

Scout ( $p, \alpha, \beta, \text{depth}, h$ )

if  $p$  is terminal then return  $\text{value}(p)$

if  $\text{depth} = 0$  then return  $\text{heuristic}(p)$

if  $p$  is a max position

$\text{best} \leftarrow -\infty$

$b \leftarrow \beta$

for each reachable position  $p'$  and while  $\alpha < \beta$

$\text{curr} \leftarrow \text{Alpha-Beta}(p', \alpha, b, \text{depth}-1, h)$  is  $p' >$  best child so far?  
already null window; don't need Scout null window (except 1st iteration)

false for 1st iteration  
( $b = \beta$ )

if  $b \leq \text{curr} < \beta$

yes; find value of  $p'$

$\text{curr} \leftarrow \text{Scout}(p', b, \beta, \text{depth}-1, h)$

already know value is  $\geq b$

$\text{best} \leftarrow \max(\text{best}, \text{curr})$

$\alpha \leftarrow \max(\alpha, \text{best})$

$\alpha = \text{best}$  unless best is bad ( $<$  original  $\alpha$ )

$b \leftarrow \max(\alpha + 1, \text{best} + 1)$

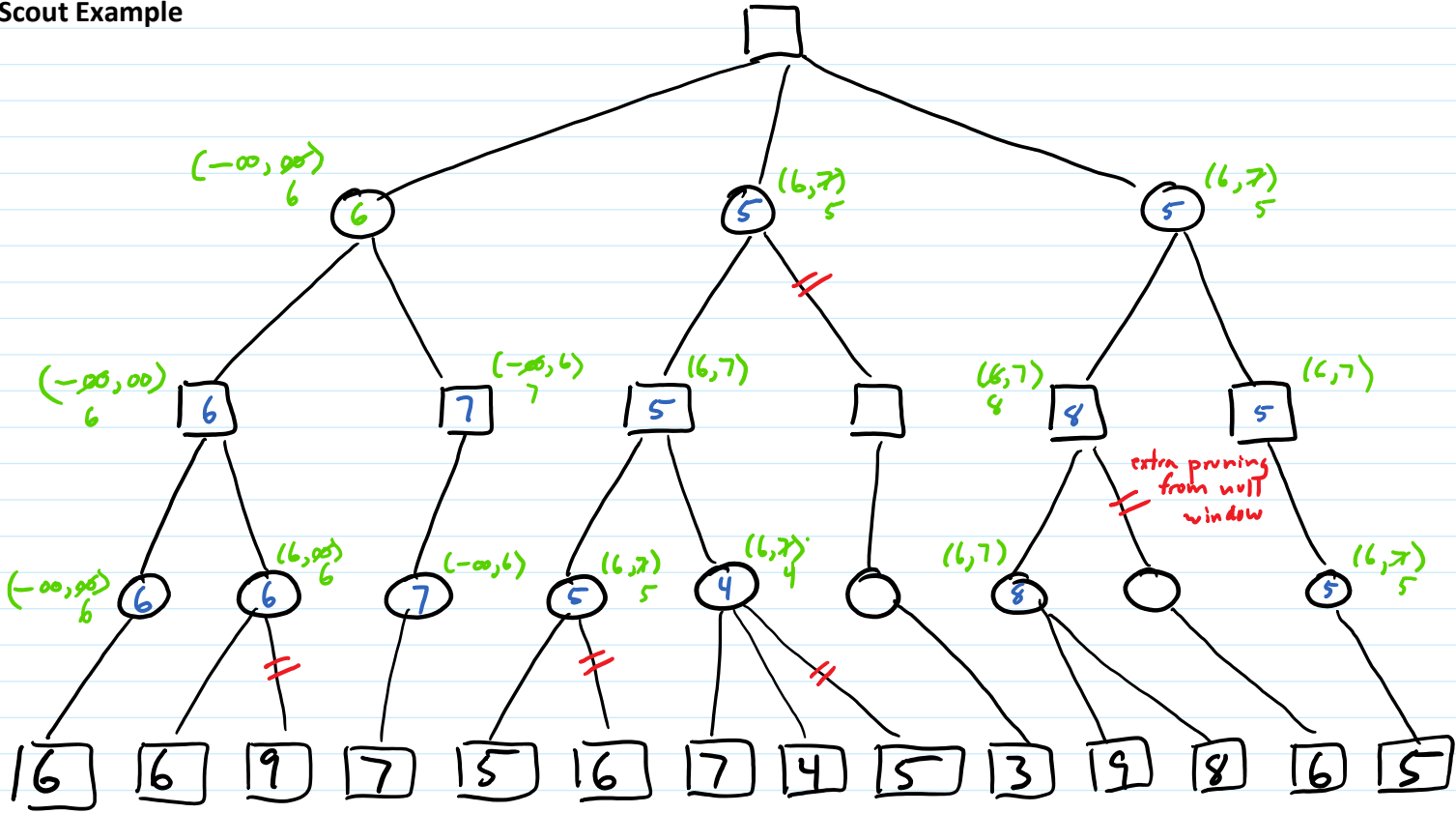
narrow window for next time around

return  $\text{best}$

else

$\vdots$  min position; symmetric  
 $\vdots$

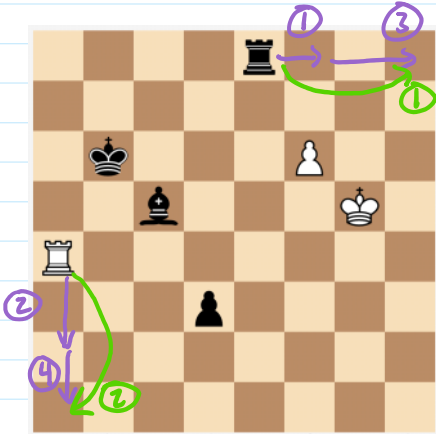
Scout Example





## Transposition Table

Positions may be reachable by multiple sequences of moves



Keep table of values for all positions examined in tree

Keys: positions

Values: (value/bound, move, depth)

lower, upper bounds  
= if exact (given h, depth)

if value/bound from shallower search, ignore

Add check at start of A-B

if pos present and searched depth  $\geq$  desired depth  
if value is exact, return value  
else if upperbound  $< \alpha$  return val  
else if lowerbound and value  $\geq \beta$  return val

Save returned values in table

fixed size - replacement policy

deepest  
largest  
newest  
two-level

# Alpha-Beta Pruning updated for transposition table

Alpha-Beta ( $p, \alpha, \beta, \text{depth}, h, \text{tt}$ ) (returns bounds on value of  $p$ , given  $h$  and  $\text{depth}$ )  
