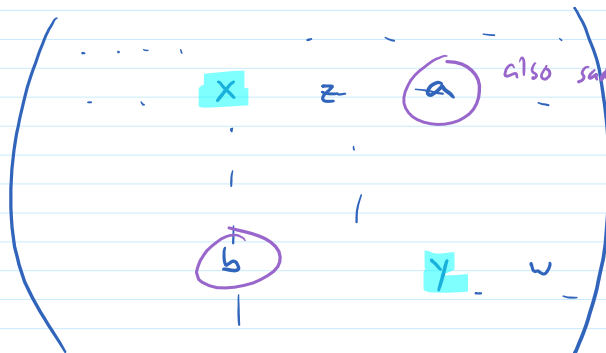
$a_{i^*j^*} \geq v^+ = v^- \geq a_{i^*j^*}$
 and \geq are = (squeeze)

$\therefore v^+ = v^- = a_{i^*j^*}$

$a_{ij} \leq a_{i^*j^*} \leq a_{ij}$ for all i, j
 (substitution)

Suppose there are 2 saddle points in pure strategies (i_1, j_1) and (i_2, j_2)
 with values $a_{i_1j_1} = v_1$ and $a_{i_2j_2} = v_2$

then $v_1 = v_2 = v^- = v^+$
 call this v ,
 value of game



$x \leq a \leq y \leq b \leq x$

since $x = 1$
 all \leq are =

$x = a = y = b = x$
 so
 $x = y$

Mixed Strategies - probability distribution

	R	P	S
Rock	0	-1	1
Paper	1	0	-1
Scissors	-1	1	0

(x_1, \dots, x_n) I plays row i w/ prob x_i

(y_1, \dots, y_m) II play col j w/ prob y_j

guess $X^* = \left(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right) = Y^*$