 ↪ (lowerBound, upperBound)

if  $p$  in  $\text{tt}$   
 ( $\text{low}, \text{up}, d$ )  $\leftarrow$   $\text{tt.get}(p)$   
 if  $d \geq \text{depth}$  don't use results of shallow searches  
 if  $\text{low} = \text{up}$  return ( $\text{low}, \text{up}$ ) had exact value (given  $h$  and  $\text{depth}$ )  
 if  $\text{low} \geq \beta$  return ( $\text{low}, \text{up}$ ) may have been result of deeper search  
 if  $\text{up} \leq \alpha$  return ( $\text{low}, \text{up}$ )

if  $p$  is terminal then return value ( $p$ ) assuming this is fast so needn't be stored

if  $\text{depth} = 0$  then  
 value  $\leftarrow h(p)$   
 $\text{tt.put}(p, (\text{value}, \text{value}), \text{depth})$  if heuristic is expensive it makes sense to save it  
 return value exact value for given  $h$

if  $p$  is a max position

bound  $\leftarrow (-\infty, -\infty)$  bounds on value of best child so far

for each position  $p'$  reachable in one move from  $p$  and while  $\alpha < \beta$  cut off if know value of  $p > \beta$   
 bound  $\leftarrow \max(\text{bound}, \text{Alpha-Beta}(p', \alpha, \beta, \text{depth} - 1, h, \text{tt}))$   
 $\alpha \leftarrow \max(\alpha, \text{bound.low})$   
 if not all  $p'$  examined then  $\text{bound.up} = \infty$   
 $\text{tt.put}(p, \text{bound}, \text{depth})$  record bounds on value of  $p$   
 return bound

else ( $p$  is a min position)

⋮ symmetric

(max lower bounds,  
max upper bounds)  
 bound = ( )  
 A-B( $p'$ ) = ( )  
 new best so far = ( )



## MTD-f

pos to eval first guess  
MTD-f (n, f, d)

lowerBound  $\leftarrow -\infty$

upperBound  $\leftarrow \infty$

g  $\leftarrow f$

while lowerBound < upperBound

B  $\leftarrow \max(\text{lowerBound} + 1, g)$

g  $\leftarrow \text{A-B}(n, \text{with TT } B-1, B, d)$

if g < B upperBound  $\leftarrow$  <sup>j-1, g most of the time</sup> g  
else lowerBound  $\leftarrow$  g

return g